





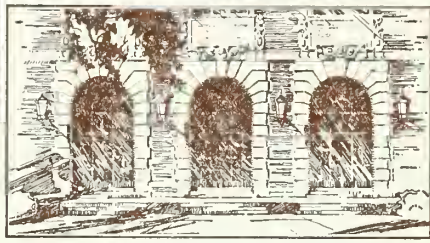
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PROGRAM MANUAL:  
NOR NETWORK TRANSDUCTION BASED ON ERROR-COMPENSATION  
(Reference Manual of NOR Network Transduction  
Programs NETTRA-E1, NETTRA-E2, and NETTRA-E3)

BY

H. C. Lai  
J. N. Culliney

June, 1975



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## ABSTRACT

Three NOR network transduction procedures based on error-compensation were implemented in the FORTRAN computer programs NETTRA-E1, NETTRA-E2, and NETTRA-E3. The general principles on which these programs are based are discussed in a separate report. The present report, however, describes the specific implementations of the three programs and serves as a reference manual for the program user. Preparation of input data is discussed in detail.

Transduction (transformation and reduction) procedures attempt to reduce given, non-optimal, multiple-output, multiple-level, loop-free, NOR-gate networks to "near-optimal" networks of fewer gates. The three programs described in this report, based on the sophisticated "error-compensation" concept, remove gates one at a time from the network and, after each removal, try to reconfigure the network, without employing additional gates, to compensate for any resultant errors caused in the network output(s).





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## 1. INTRODUCTION

This manual is intended to instruct the reader in the use of the FORTRAN programs 'NETTRA-E1,' 'NETTRA-E2,' and 'NETTRA-E3,' and also to enable a moderately detailed understanding of how these programs actually realize their respective algorithms. The principal algorithm upon which these programs are based is described in detail in [5], and this manual will assume a knowledge of the definitions and algorithm descriptions in [4] and [5].

NETTRA-E1, -E2, and -E3 represent only three out of a whole system of programs developed at the University of Illinois by the logical design group of S. Muroga. The generic name 'NETTRA' (for NETWORK TRANSformation) designates the whole collection of programs comprising the system. All of the programs in the NETTRA system either transform or assist in transforming networks of interconnected NOR gates realizing various functions (either completely or incompletely specified) of their respective sets of input variables. By these transformations, a large, non-optimal network of NOR gates realizing one or more various functions can often be reduced to a smaller, less expensive (in terms of the number of required gates and interconnections, for example), near-optimal network realizing the same function(s). In general, such a transduction (transformation and reduction) could involve a complete reorganization of the network: the addition and/or deletion of gates; the addition and/or deletion of connections among gates;

and/or the substitution of certain connections for various others. The transduction procedure realized by NETTRA-E1 (the same procedure forms the basis of NETTRA-E2 and -E3 also) can accomplish any of these changes, with the exception of adding gates to the network.

The present three programs employ a transduction procedure much more powerful than those found in NETTRA-P1, -P2, -PG1, -G1, -G2, -G3, and -G4 (see [1], [2], [3], [6], [7], [8]). The transduction procedures embodied in these earlier programs never make a change to a network which would alter its output functions. However, the current procedure is able to go through a long series of networks representing intermediate stages of the transformation, none of which realize the correct output function(s), in order to finally obtain a less expensive network correctly realizing the desired outputs. In this sense, this procedure is more "far-sighted" than the transduction procedures realized by the earlier programs. Although the feature is not exploited in the programs explained here, it is possible for the current procedure to use an initial network which does not even realize the desired function(s) (hopefully, though, it does realize a function "reasonably close" to the desired one).

NETTRA-E1, -E2, and -E3 are primarily intended to reduce the number of gates in a given network. No serious attempts are made by these programs to minimize the number of connections. However, other transformation programs (e.g., NETTRA-P1, -P2, or -G1) can be applied after NETTRA-E1 (or -E2 or -E3) to try to further reduce the number of connections in a network.

The following section, Section 2, explains the NETTRA-E1 program which applies just once the transduction procedure discussed in [5] to a given network. This single application attempts to eliminate just a single

gate from the network. In order to eliminate several gates, the procedure must be applied several times. Two different methods of doing this (corresponding to the programs NETTRA-E2 and NETTRA-E3) are discussed in Section 3. Section 4 briefly describes all of the subroutines used in the programs NETTRA-E1, -E2 and -E3. In Section 5, instructions are given on the preparation of the input data for the three programs. Finally, a listing of all of the FORTRAN subroutines used in NETTRA-E1, -E2, and -E3 is given in the appendix.



## 2. ERROR-COMPENSATION PROCEDURE

In this section, the NOR-network transduction procedure realized by the FORTRAN program designated NETTRA-E1 is discussed. When it is applied in an attempt to transform a network, the number of gates in the network will either be reduced or left unchanged - it will never be increased.

In contrast to the earlier programs (NETTRA-P1, -P2, -PG1, -G1, -G2, -G3, and -G4) which never transform a network so that it produces incorrect output functions, NETTRA-E1 causes "errors" to appear in the outputs of a network by deliberately removing a necessary (to the correct operation of the network) gate of the network. Then it attempts to compensate for these errors by adding and rearranging connections in the remaining network.

Actually, NETTRA-E1 "memorizes" the original network and removes each of its gates in turn, trying to compensate for the errors in the new networks which each have one less gate than the original.

If the program is successful in compensating for any of these removed gates, it prints out the solution (i.e., the transformed, reduced network of a smaller number of gates than the network originally given as input) and halts.

The input to this program and NETTRA-E2 and -E3 is a description of a particular NOR network under consideration. This description (explained in great detail in Section 5) consists of a set of variables and arrays containing various network parameters.

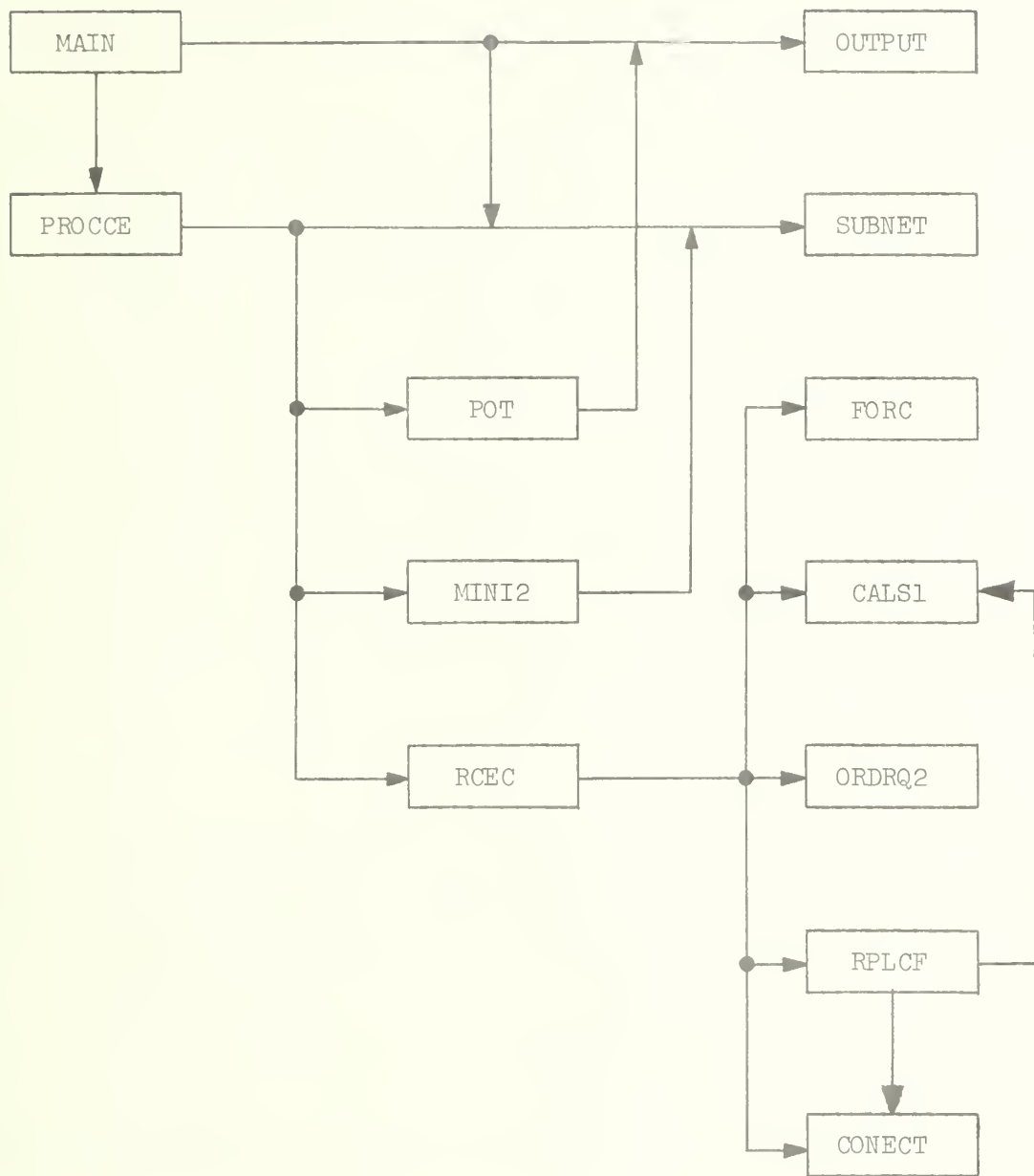


Figure 2.1 General organization of the programs NETTRA-E1 and NETTRA-E2.

The entire NETTRA-El program requires 163K bytes of core storage, about 78K being occupied by the actual program instructions and about 85K by the stored data.

The following subroutines, written in FORTRAN IV for the IBM 360/75, constitute the program NETTRA-El: CALS1, CONECT, FORC, MAIN, MINI2, ORDRQ2, OUTPUT, POT, PROCCE, RCEC, RPLCF, and SUBNET. Two system-supplied timing routines, STIMEZ and KTIMEZ are also assumed to be available, but if they are not, their use can be omitted from the program, or another suitable timing routine substituted, without harming the procedure itself.

The general organization of the program NETTRA-El is shown in Figure 2.1. An arrow from block i to block j represents the fact that the subroutine represented by block i calls the subroutine represented by block j.

## 2.1 General Procedure and Flowchart

The general execution of the error-compensation procedure is carried out by the subroutine PROCCE (for: transduction PROCedure by the Compensation of Errors) which, while quite simple itself, controls the calling of the major subroutines (explained in more detail in Sections 2.2 and 2.3) that actually execute the complex details of the procedure. The following discussion of PROCCE will assume a knowledge of the information contained in [4] and [5].

Explanations of the purposes of the variables and arrays actually appearing in the subroutine can be found in the program listing of PROCCE in the appendix. It is, however, convenient to define some of the variables at this point in order to discuss the flowchart of PROCCE which appears in Figure 2.1.1:



N is the number of external variables,  $n$ , if only uncomplemented variables are allowed as inputs. If both complemented and uncomplemented are available (i.e.,  $n$  variables and their  $n$  complements) then  $N$  is equal to  $2n$ . Note that this is strictly the representation internal to the program; for input-output purposes (as described in Section 5)  $N$  and  $n$  are always equal.

R is the number of gates specified by the input data to the program. It includes all gates declared to be present by the input data, even though some of them may be isolated (i.e., not connected to other gates in the network). Internally, the program represents the gates 1, 2, ...,  $R$  by the labels  $N + 1$ ,  $N + 2$ , ...,  $N + R$ . (External variables are labeled 1, 2, ...,  $N$  internally.)

NR is equal to the sum  $N + R$ . It is often convenient to treat both external variables and gates in a similar manner. External variables being labeled 1, 2, ...,  $N$  and gates being labeled  $N + 1$ , ...,  $N + R$  (internally), the number  $N + R$  is frequently required.

GSMALL is a two-dimensional array used to store intermediate and final calculated compatible sets.<sup>†</sup> GSMALL ( $i$ ,  $j$ ) contains (or rather will contain by the end of a procedure) the  $j$ -th component of the vector representing the compatible set of permissible functions for gate or external variable  $i$ .

---

<sup>†</sup> For simplicity, sometimes just the words "compatible sets" will be used to denote compatible sets of permissible functions.

NEPMAX is a variable limiting the maximum number of "error-positions" which will be tolerated in a network. NEPMAX is either specified by the input data, or, if left unspecified on the input cards, it is set to the value  $2^{(n-1)}$  (for example, in the case of an implicit specification of external variables). An error-position is an index number  $i$  such that at least one output gate of a network (with errors) has an incorrect output for the  $i$ -th network input vector (i.e., the  $i$ -th combination of 0 and 1 assignments to the network inputs). The removal of a gate from the network, as occurs during the execution of the procedure, usually causes errors to appear in several positions. If the number of these error-positions is too great, the chance of compensating all of them is generally low. In the interest of efficiency, if the number of error-positions exceeds NEPMAX after the removal of a certain gate, PROCCE does not attempt to compensate for the errors. Instead, it restores the original network and moves on to remove another gate.

In addition to these variables which appear in the FORTRAN program itself, a few concepts from [4] and [5] should be recalled:

$G_c(v_i)$  was defined in [4], and  $G_E(v_i)$  was defined in [5]. The expression  $G_c(v_i)$  denotes a vector representing a compatible set of permissible functions (CSPF) of a gate or external variable  $v_i$ <sup>†</sup>. This concept

---

<sup>†</sup> In [5], the  $v_i$  represent input terminals and gates, the concept of input terminals being introduced for theoretical completeness. Since the distinction between input terminals and external variables is unnecessary for the purposes of this paper, the  $v_i$  are considered to represent external variables and gates.

was frequently used in programs realizing earlier procedures: NETTRA-PG1, NETTRA-G1, NETTRA-G3, and NETTRA-G4. For the error-compensation procedure, however, CSPF's must be extended to the concept of compatible sets of permissible functions with errors (CSPFE's) as defined in [5]. For each gate or external variable  $v_i$ , the corresponding CSPFE is denoted  $G_E(v_i)$ . The notation  $G_C^{(d)}(v_i)$  or  $G_E^{(d)}(v_i)$  refers to the  $d$ -th component of the respective corresponding vector.

Whereas the components of the CSPF vectors were only of three types: 0, 1, or \* (don't-care), the components of the CSPFE vectors may be any of five (logical) types: \*, 0, 1, 0 (a 0 error), or 1 (a 1 error).

Blocks 1 through 4 of the flowchart of PROCCE (Figure 2.1.1) perform some preliminary steps in preparation for the main part of the procedure realized in blocks 5 through 20.

Block 1 calls the subroutine MINI2 (described in detail in [7]). This serves two purposes. First of all, MINI2 realizes a "pruning" transduction procedure, and it may be able to quickly eliminate some unnecessary gates from the original network. Secondly, MINI2 will calculate CSPF vectors ( $G_v(v_i)$ 's) for all of the gates remaining in the network.

The information is examined in block 2 where the number of 1's in the CSPF vector of each gate is determined.

Block 3 creates an ordering of gates 1 through R based on an increasing number of 1's in their respective CSPF vectors. The ordering is stored in the array PORDER such that the gate stored in the location PORDER(1) has a minimum number of 1's in its CSPF vector and the gate stored in PORDER(R) has a maximum number of 1's. In a rough sense, the number of

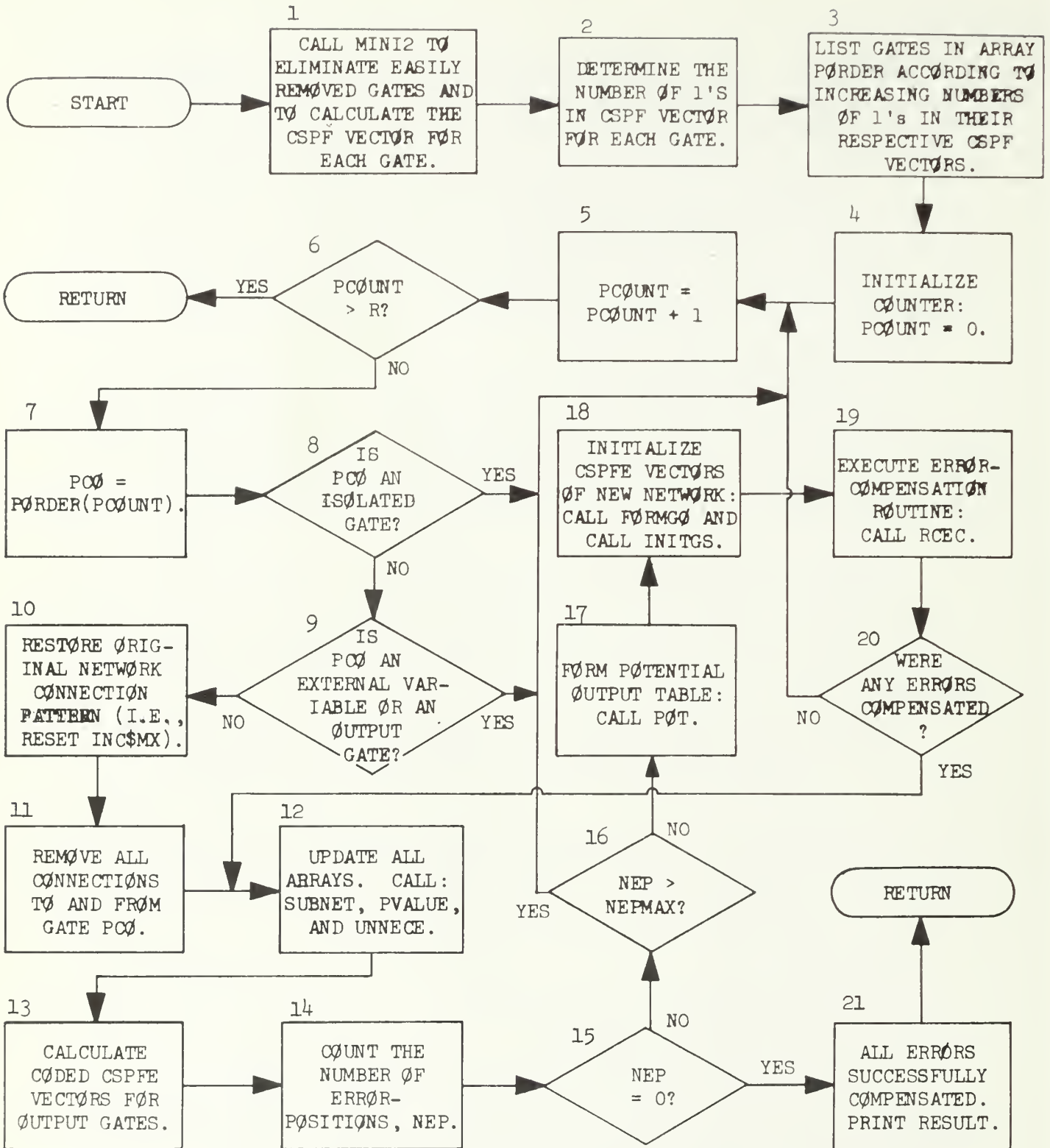


Figure 2.1.1 Generalized flowchart of PROCCE.

1's in a gate's CSPF vector reflects the relative "importance" of that gate in the network. In general the removal of a gate with many 1's in its CSPF vector from a given network is likely to produce more error-positions in the output gates than would the removal of a gate with fewer 1's. Thus, the ordering in the array PORDER roughly lists gates which are, from PORDER(1) to PORDER(R), increasingly more difficult to compensate for (the errors caused) when they are removed from the original network. Trying to remove the most easily compensated gates (i.e., those whose removals cause fewer error-positions) first (as is done in blocks 5 through 20) will generally lead to a quicker solution.

In block 4 a counter, PCOUNT, is initialized to the value 0. PCOUNT will be incremented by 1 during every pass through block 5. If PCOUNT exceeds the value R (i.e., the number of gates in the network), a condition tested for in block 6, it means PROCCE was unable to successfully compensate for the removal of any gate from the network, and PROCCE returns to the calling subroutine (MAIN).

If PCOUNT was found to be less than or equal to R, block 7 will set the variable PCO equal to PORDER(PCOUNT). PCO is then the label of the next gate to be removed from the network by the error-compensation procedure.

Block 8 checks to see if PCO is an isolated gate. If it is, there is, of course, no purpose in removing it from the network, and so the program would move on to the next gate to be removed (by returning to block 5).

If PCO is not isolated, block 9 further tests to see if PCO is an external variable (PROCCE assumes that all n variables are essential to the network) or an output gate. In either case, it cannot be removed from the network, and so the next gate is selected for removal instead (by returning



to block 5).

If PCO passes these tests, PROCCE enters block 10. Here an array INC\$MX is initialized to contain the connection pattern of the original network (i.e., the network connection pattern as it existed immediately after the transformation by MINI2 in block 1). The value  $\text{INC\$MX}(v_i, v_j)$  indicates the presence or absence of a connection from gate or external variable  $v_i$  to gate  $v_j$ .

This initialization is done in preparation for the removal of PCO from that original network. During the calculation, the array INC\$MX always contains the most recently updated version of the network connection pattern.

Block 11 removes PCO from the original network by removing all of its input and output connections. This is done by changing some of the values of INC\$MX.

The removal of PCO also causes changes in many other arrays storing various information related to the network configuration. These are updated by calling the subroutine SUBNET and its entry point PVALUE in block 12. Another entry point of SUBNET, UNNECE, is called to eliminate from the network any gates which may have been left, after the removal of PCO, with no paths to any of the network output gates. It is possible that leaving such gates in the network might actually be more beneficial, but an argument could be made either way. Programming considerations tipped the scale in favor of their removal.

Although there are only 5 logical types of components permitted in the CSPFE vectors, \*, 0, 1, 0, and 1, the variety of the actual codings employed to realize these 5 types is somewhat larger (this will be discussed further in Section 2.2.1). In block 13 the outputs of all of the gates

remaining in the new network are recalculated. The actual (new) network outputs are compared with the desired outputs, and from this information the (initial) coded entries of the CSPFE vectors of the network output gates can be directly determined.

The number of error-positions in this new network, NEP, are counted in block 14. If the number of error-positions is 0, the new network realizes all of the desired functions correctly, and block 15 directs control of the program to block 21.

Otherwise, the program enters block 16 where NEP is checked to see if it exceeds the maximum number of allowable error-positions, NEPMAX. NEPMAX is a parameter which may be varied by the user (see Section 5). If NEPMAX is exceeded, PROCCE abandons all hope of compensating for errors in so many positions, and control goes back to block 5 for the selection of a new PCO.

If NEPMAX was not exceeded, block 17 forms the "potential output table" by calling the subroutine POT (POT will be discussed in more detail in Section 2.2.2). Essentially, the potential output table lists all (theoretically all - but actually just a "great many") functions which either exist or can be "manufactured" (by adding the appropriate connections to the network) by a certain algorithm. This table is used to assist in error-compensation by providing a list of functions which may be connected as new inputs to certain gates during the procedure.

This is followed by another preparatory step just before the execution of the main error-compensation subroutine, RCEC. The CSPFE vectors are initialized in block 18 by calling INITGS (an entry point of MINI2). (The call to FORMGO merely recalculates GORDER in preparation for the execution of

INITGS.) This initialization sets the values of some components of certain CSPFE vectors which can be predetermined. This initialization generally improves both the efficiency and effectiveness of the procedure.

Block 19 is the most important part of the flowchart. Here, the heart of the procedure, subroutine RCEC, is called. The subroutine RCEC (discussed in detail in Section 2.3) performs a function somewhat similar to that of the subroutine PROCII in [1], except that it must also deal with the added difficulty of "errors" (i.e., undesired outputs of certain gates for certain input combinations) present in the network. The main purpose of RCEC is to rearrange the network connection pattern in a manner such that all of these "errors" can be compensated. This cannot be done in just a single call to RCEC however. When RCEC corrects just a single error of a certain type, it must return to PROCCE (block 12) to have all of the necessary arrays updated before it can continue its error-compensation task. If the correction of that error causes the network to function correctly, this fact will be detected in block 15.

If no errors of any kind were compensated by RCEC, it means that the compensation of further errors is impossible (at least, by following the algorithm realized by RCEC), so control of the program is sent back to block 5 to select a new PCO. The test is made in block 20.

When block 15 detects the presence of an error-free network (i.e., all of the desired output functions are correctly realized by the network) with the gate PCO removed, control goes to block 21 which prints out the new result. PROCCE then returns to the calling subroutine.

## 2.2 Major Support Subroutines

This section will discuss some of the details of two important support subroutines called by PROCCE, MINI2 and POT. A knowledge of the functions of these two subroutines is necessary in order to understand the error-compensation subroutine, RCEC (discussed in Section 2.3).

### 2.2.1 Subroutine MINI2

The subroutine MINI2 was discussed in an earlier report, [7], with an emphasis on those points of the subroutine most relevant to the operation of the program NETTRA-PG1. In block 1 of the flowchart of PROCCE (see Figure 2.1.1), MINI2 is used in a capacity similar to that in NETTRA-PG1. Since this function has already been discussed adequately in [7], it will not be repeated here.

Of greater interest with regard to the error-compensation procedure is the function of MINI2 in block 18 of the flowchart of PROCCE. In that block, an entry point, INITGS, of MINI2 is called which performs an initialization of the CSPFE vectors of the network. Although this aspect of the subroutine MINI2 was mentioned in [7], it was not discussed in detail. The rest of this section will explain the results of calling INITGS.

During the main body of the error-compensation procedure (i.e., during the call to RCEC in block 19 of the flowchart of PROCCE) the compatible sets of permissible functions are determined by assigning logical \*'s, 0's, 1's, 0's, and 1's to elements of GSMALL (recall that the array GSMALL is used to store the CSPFE vectors). There is usually considerable freedom in making these assignments, and consequently a large variety of collections of compatible sets for the entire network can be produced. Despite this freedom though,

there are certain assignments of logical 1's and 0's to certain entries of GSMALL which are predetermined by the configuration of the network and by the algorithm to be applied. These necessary assignments are made during the initialization step in order to avoid making certain unnecessary assignments (to particular elements of GSMALL) later, as would otherwise occur.

It is important, however, to remark that the existence of such predetermined entries is only possible under the assumption that only a single call to RCEC will follow the initialization. Partly for this reason, INITGS is always called to re-initialize the CSPFE vectors immediately before each call to RCEC (see blocks 18 and 19 of the flowchart of PROCCE). Two or more successive calls to RCEC (without changing the initialization assignments) might change the network configuration so that any such initialization would be invalid. In any event, a call to RCEC usually alters the network configuration so that all of the CSPFE vector entries calculated during that call by RCEC itself must be recalculated (in the next call to RCEC), so some sort of initialization must be performed before each call anyway.

As previously mentioned, five different logical entries, \*, 0, 1, 0, and 1, may appear as components of the different CSPFE vectors. The precise meanings of these five logical entries are discussed in [5], and they will also become clear during the discussion of the error-compensation subroutine (Section 2.3). In general terms, however, the meanings are as follows:

- \* - indicative of a "don't-care" component. In other words, every permissible function of the corresponding CSPFE may

have either a 0 or a 1 for this component.

- 0 - indicative of a required 0 component. Every permissible function of the corresponding CSPFE must have a 0 for this component position.
- 1 - indicative of a required 1 component. Every permissible function of the corresponding CSPFE must have a 1 for this component position.
- 0 - indicative of currently required 0 that should preferably be a required 1 instead (this is called a "0-error"). Although every permissible function of the corresponding CSPFE currently has a 0 for this component, it would be desirable to change the network configuration so that a 1 could be demanded instead.
- 1 - indicative of currently required 1 that should preferably be a required 0 instead (this is called a "1-error"). Although every permissible function of the corresponding CSPFE currently has a 1 for this component, it would be desirable to change the network configuration so that a 0 would be demanded instead.

To understand the coding of these logical values into the actual values used by the program, the concept of "covering" must be understood. Due to the nature of NOR gates, a 1 appearing on any of the input lines to a gate will cause a 0 output of that gate. Such a 1 is called a cover of that 0, and the 0 is said to be covered by that 1. Although a 0 output may be covered by several 1's (appearing on different input lines), only a single cover is actually required to guarantee the 0 output. For a given



gate, certain 0 components of its CSPFE vector (assuming that its CSPFE vector has already been calculated somehow - at least partially), representing required 0 outputs of the gate, each have only a single 1 cover. In each such case, the single 1 covering the 0 is called an essential 1. These essential 1's form the basis of the initialization steps performed by calling INITGS.

The correspondence between the logical values \*, 0, 1, 0, and 1 and the actual (coded) values used in the program is shown in Table 2.2.1.1 (the symbols  $\alpha$  and  $\beta$  used in the table represent variable integer values which may appear as the computer's internal representation of the logical value 0 of a CSPFE component).

<u>LOGICAL</u> <u>VALUE</u> <u>OF CSPFE</u> <u>COMPONENT</u>		<u>ACTUAL</u> <u>VALUE</u> <u>OF CSPFE</u> <u>COMPONENT</u>
---	--	--

*	→	0
0	→	$\alpha$ , where $0 > \alpha > -1000$

$\alpha = -100$  means this component has an unknown number of covers

$\alpha = -200$  means this component has 2 or more covers.

$\alpha = -\beta \neq -100, -200$  means this component has an essential 1 cover from gate or external variable  $\beta$

1	→	1
<u>0</u>	→	-1100
<u>1</u>	→	1001

Table 2.2.1.1 Coding of logical values for components of CSPFE vectors.

The sequence of events when INITGS is called is shown in Figure 2.2.1.1. INITGS assumes certain preliminaries have already been completed before it is called. For example, GORDER is assumed to be updated and

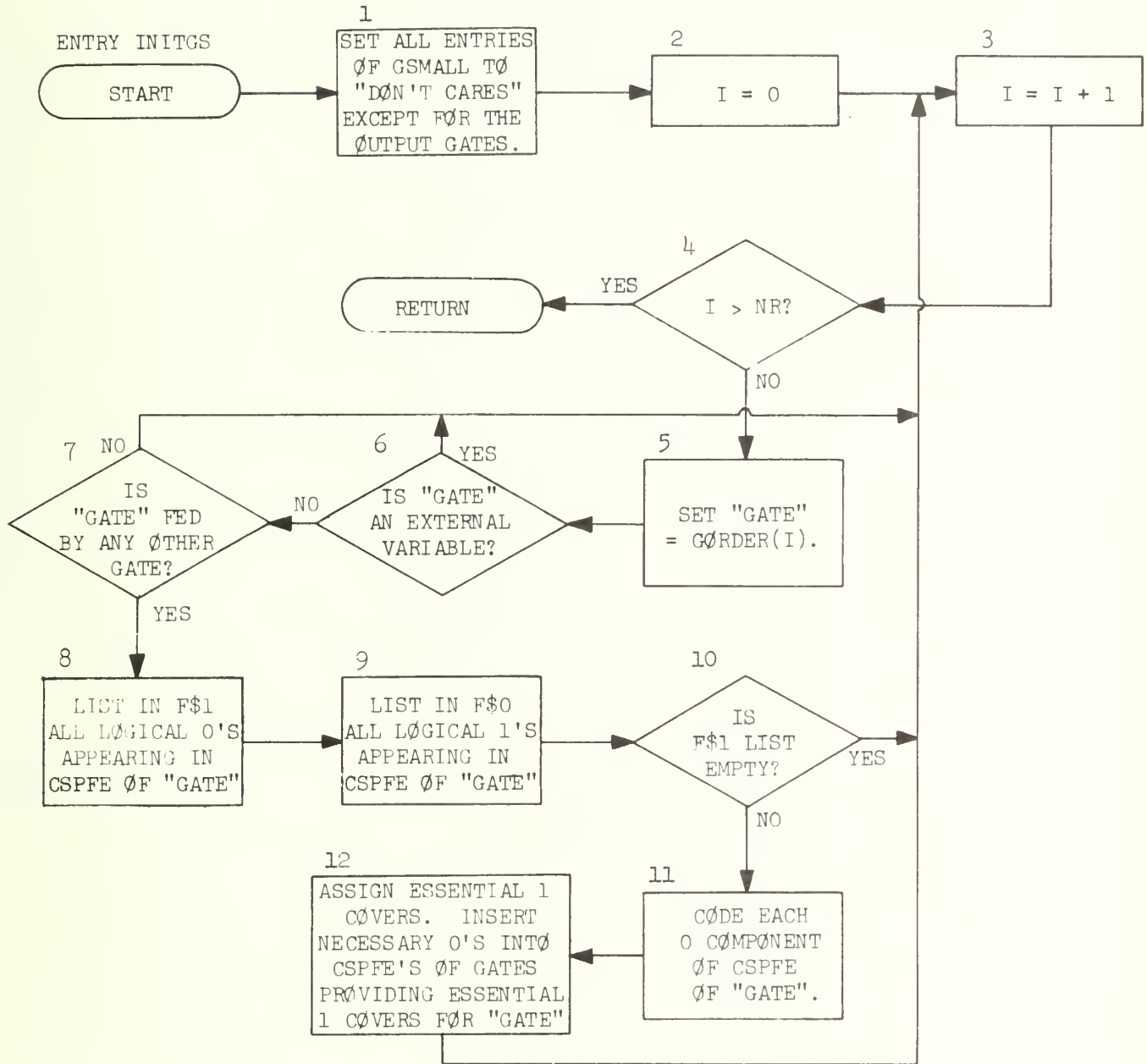


Figure 2.2.1.1 Flowchart of a section of the subroutine MINI2 (beginning at entry point INITGS).

GSMALL is assumed to already contain the correct values for the CSPFE vectors of the output gates (this is done by PROCCE).

In block 1 of the flowchart in Figure 2.2.1.1, all of the entries in GSMALL, except for those corresponding to the network output gates (since these have already been determined), are set to "don't cares" in preparation for the initialization.

Block 2 sets a counter, I, to the value 0. Thereafter, each pass through block 3 will increase I by 1. If I should exceed NR (the total number of gates and external variables), as tested in block 4, a return is made to PROCCE.

Otherwise, block 5 sets the variable GATE equal to GORDER(I). GATE is the current gate under consideration. Elements of the CSPFE vectors of the gates which feed GATE will be initialized according to the already initialized CSPFE vector of GATE itself (the ordering contained in the array GORDER guarantees that the CSPFE vector of GATE will be completely initialized before the CSPFE vectors of any of its predecessors).

Blocks 6 and 7 test whether or not GATE is an external variable or an isolated gate. If it is either, the program returns to block 3 to increase I and select a new GATE.

A list is made in block 8 of the component positions of all the logical 0's currently appearing in the CSPFE of GATE. These are stored in the 1-dimensional array F\$1.

Similarly, block 9 lists the positions of all of the logical 1's and stores them in the array F\$0.

If F\$1 contains an empty list, GATE has no logical 0's in its CSPFE to be covered by its predecessors, and the covering step in block 12 can be skipped. The test is made in block 10.

By checking the number of predecessors which cover each logical 0 in GATE's CSPFE vector, block 11 determines the actual coding for each of these logical 0's individually (according to Table 2.2.1.1).

This coding is used by block 12 which assigns values to (i.e., initializes) certain components of the CSPFE vectors of GATE's predecessors. Each logical 0 in GATE's CSPFE vector whose actual coding is of the  $-\beta$  ( $\beta \neq 100, 200$ ) type has an essential 1 cover. In other words, each logical 0 of this type is covered by the function from only a single gate.

Making the assumption that every logical 0 or 1 appearing in the CSPFE vector of GATE can be shown to be a required value assignment regardless of the finally chosen sets of CSPFE vectors (i.e., for any choice of compatible sets of permissible functions for the gates of the network, these components of the CSPFE vector of GATE will always have the same assigned value), one can assert that the essential 1 covers of those logical 0's are also required regardless of the finally chosen CSPFE vectors (assuming no change in the network configuration). Thus justified, an assignment of a logical 1 is made to the CSPFE vector of a predecessor, PRED, of GATE in a component position corresponding to an essential 1 cover provided by PRED for a 0 in GATE's CSPFE vector. This is done by block 12 for every 0 in the CSPFE vector of GATE.

Furthermore, since every predecessor, PRED, of GATE which provides an essential 1 cover for at least one 0 of GATE's CSPFE vector cannot be disconnected from GATE without introducing a new error into the network, it is clear that the CSPFE vector of PRED must also contain a logical 0 in every component position where a logical 1 appears in GATE's CSPFE. These assignments are also made in block 12.

Generally the initialization of GSMALL (i.e., the CSPFE vectors) can be seen as a propagation of required values of certain CSPFE vector components from the output gates through the gates feeding them, through the gates feeding these gates, etc.: The complete CSPFE vectors for the network output gates are known completely from the beginning. For the gates providing essential 1 covers for certain 0's in the CSPFE's of the output gates, one can make assignments of 0's and 1's to their own CSPFE vectors which are valid independent of whatever algorithm is used later (after the initialization step) to make a complete assignment to all components of all CSPFE vectors. This step is then repeated for the gates providing essential 1 covers for the gates providing essential 1 covers for the output gates, and so on.

### 2.2.2 Potential output table (POT)

In order to compensate for error-components at the CSPFE of a gate, functions currently realized at other gates and external variables may be used according to the basic transduction procedure using CSPFE's discussed in [5]. In order to make the error-compensation more flexible, the concepts of potential outputs and potential output tables (POT) were introduced in [5], and a procedure utilizing a potential output table was also given. This section will briefly explain these concepts and discuss in some detail the implementation of the potential output table in programs NETTRA-E1, -E2, and -E3. (For details of these concepts see [5]).

A potential output from a gate GI is a function realized at GI by connecting additional inputs to GI. This potential output of GI differs from the function currently realized at GI, and therefore can be used to

compensate for errors in a certain gate to which the current output function at GI is not connectable (i.e., connecting GI to that gate will change some non-error components in its CSPFE). There may be two major problems in using potential outputs. One is the problem of over-compensation caused by using too sophisticated potential outputs (i.e., too many modifications are required in producing potential outputs) that the number of error components in other gates may increase and it becomes too difficult to completely compensate for them. This problem can be avoided by restricting the types of modifications used in producing potential outputs. As explained in [5], only potential outputs which can be realized by adding connections satisfying triangular conditions are allowed in the error compensation procedure. Another potential problem is the additional computational complexity required by searching for potential outputs. This problem can be serious if a complete searching of potential outputs is required each time when a potential output is desired. This problem of time-consuming search cannot be completely avoided (if potential outputs are to be used) but can be reduced to a certain degree by employing a potential output table (POT). The POT is a list of selected potential outputs and the information on how to construct each potential output. For the sake of convenience, functions which are currently realized at input terminals (external variables) and gates are also listed. In the beginning of the error compensation process, all potential outputs satisfying certain triangular conditions are searched and listed in the form of a table. During the error-compensation process for a particular gate GI, the candidates for new inputs to GI (strongly effectively E-connectable functions with respect to the CSPFE of GI) are selected from POT. Therefore,



at the expense of memory space (for storing POT), the relatively time consuming process of searching for potential outputs can be replaced by a simple table look-up process.

The potential output table used in NETTRA-E1, -E2, and -E3 is organized as follows.

(1) The potential output table (POT) is stored in a two dimensional array POTAB(200,42). The first argument of POTAB is the entry number. For example, the information on the PTR-th entry of POT is stored in POTAB(PTR, 1) ~ POTAB(PTR, 42).

(2) Each entry, corresponding to one potential output, contains the following information:

- (a) POTAB(PTR, 1) ~ POTAB(PTR, 32): the actual function (vector) of the PTR-th potential output in a truth table form.
- (b) POTAB(PTR, \$GT): the gate label of gate GI at which the PTR-th potential output is realized (\$GT is an integer constant and has the value 33).
- (c) POTAB(PTR, \$LTH): the number of connections to be added to gate GI in order to realize the PTR-th potential output (\$LTH is an integer constant and has the value 34). POTAB(PTR, \$LTH) may have the value 0 through 6.
- (d) POTAB(PTR, \$LTH+1) ~ POTAB(PTR, \$LTH+6): the list of external variables and/or gates whose outputs must be connected to GI in order to realize the PTR-th potential output.
- (e) POTAB(PTR, \$PW) and POTAB(PTR, \$NOE): the preference weight and the number of one errors, in the PTR-th

potential output, respectively. Since these two values are defined with respect to the CSPFE of a particular gate, these two fields are used only during the error-compensation process when a gate has been selected.

(\$PW and \$NOE are integer constants and have the values 41 and 42, respectively).

(3) Entries of POT are divided into several blocks of consecutive entries. Each block contains the potential outputs realized at one particular gate. In other words, every entry in the same block has the same POTAB(PTR, \$GT), and for every gate or external variable in the network there is a corresponding block containing the potential outputs realized at that gate or external variable (a block may consist of only one entry).

(4) For each gate or external variable in the network, there are two pointers PPOTAB and LPOTAB both of which are one dimensional arrays. PPOTAB(GI) indicates the starting entry of the block corresponding to GI, whereas LPOTAB indicates the last entry of the same block.

(5) In the block corresponding to gate or external variable GI, the entries are listed in the following order:

- (a) The function currently realized at GI, and therefore  $POTAB(PPOTAB(GI), \$LTH) = 0$
- (b) The functions realized at GI with one additional connection (i.e.,  $POTAB(PTR, \$LTH) = 1$ ). For the sake of convenience, these entries are called simple entries.
- (c) Entries realized at GI with  $POTAB(PTR, \$LTH) \geq 2$ .

An entry with  $POTAB(PTR, \$LTH) \geq 2$  is called a composite entry because it can be obtained by an operation with two other entries with the

same  $POTAB(PTR, \$GT)$ . For example, if the  $PTRA$ -th and  $PTRB$ -th entries satisfy  $POTAB(PTRA, \$GT) = POTAB(PTRB, \$GT) = GI$ , a composite entry  $PTR$  can be generated as follows.

- (1)  $POTAB(PTR, I) = POTAB(PTRA, I) \wedge POTAB(PTRB, I)$ , for  $I = 1, 2, \dots, 32$ .
- (2)  $POTAB(PTR, \$GT) = GI$ .
- (3)  $POTAB(PTR, \$LTH) = POTAB(PTRA, \$LTH) + POTAB(PTRB, \$LTH)$ .
- (4)  $POTAB(PTR, \$LTH + I) = POTAB(PTRA, \$LTH + I)$ , for  $I = 1, \dots, POTAB(PTRA, \$LTH)$ .
- (5)  $POTAB(PTR, \$LTH + I) = POTAB(PTRB, \$LTH + I - POTAB(PTRA, \$LTH))$ , for  $I = POTAB(PTRA, \$LTH) + 1, \dots, POTAB(PTRA, \$LTH) + POTAB(PTRB, \$LTH)$ .

In programs NETTRA-E1, -E2, and -E3, the potential output table for a given network is generated by subroutine POT which realizes the following procedure.

### Procedure for Generating Potential Output Table

#### Step 1. Initialize

Set  $PPOTAB(GI) = 0$  for  $GI = 1, \dots, N + R$ , where  $N$  and  $R$  are the numbers of external variables and gates in the network, respectively. Set  $POINTR = 1$  ( $POINTR$  is a pointer indicating the next entry to be generated).

#### Step 2. Selection of gates

According to the level assigned to each gate and external variable (i.e., a gate in a higher level precedes a gate in a lower level), select a gate or an external variable,  $GI$ , in the network. If all gates and external variables have been considered, the procedure terminates.

Step 3. Set the first entry for GI

Set  $POTAB(GI) = POINTR$ .

Copy the present function realized at GI into  $POTAB(POINTR, 1) \sim POTAB(POINTR, N2)$ , where  $N2$  is the number of input combinations (if completely specified,  $N2 = 2**N$ ).

Set the following values:

$POTAB(POINTR, \$GT) = GI$ ,

$POTAB(POINTR, \$LTH) = 0$ , and

$POINTR = POINTR + 1$ .

If  $GI \leq N$ , go to Step 6.

Step 4. Find the simple entries realized at GI

Step 4-1 Select a gate or an external variable, GJ, whose level number is not lower than that of GI. If all such GJ's have been considered, go to Step 5.

Step 4-2 Check whether or not GJ is connected to all immediate successors of GI. If not, go to Step 4-1.

Step 4-3 Check whether or not connecting GJ to GI produces a function which differs from the function of GI and from the functions produced by connecting previous GJ's to GI. If not, go to Step 4-1.

Step 4-4 Make the  $POINTR$ -th entry as follows.

$POTAB(POINTR, \$GT) = GI$

$POTAB(POINTR, \$LTH) = 1$

$POTAB(POINTR, \$LTH + 1) = GJ$

$POINTR = POINTR + 1$

Go to Step 4-1.

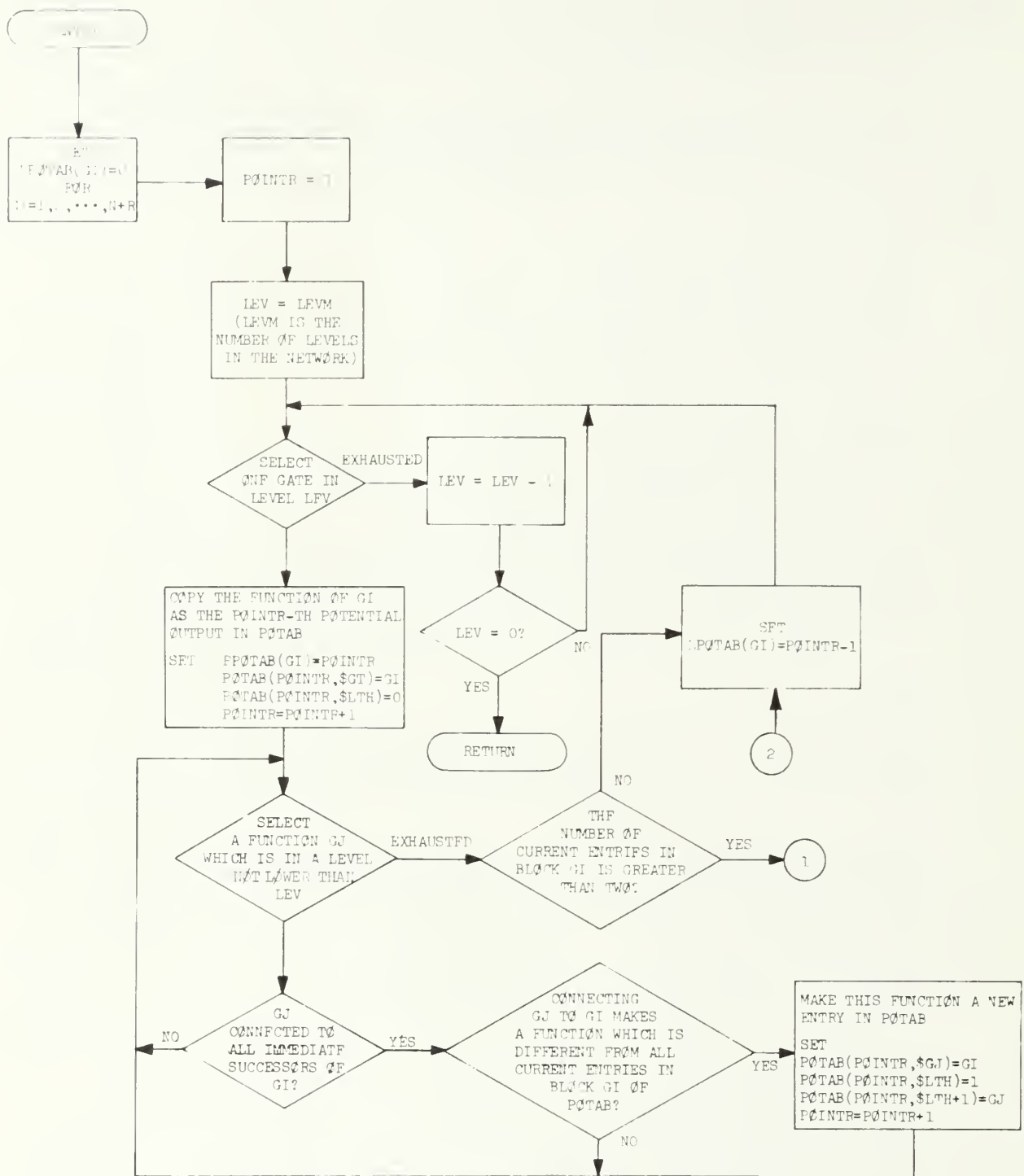


Figure 2.2.2.1 Flowchart of subroutine POT

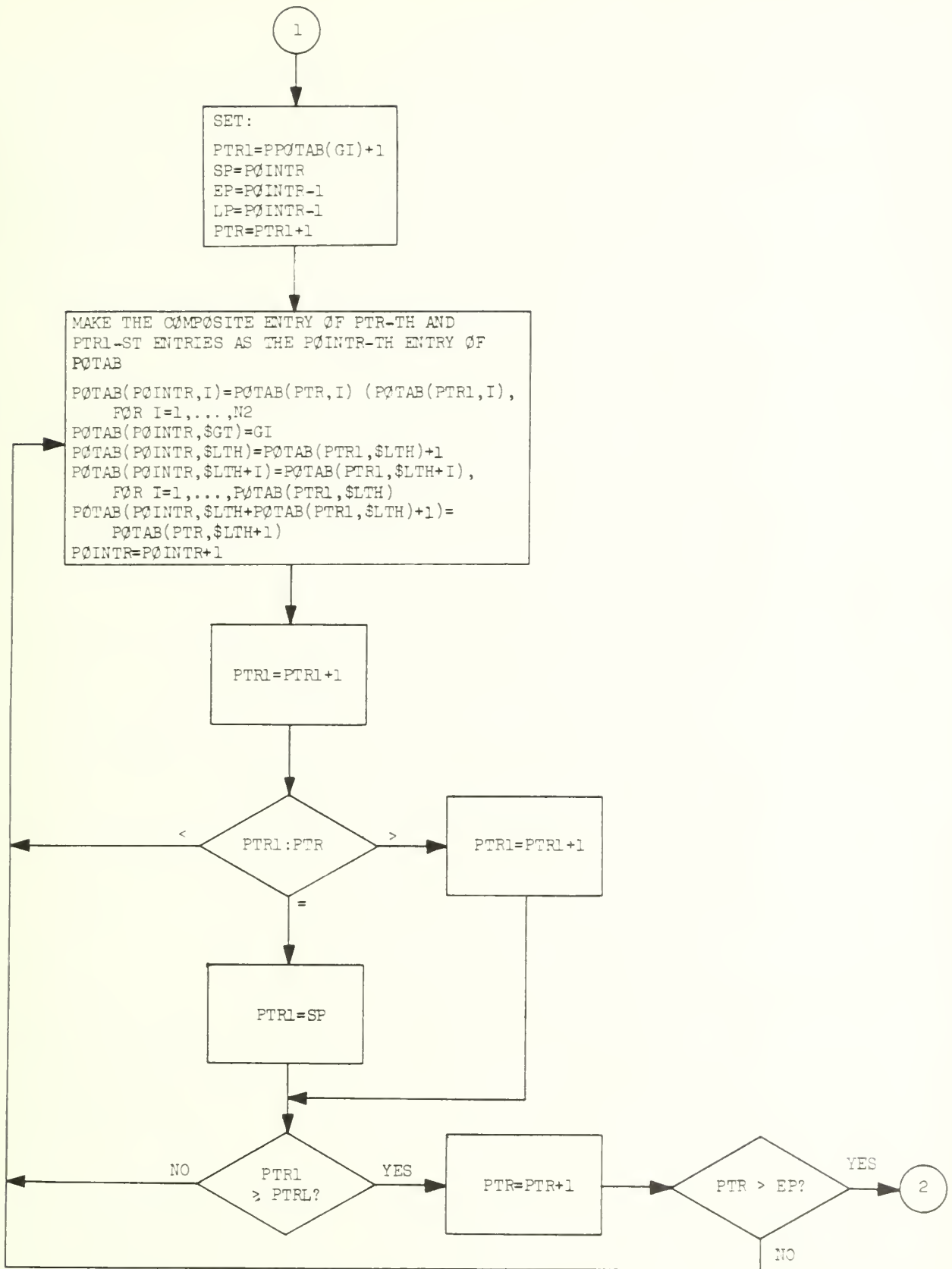


Figure 2.2.2.1 Continued.

Step 5. Generate composite entries to be realized at GI

If there is no more than one simple entry for GI, go to Step 6; otherwise generate composite entries as follows. (S denotes the total number of simple entries, which are the entries  $PPOTAB(GI) + 1$  through  $PPOTAB(GI) + S$  in the potential output table, and POINTR always points at the next entry to be entered).

Step 5-1 Set  $I = 2$  and  $E = S$ .

Step 5-2 Generate a composite entry from the  $(PPOTAB(GI) + I)$ -th and the  $(PPOTAB(GI) + J)$ -th entry, for  $J = 1, \dots, I - 1$ .

Step 5-3 Generate a composite entry from the  $(PPOTAB(GI) + I)$ -th entry and the  $(PPOTAB(GI) + J)$ -th entry, for  $J = S + 1, \dots, E$  (if  $E \leq S$ , this step is skipped).

Step 5-4 Set  $I = I + 1$ . If  $I > S$ , go to Step 6; otherwise set  $E = POINTR$  and go to Step 5-2.

Step 6. Set the pointer indicating the last entry in block GI

Set  $LPOTAB(GI) = POINTR - 1$ , and go to Step 2.

The flowchart of this procedure is shown in Figure 2.2.2.1.

### 2.3 Error-Compensation Subroutines

Error compensation procedure is coded into a FORTRAN subroutine named RCEC which stands for Replacement of Connections for Error-Compensation and five supporting subroutines: CALS1, RPLCF, FORC, ORDRQ2, and CONECT. RCEC is the central part of the transduction programs NETTRA-E1, -E2, and -E3. This subroutine needs, as parameters, a NOR network which does not realize the desired output functions along with the initialized CSPFE's for each gate



(explained in Section 2.2.1). When it is entered, it selects one gate at a time whose CSPFE has been completely calculated and tries to compensate for error-components in the CSPFE of the selected gate. As a result of the error-compensation for this particular gate, if some error-components are compensated or output functions of some gates are changed, this subroutine will return to PROCCE, the calling subroutine in NETTRA-E1, -E2, and -E3. PROCCE will then recalculate the output functions of the network to check whether or not this new network realizes the desired functions. If it does, the original network has already been transduced to a desired new network; otherwise PROCCE will recalculate the potential output table and call subroutine RCEC again to apply the error-compensation procedure to this slightly modified network. On the other hand, if no error-components in the CSPFE of the selected gate can be compensated, the CSPFE of the selected gate will be propagated to its inputs. If all gates have been considered in this manner, it means no error-components can be compensated in this application of RCEC (otherwise it would have returned to PROCCE). In this case, the subroutine returns to PROCCE unsuccessfully.

The error-compensation procedure consists of several subprocedures. These subprocedures and two supporting subroutines will be discussed in some detail in Sections 2.3.1 and 2.3.2, respectively. Section 2.3.3 will be devoted to the discussion of the propagation of CSPFE's, and Section 2.3.4 will summarize the entire procedure and present the flowchart.

### 2.3.1 Compensation of error-components for a particular gate

The error-compensation procedure considers only one gate at a time. When a gate is selected, its CSPFE must be completely calculated. Therefore,

at the beginning of the procedure, only first level output gates can be selected. As the calculation goes on, a gate becomes selectable only when all its immediate successors have been selected. Thus the ordering of selection can be made according to the gate level assigned to each gate. The ordering also could take into consideration the number, type and degree of the error-components in the CSPFE of each gate when more than one gate is selectable as explained in [5].

After a gate GI has been selected, the procedure concentrates on compensating for error-components in the CSPFE of gate GI by (1) removing redundant input connections, (2) substituting for input connections, and (3) adding connections to compensate for 1-error-components. The first two types of operations are aimed at compensating for 0-error-components whereas the third one 1-error-components. In all cases, the number of error-components in the CSPFE of gate GI will never increase after applying these operations. In addition to the number of error-components, the degree of an error-component is also an important criterion in deciding which connections are to be added or disconnected. The degree of an error-component is defined only for 0-error-components to indicate how difficult this 0-error-component is to be compensated. If a 0-error-component in the CSPFE of gate GI is covered by only one input connection of the gate, this error is considered easier to be compensated since the proper substitution for this input connection or the compensation for the corresponding 1-error-component in that input can compensate for this 0-error-component. This type of 0-error-components are called primary 0-error-components<sup>†</sup>. On the other

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<sup>†</sup> The Definition 4.3.2.3 of [5] defines primary 1-error-components which are the ones which cover primary 0-error-components defined here. For the convenience of discussion both are regarded as primary error-components in this paper.

hand, if a 0-error-component is covered by more than one input , it is considered more difficult to be compensated either at this stage when the gate having this 0-error-component is under consideration or later when the immediate predecessors covering this 0-error-component are under consideration. This type of 0-error-components are called secondary error components.

Among the three types of operations mentioned earlier, the last two require the connection of functions to the selected gate GI. The set of candidates for those functions must be strongly effectively E-connectable functions with respect to the CSPFE of gate GI. Among these candidates the functions which have been actually selected to be connected to GI as inputs must also satisfy the prohibition conditions (Lemma 3.2 of [5]) with each another. These conditions are examined each time when a function has been selected, and the functions which do not satisfy one of these conditions with this selected function will be prohibited from being selected.

The above conditions are required for the functions to be connected to gate GI both in the second (substitution of connections) or the third type (addition of connections) operation. In addition to these conditions, the replacing functions and the replaced functions must satisfy the following condition. That is, the addition of the entire set of replacing functions will make the set of to-be-replaced functions E-disconnectable with respect to the CSPFE of gate GI. This condition guarantees that the number of 0-error-components never increases. In the actual procedure, the operation of the substitution for input connections consists of three subprocedures. Each subprocedure is aimed at compensating for a particular type of

error-components. Additional conditions are required in each subprocedure in order to actually compensate for some error-components, and will be discussed later in this section. For the third type of operation, each of the added connections must have at least one 1-component which covers a 1-error-component in the CSPFE of gate GI since this operation is aimed at compensating for 1-error components.

The input connection substitution problem is essentially a covering problem. All 0-components in CSPFE of gate GI must still be covered after the substitution, and the number of 0-error-components covered by the new input set (the replacing functions and the remaining functions) is to be minimized. Although the optimal solutions can be obtained by solving this covering problem, the time required for deriving optimal solutions would usually be too excessive, especially when the number of candidates for substituting functions is very large. Furthermore, the local optimization, i.e., the optimization concerning only the selected gate GI is not necessarily the optimization with respect to the entire error compensation procedure. Based on these considerations, a heuristic method for substitution subprocedures is used.

The error compensation procedure contains the following six subprocedures, among which, subprocedures (2), (3), and (4), are substitution subprocedures.

(1) Remove redundant connections.

The redundant connections are the input connections which are E-disconnectable with respect to the CSPFE of the selected gate GI. Since it has no essential 1-components to cover 0-components in the CSPFE of GI, the removal of this connection will not increase the number of error-components.

The number of 0-error-components or the degrees of some 0-error-components in the CSPFE of the selected gate may decrease resulting from the removal of redundant connections. It should be noted that the redundant connections are removed before substitution is considered but the removed input connection may be reconnected in the substitution subprocedures.

- (2) Substitution for input connections from external variable with error-components.

If an input connection from an external variable has a 1-component which covers a 0-error-component in the CSPFE of gate GI, this 0-error-component can never be compensated unless this input connection to gate GI is removed. Therefore, this type of input connections must be considered for substitution prior to other input connections. Since the main purpose of this subprocedure is to replace the input connections having uncompensatable errors, the strongly effectively E-connectable functions (for simplicity, strongly effectively E-connectable functions will henceforth be referred to as connectable functions) without uncompensatable error-components, i.e., the connectable function from external variables without error-components and connectable function from gates with or without error-components may be used as the candidates for the substituting functions.

- (3) Substitution of input connections from gates with primary errors.

If a 0-error-component in the CSPFE of the selected gate is covered by only one input function, it is called a primary error as it may be corrected by substituting for that input function only. In this case, the connectable functions which have a 0-component corresponding to the primary 0-error-component under compensation are candidates for the substitution. Once a primary error is eliminated by substitution, the



candidates for later substitutions should be limited to those functions which have 0-components corresponding to the compensated primary 0-error-component.

(4) Substitution for input connections by functions without error-components.

This subprocedure substitutes connectable functions which have no 1-components corresponding to 0-error-components in the CSPFE of the selected gate GI for functions which have at least one 1-component corresponding to a 0-error-component in the CSPFE of the gate GI. As a result of the substitution the degrees of some 0-error-components may be reduced and this may make the total number of error-components in the input functions smaller. If the degree of an error is reduced to zero, this error of the selected gate has been compensated.

(5) Adding connections to compensate for 1-error-components.

A 1-error-component in the CSPFE of the selected gate GI is easy to be compensated if there is a connectable function whose corresponding component is a 1. Some of the 1-error-components may have been already compensated during the compensation for 0-error-components by substitutions. Since earlier subprocedures are aimed at compensating for 0-error-components, however, there are still possibilities to compensate for 1-error-components especially when there is no 0-error-component in the CSPFE of gate GI. In this subprocedure, the candidates are those functions which have a 0-component corresponding to each 1-component (including the corrected 0-error-components if any) in the CSPFE of gate GI.

(6) Adding redundant input connections from external variables.

This subprocedure does not belong to any of the three types of operations for error-compensation since it is not aimed at compensating for



error-components in the CSPFE of the selected gate GI, but rather, at loosening the requirements of the predecessors of gate GI to make error-compensation at later steps easier.

In subprocedure (2), the inputs from external variables with error-components have been considered for replacement in the first place since an error-component in an external variable can never be compensated unless it is removed. A redundant input from an external variable without error-components, however, will help to loosen the requirements for the predecessors of the selected gate GI if this input from external variable has some 1-components corresponding to the 0-components in the CSPFE of gate GI. Therefore, this type of redundant external variables (i.e., those without anticipated error-components) should be connected to gate GI. In addition, the external variables with some anticipated 1-error-components which cover the corresponding 0-error-components are also added if these 0-error-components are not primary errors (i.e., the addition of these redundant inputs from external variables will not decrease the number of primary 0-error-components in the CSPFE of the selected gate). This seems to contradict the objective of subprocedure (2), i.e., removing external variables with error-components by ~~any~~ means, but it should be noted that if some error-components can be compensated later, this gate will be considered again when subroutine RCEC is reentered. At that time, the redundant external variables, if still redundant, will be removed by subprocedure (1), and may not be added again if it would cover primary 0-error-components.

These subprocedures will be explained in some detail in Section 2.3.3.

### 2.3.2 Supporting subroutines for substitution

In the error-compensation procedure, there are three subprocedures which substitute a subset of candidates for a subset of input functions of a selected gate GI. Since these subprocedures are similar in nature, they are coded to share two subroutines, CALS1 and RPLCF. Subroutine CALS1 calculates the to-be-replaced subset of certain input functions when a set of candidates are given. Let the set of input functions to be removed be denoted with S, and the set of candidates for the functions to be added with  $S_2^\dagger$ . The subroutine CALS1 will calculate a subset S1 of S which can be replaced by S2. For example, in subprocedure (2) the set S contains all external variable inputs which cover some 0-error-components, and the set S2 contains all connectable functions except the connectable external variables which have 1-components corresponding to some 0-error-components in the CSPFE of the selected gate GI. The subroutine CALS1 checks whether or not every essential 1-component in each function of S can be covered by S2. If all essential ones in a function of S can be covered by set S2, this function is replaceable, and therefore should be placed in set S1. Along with set S1, the components which must be covered by one of the substituting functions are stored in set T1, which will be referred to by subroutine RPLCF.

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<sup>†</sup> In reference [5], symbols S,  $S_1$ ,  $S_2$ , and  $S_3$  were not used in order to avoid the possible confusion with  $S(v_i)$ , the set of successors for input terminal or gate  $v_i$ . Instead of S,  $S_1$ ,  $S_2$ , and  $S_3$ , symbols Q,  $Q_1$ ,  $Q_2$ ,  $Q_3$  were used to indicate a set of input functions containing the candidates to be replaced, a subset of Q which can be actually replaced, a set of candidates replacing  $Q_1$ , and a subset of  $Q_2$  which can actually replace  $Q_1$ , respectively. For better correspondence with the notation in the actual programs NETTRA-E1, NETTRA-E2, and NETTRA-E3, the symbols S, S1, S2 and S3 are used in this paper in place of Q,  $Q_1$ ,  $Q_2$ , and  $Q_3$ , respectively.

The procedure of subroutine CALS1 is as follows.

Step 1. Selection of functions

Take a function GP from set S. If all functions in S have been considered, return to the calling procedure (subroutine RCEC in NETTRA-E1, -E2, or -E3).

Step 2. Check replaceability

For each 0-component in the CSPFE of GI, the gate under consideration, which is covered only by input function GP, check whether or not set S2 covers it. If not, GP is not replaceable and go to Step 1.

Step 3. List essential ones

Add the positions of all essential 1-components in the function of GP into set T1, which is the list of the positions to be covered by the substituting functions.

Step 4. Update

Add GP into set S1. Remove GP from the list of current input functions of GI. Remove GP from set S. Go to Step 1.

The flowchart of subroutine CALS1 is shown in Figure 2.3.2.1.

Subroutine RPLCF is called immediately after CALS1 has returned. RPLCF selects a subset S3 of S2 which is needed to replace functions in set S1. Therefore, sets S3 and S1 must satisfy the conditions for substitutions described in Section 2.3.1. Since the selection of S3 is essentially a covering problem, as mentioned previously, a heuristic procedure is used. The candidates for the replacement of S1 (i.e., set S2) are stored according to the ordering based primarily on the number of

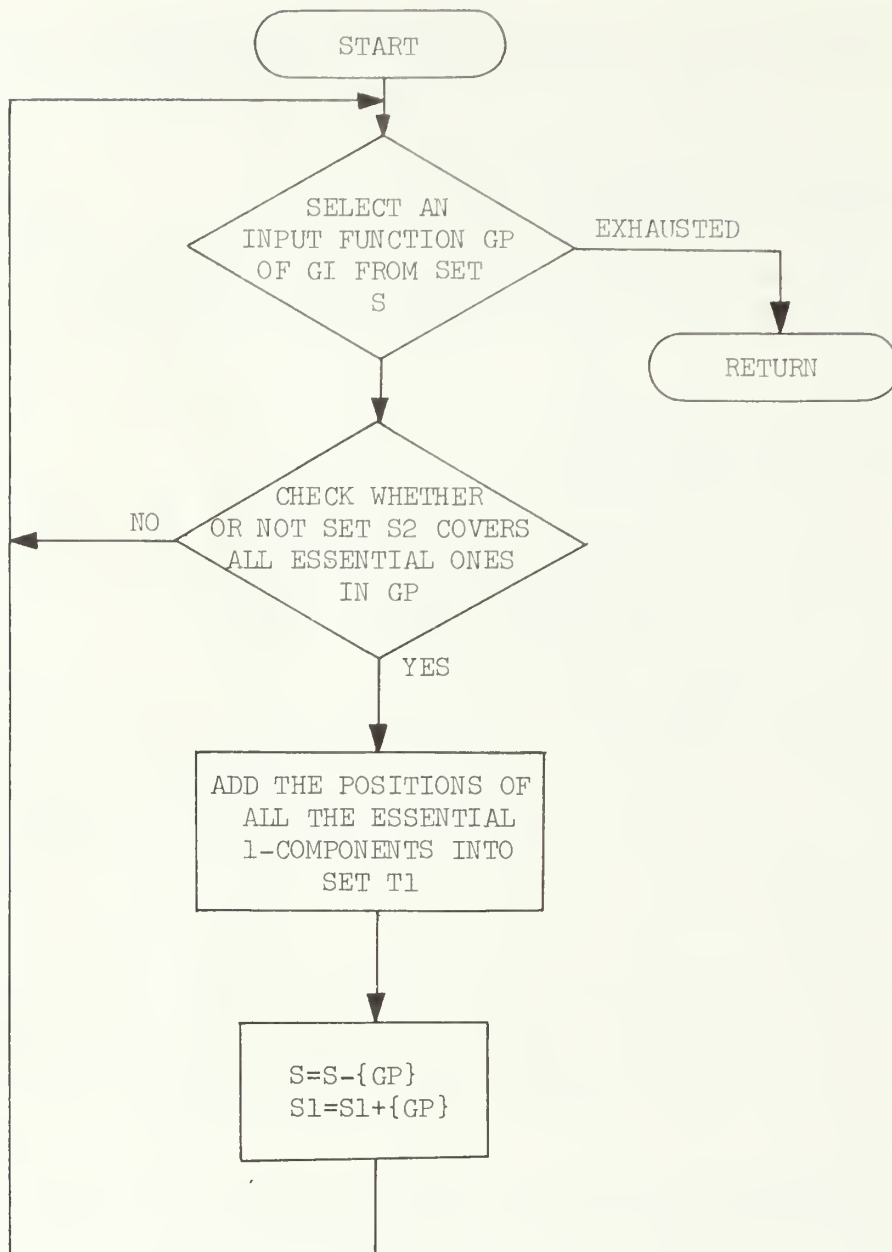


Figure 2.3.2.1 Flowchart of subroutine CALS1.

anticipated error-components (if connected) and secondarily the number of 1-components covering 0-components or 1-error-components. The procedure examines each candidate according to this ordering to check if it covers any components in T1 which are not covered by any selected candidates. If it does, the candidate will be added into set S3 which is the set of functions substituting for inputs in set S1. As mentioned in Sections 2.2.2 and 2.3.1, the candidates in set S2 may not be compatible with one another, so each time when a candidate is added into set S3, all other candidates which are not compatible with it are prohibited from being selected. Because of the prohibition, the remaining functions in set S2 and functions in set S3 may no longer cover set T1. In this case, the functions prohibited by the latest selection are permanently prohibited, and the set S2 is reconstructed. The procedure then calls subroutine CALS1 (from subroutine RPLCF) and reenter RPLCF itself to repeat this procedure of substitution. As a final result, a subset of input functions may be replaced by a subset of candidates in set S2, and the conditions listed in Section 2.3.1 are satisfied.

The procedure of subroutine RPLCF is as follows.

Step 1. Initialize

Set  $S3 = \emptyset$ .

Step 2. Selection of a candidate

Select a function, PTR, from set S2 which is the PTR-th entry of POT and is realized at GT.

Step 3. Check usefulness

Check whether or not this function covers any components in set T1 which are not covered by any functions already in set S3. If not, go to

Step 2.

Step 4. Prohibition

Prohibit functions which are realized either at gate GT or at some other gate GX requiring the addition of the connection from GT to GX.

Step 5. Update

Temporarily add this function into set S3. Remove components covered by this function from set T1. If  $T1 = \emptyset$ , go to Step 8.

Step 6. Check replaceability

Check whether or not remaining functions in set S2 still cover all remaining components in set T1. If yes, set  $S3 = S3 \cup \{PTR\}$  and go to Step 2.

Step 7. Recalculate set S1

Restore set S2 but delete from S2 the functions just prohibited in Step 4. Call CALS1 to recalculate S1 and associated T1. Go to Step 1.

Step 8. Substitution

Connect functions in S3 to gate GI. Disconnect the input functions in set S1 from gate GI. Return to the calling procedure (subroutine RCEC in NETTRA-E1, -E2, and -E3).

The flowchart of this subroutine is shown in Figure 2.3.2.2.

Beside the two subroutines mentioned above, there are three other supporting subroutines: FØRC, ØRDRQ2, and CONECT.

FORC has an argument GJ. When it is called, the connection from GJ to GI which is the gate under consideration is examined. If the connection from GJ to GI is E-disconnectable with respect to the CSPFE of GI, it will be removed from the network; otherwise it will do nothing.



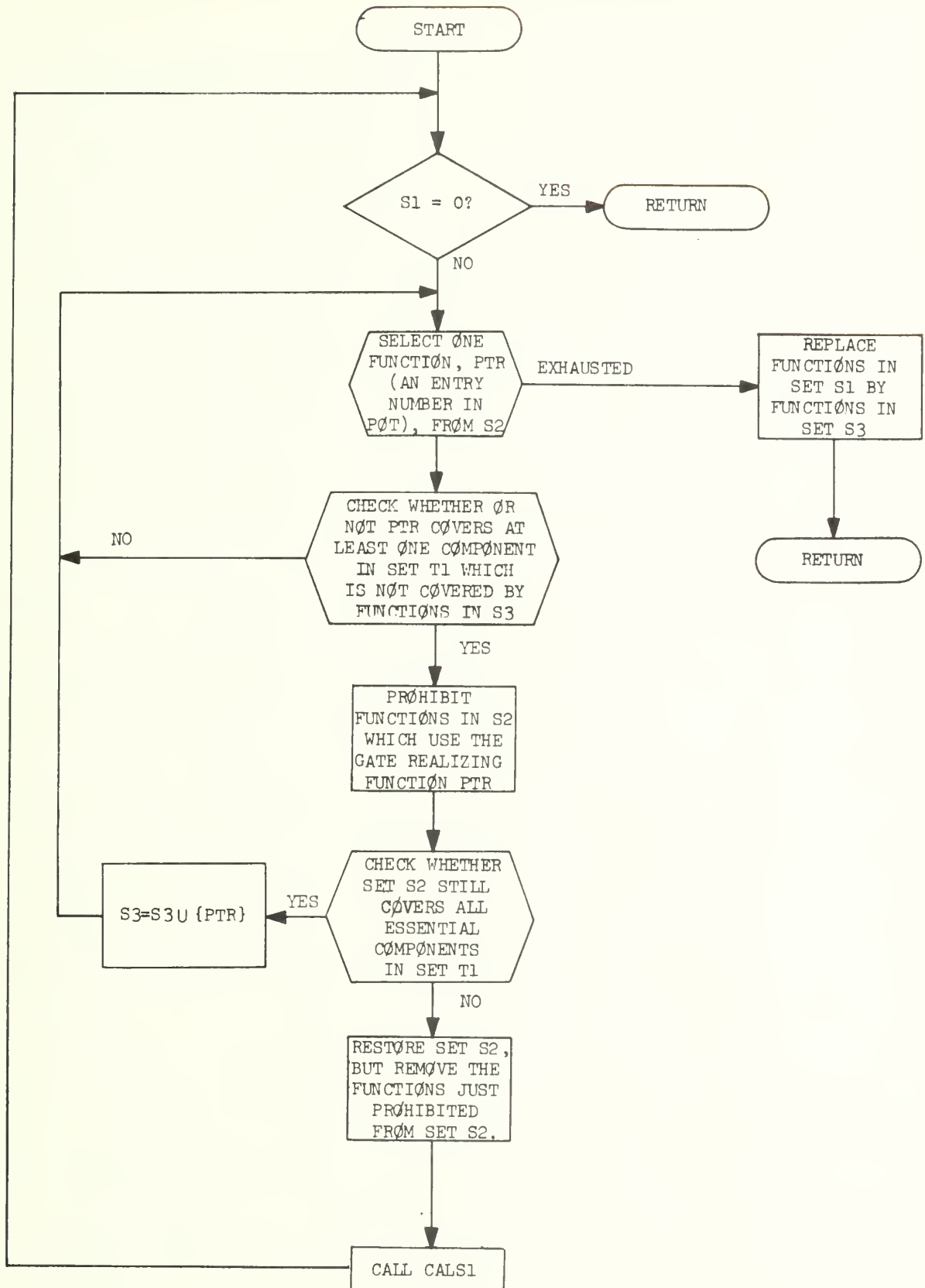


Figure 2.3.2 Flowchart of subroutine RPLCF.

Subroutine ORDRQ2 classifies the input functions except external variables of the selected gate GI into three groups. The first group contains input functions which have no 1-components corresponding to 0-error-components in the CSPFE of GI. The second group contains the functions which have at least one 1-component corresponding to a primary 0-error-component in the CSPFE of GI. The functions which have 1-components corresponding to some 0-error components but none of them are primary error-components are classified as the third group. Since the functions in the first group are not subjects of substitution, they are not listed. The other two groups are put into a 2-dimensional array working as two 1-dimensional arrays, each of which contains the sorted list of the functions in second and third groups, respectively. The sorting is based on the number of 1-components corresponding to the 0-error-components in the CSPFE of GI. These two lists are referenced by subroutine RCEC as the lists of candidates to be replaced.

Subroutine CONECT(PTR) realizes a simple procedure which connects the PTR-th function in the POT to gate GI, the gate under consideration. As the realization of the PTR-th function in the POT may require additional connections as explained in Section 2.2.2, subroutine CONECT also connects these connections.

### 2.3.3 Propagation of CSPFE's

As discussed in Section 2.3.1, if no actual substitution has taken place and no error-component has been compensated during the error-compensation concerning a selected gate GI, the procedure will propagate the CSPFE of GI to the input connections of GI and will select another gate.

The propagation of CSPFE's is a procedure for calculating CSPFE's. As explained in Section 2.1 of [5], the CSPFE of a particular gate depends on the CSPFE's of its output connections, and the operation for calculating the CSPFE of a gate from the CSPFE's of its output connections is commutative (i.e., the order in which connections are considered is not important). In the procedure realized by subroutine RCEC, the CSPFE for each connection is not stored. Instead, the intermediate CSPFE for each gate is stored, and it will remain an intermediate CSPFE for that gate until all its output connections have been considered. The propagation of the CSPFE of gate GI is essentially the calculation of CSPFE's of its immediate predecessors contributed by gate GI.

Since there are five types of components in the CSPFE, the propagation can be classified into the following five categories.

(1) Don't-care components (\*)

The corresponding input components can be either 0 or 1. Since it is a don't-care component, no change is required for the intermediate CSPFE's of the immediate predecessors of gate GI.

(2) 1-components

A 1-component is realized when all corresponding input components are 0. Since the 1-component is required for the CSPFE of gate GI, all the corresponding components in the CSPFE's of GI's immediate predecessors should be assigned to 0 unless the input is from an output gate and the corresponding component has been assigned as a 0-error-component.

(3) 1-error-components

Since a 1-error-component in the CSPFE of GI indicates the actual output is 1 but the preferred value is 0, it should be propagated to its

immediate predecessors as 0-error-components if possible. If the corresponding component in the intermediate CSPFE of one of GI's immediate predecessors has already been assigned to 0 before this propagation, it should remain 0 since the assigned 0 means that an earlier propagation from another gate requires that it be 0.

#### (4) 0-error-components

In this case, the corresponding components of the inputs of GI could be either 0 or 1. If the component of an input corresponding to a 0-error-component in the CSPFE of GI is 0, the corresponding component of the CSPFE of that input is assigned to 0 in order not to increase the degree of the 0-error-component under consideration. On the other hand, if the component corresponding to a 0-error-component is a 1, it should be assigned as a 1-error-component unless it has already been assigned as a 1-component in an earlier step.

#### (5) 0-components

The propagation of 0-components is the most complicated case among those of five different types of components. For each 0-component in the CSPFE of GI, only one corresponding component of GI's inputs should be assigned to 1, and all others need not change. If a 0-component is covered by an external variable or a gate whose corresponding component in its CSPFE has already been assigned to 1, the above condition has already been satisfied and therefore no other changes are necessary. Since the error-compensation procedure depends very much on how to choose the 1-component to cover a 0-component, the ordering is carefully calculated according to the following.

##### (a) Inputs which covers fewer primary 0-error-components in the

CSPFE of GI have higher priority in the selection.

(b) If there is a tie based on (a), the input which covers fewer 0-error-components has the higher priority in the selection.

(c) If two inputs have the same priority based on (a) and (b), the one which has more output connections has the higher priority.

Rules (a) and (b) are based on the consideration of easier error-compensation for the immediate predecessors of GI. Since a primary 0-error-component in the CSPFE of GI can be compensated if the 1-error-component covering it is compensated, every effort should be made to make this 1-error-component easier to be compensated. Rules (a) and (b) tend to assign more don't-care-components to the inputs which cover more primary 0-error-components. This explains rules (a) and (b) since the more don't-care-components there are in the CSPFE of a gate, the more connectable functions are likely, and therefore the more chances there are to compensate for these error-components. Rule (c) is based on the number of output connections. If a gate which has more output connections than others is selected and assigned to 1, this 1 is likely to cover 0-components in other immediate successors of this gate.

In the actual program, the priority depends on the weight associated with each input connection which is calculated according to the following formula.

$$W = 10^6 - 10^4 * (\# \text{ primary error-components}) \\ - 10^2 * (\# \text{ error-components}) + (\# \text{ immediate successors})$$

A special case should be noted. If a 0-component is covered only by an output gate of the network and the corresponding component in the CSPFE of that gate has already been assigned as a 1-error-component, this

component should not be changed and the corresponding components of other inputs should be assigned as 0-error-components, if possible, in order to cover the 0-component in the CSPFE of GI when the corresponding 1-error-component in the CSPFE of the output gate is compensated.

#### 2.3.4 Algorithm of error-compensation procedure and flowchart

As a summary of the above discussions, a brief description of the procedure realized by subroutine RCEC is presented in this section. For the detail of the procedure, however, see the flowchart in Figure 2.3.4.1.

#### The Procedure of Error Compensation Realized by RCEC

##### Step 1. Selection of gate

Select a gate GI according to the ascending order of the level number assigned to each gate, i.e., a gate in the lowest level is selected first. If all gates have been considered, return (RETURN 1) to the calling subroutine (PROCCE in NETTRA-E1, -E2, and -E3)(this is an unsuccessful compensation).

##### Step 2. Removal of redundant connections

For each immediate predecessor GJ of GI, call subroutine FORC(GJ) to check whether or not the connection from GJ to GI is redundant. If so, remove this connection.

If no error-component is in the CSPFE of GI, go to Step 9.

##### Step 3. Selection of connectable functions from POT

Select strongly effectively E-connectable functions for GI from the potential output table. If there is no such connectable function for GI, go to Step 9.



#### Step 4. Classification of connectable functions

Classify connectable functions for GI into four categories: (1) external variables without error-components (set DIO), (2) gates without anticipated error-components (set BIO), (3) gates with anticipated error-components (set BI), and (4) external variables with error-components (set DI). Sort BI according to the number of 0-error-components the function covers.

#### Step 5. Substituting for external variables with errors

Let  $S2 = DIO \cup BIO \cup BI$ ,  $S$  = external variable inputs of GI, which cover 0-error-components (these inputs are found and sorted by calling subroutine ORDRQ2). Call subroutine CALS1 and RPLCF to substitute a subset  $S3$  of  $S2$  for a subset  $S1$  of  $S$ , if possible.

#### Step 6. Substituting for functions covering primary error-components

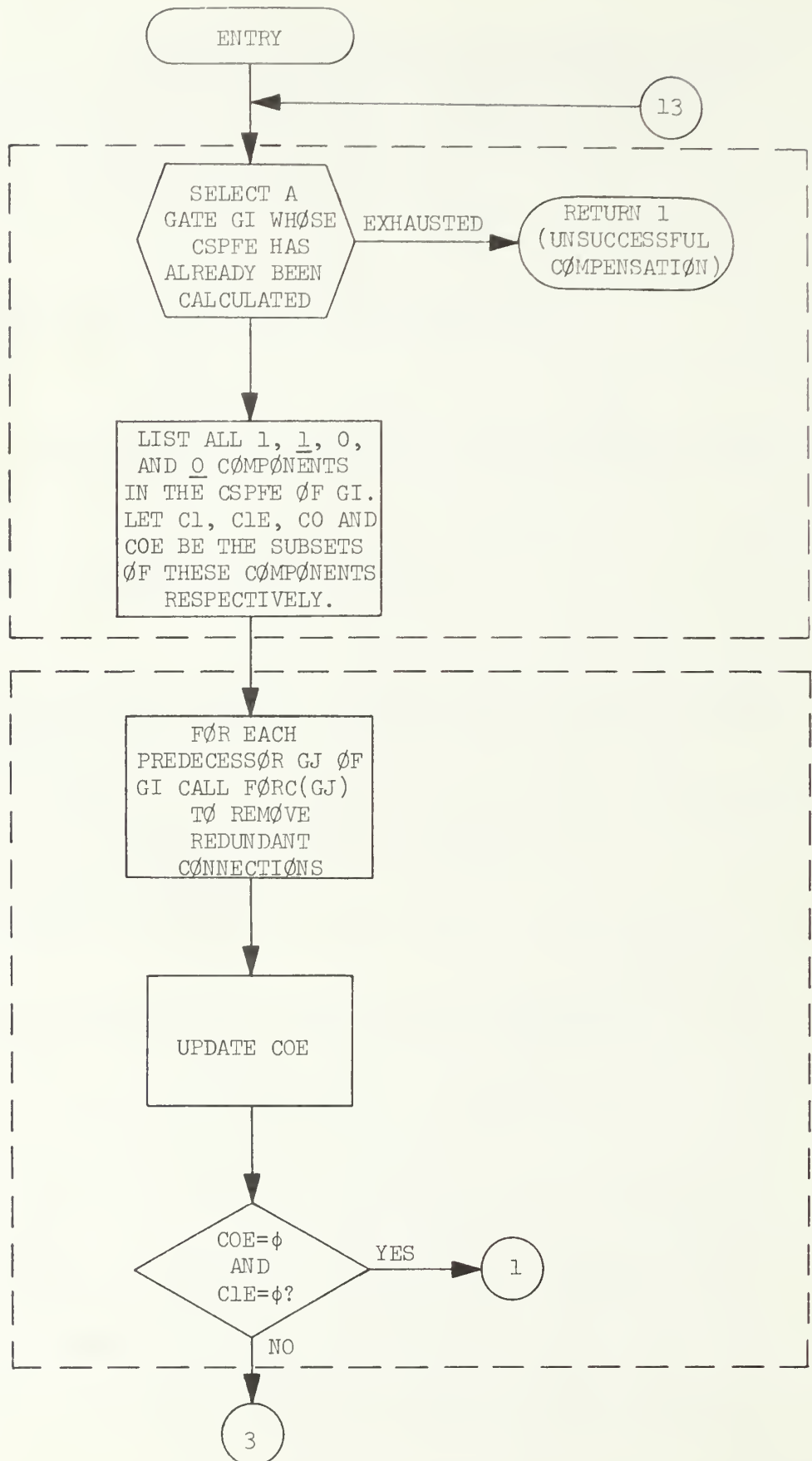
For each input of GI covering primary error-components (these inputs are found and sorted by calling ORDRQ2 in Step 5), check whether or not it can be replaced by functions in DIO, BIO and BI with some primary error-components corrected. Call CALS1 and RPLCF to complete this substitution.

#### Step 7. Substituting for functions by functions having no error-components

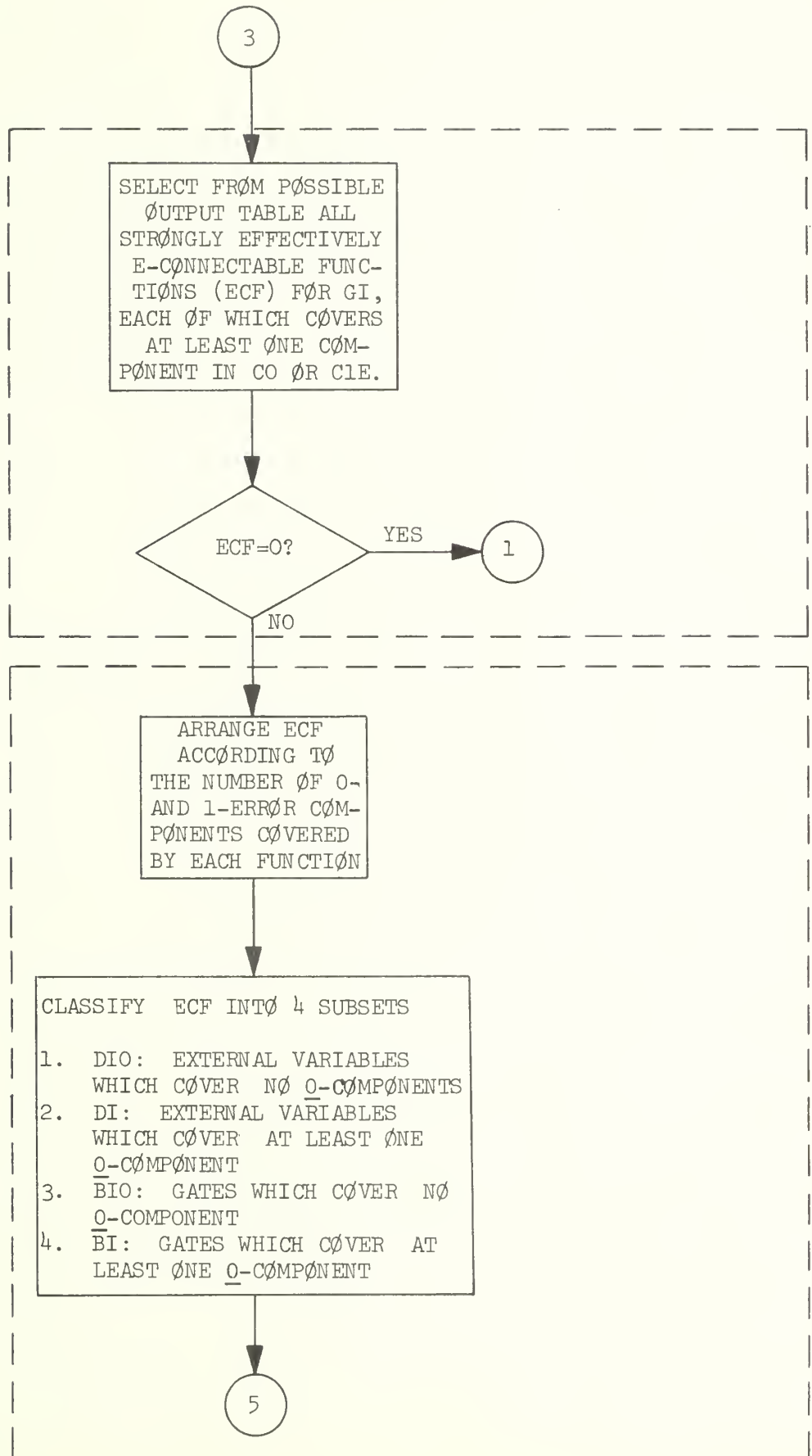
Let  $S$  be the set of remaining inputs of GI which have error-components, and let  $S2$  be the set of the remaining functions (exclude already selected and/or prohibited ones) in sets DIO and BIO. Call CALS1 and RPLCF to replace as subset  $S1$  of  $S$  by functions in a subset  $S3$  of  $S2$ , if possible.

#### Step 8. Compensating for 1-error-components in the CSPFE of GI

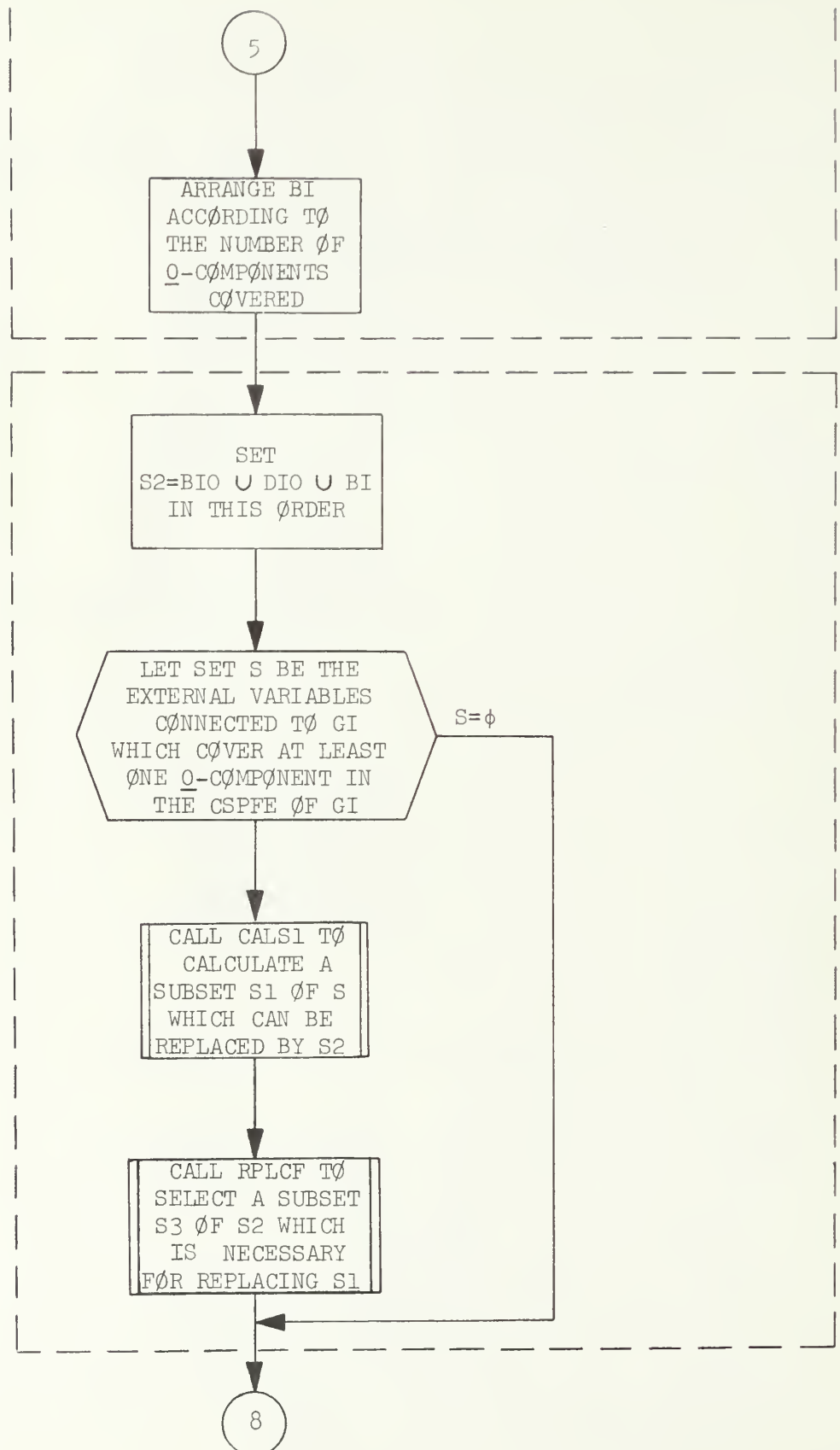
If some of the remaining functions in DIO, BIO and BI have 1-components corresponding to some 1-error-components in the CSPFE of GI,



Step 3

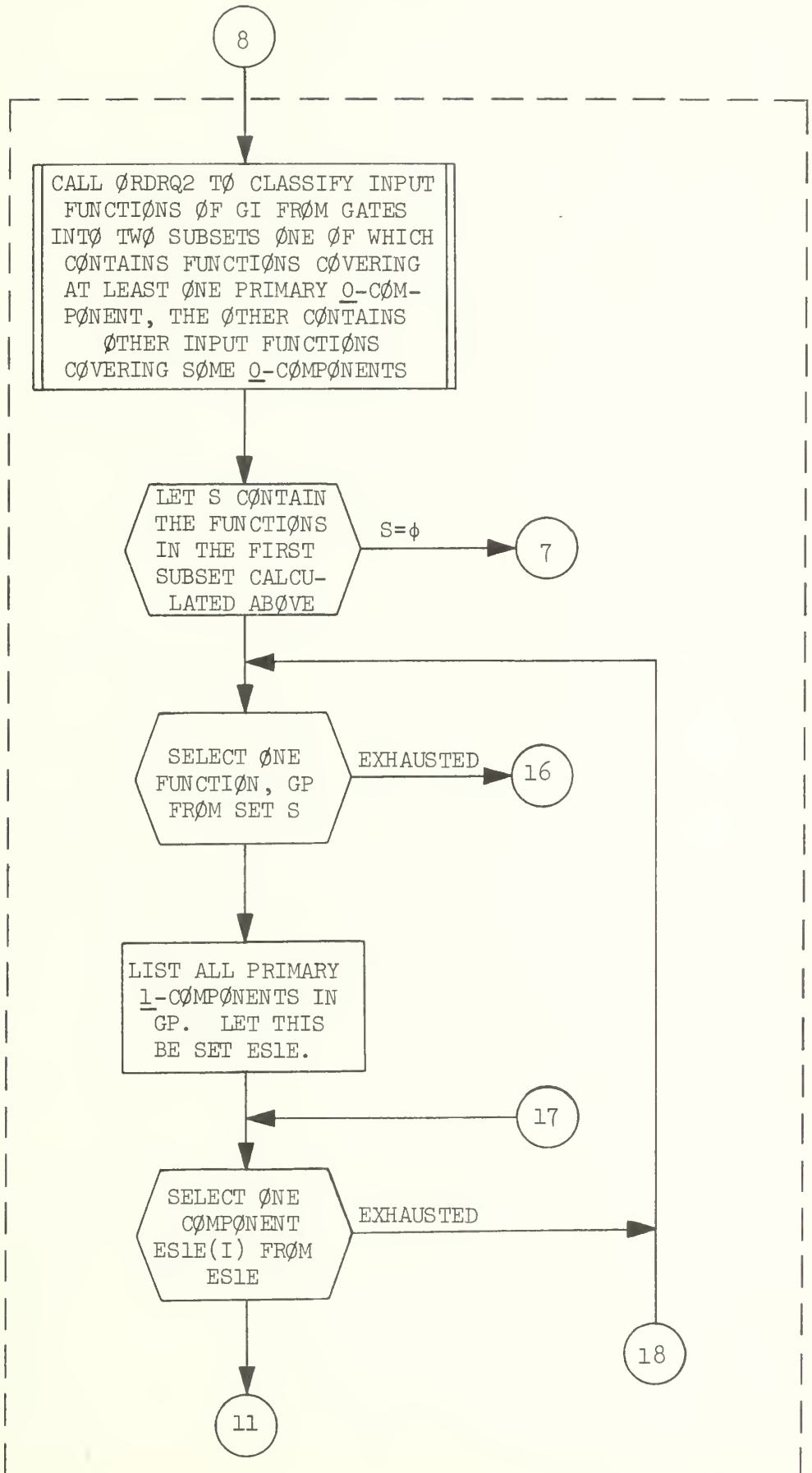


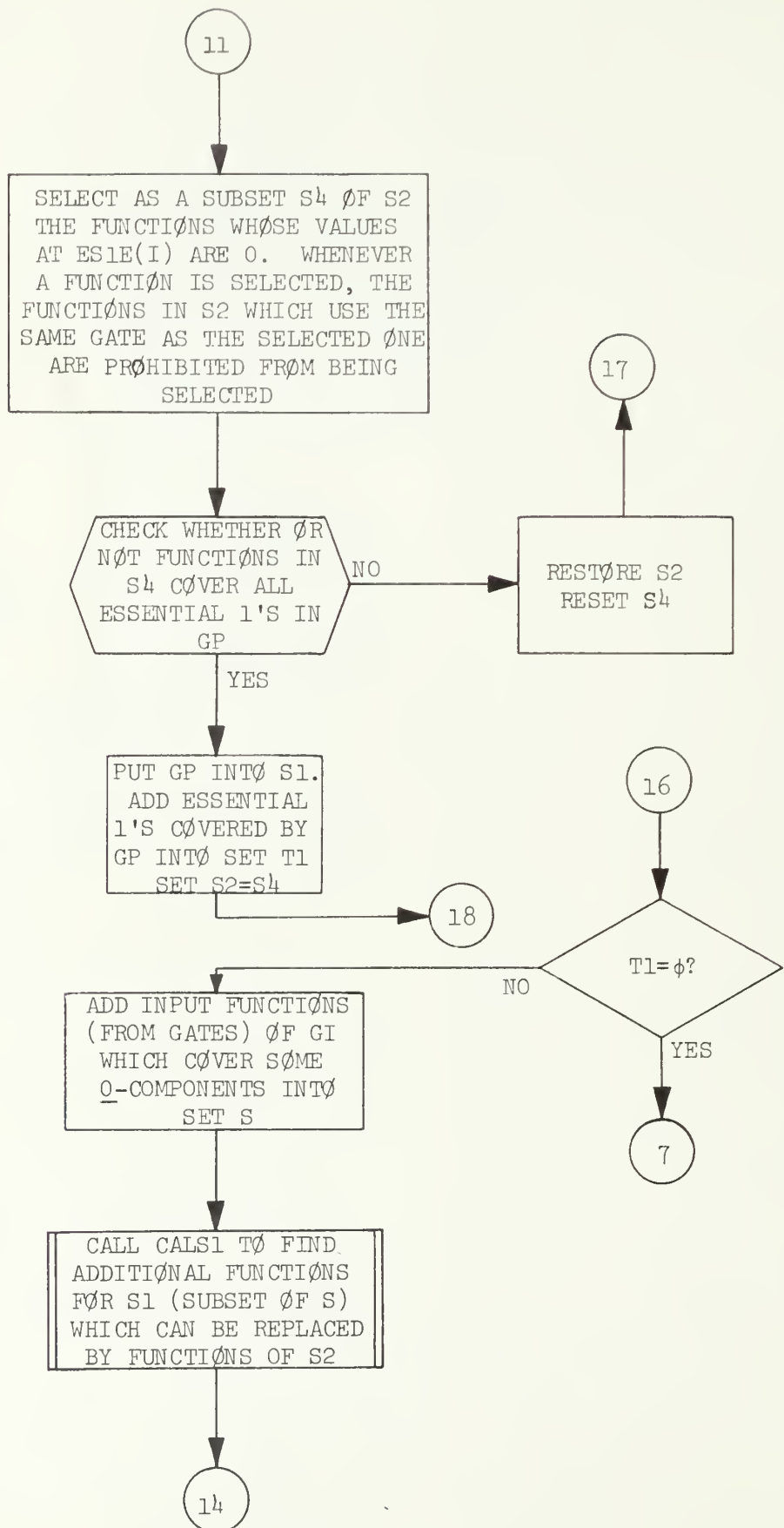
Step 4



Step 5

Step 6

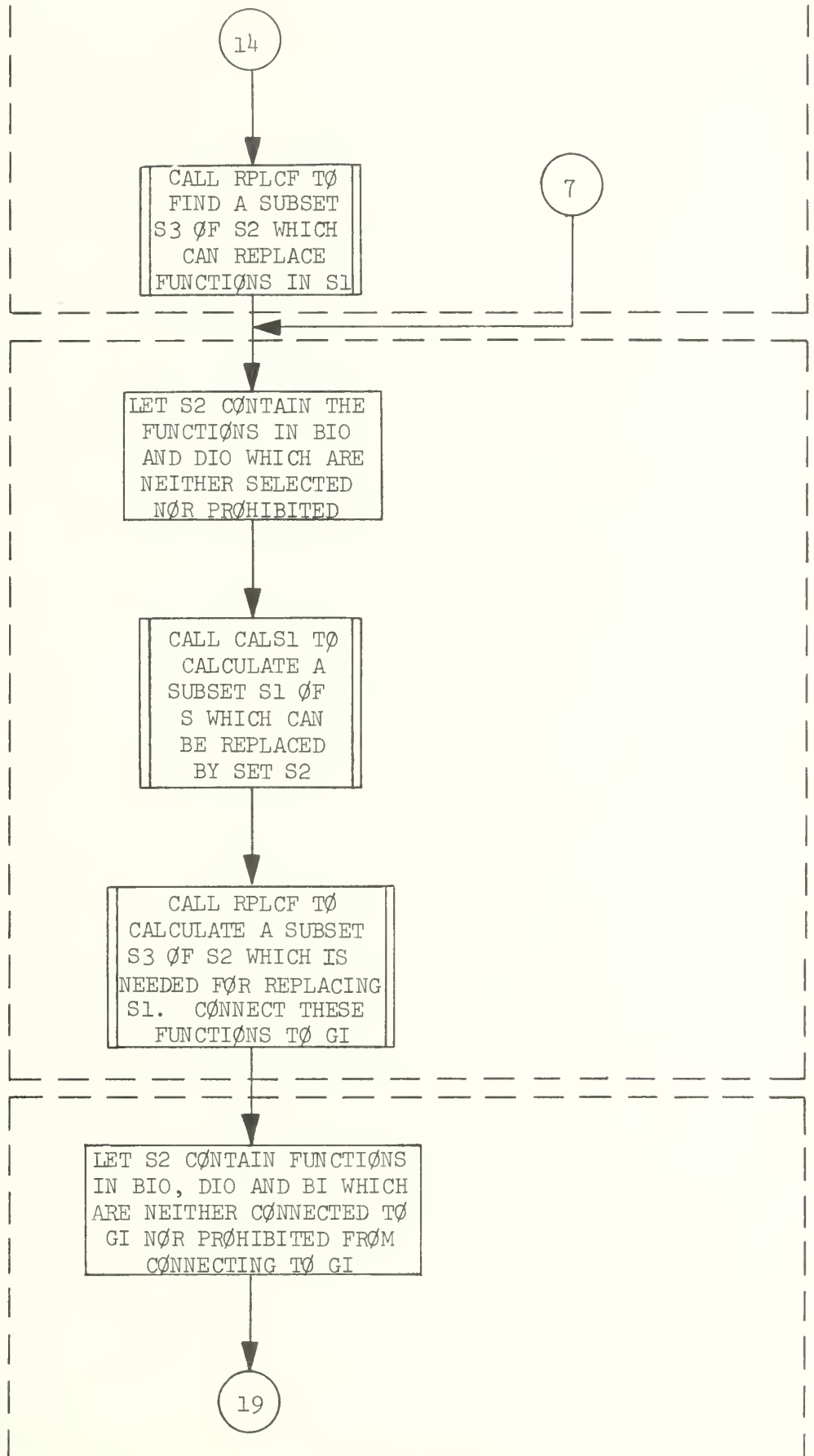


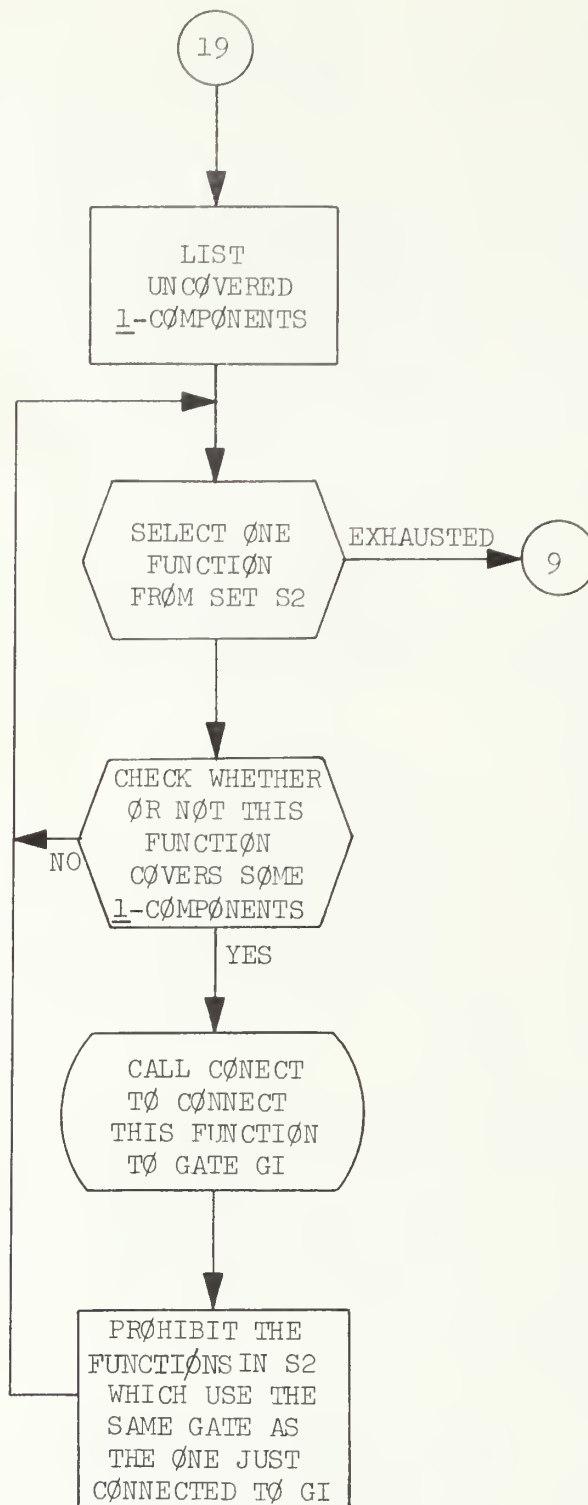




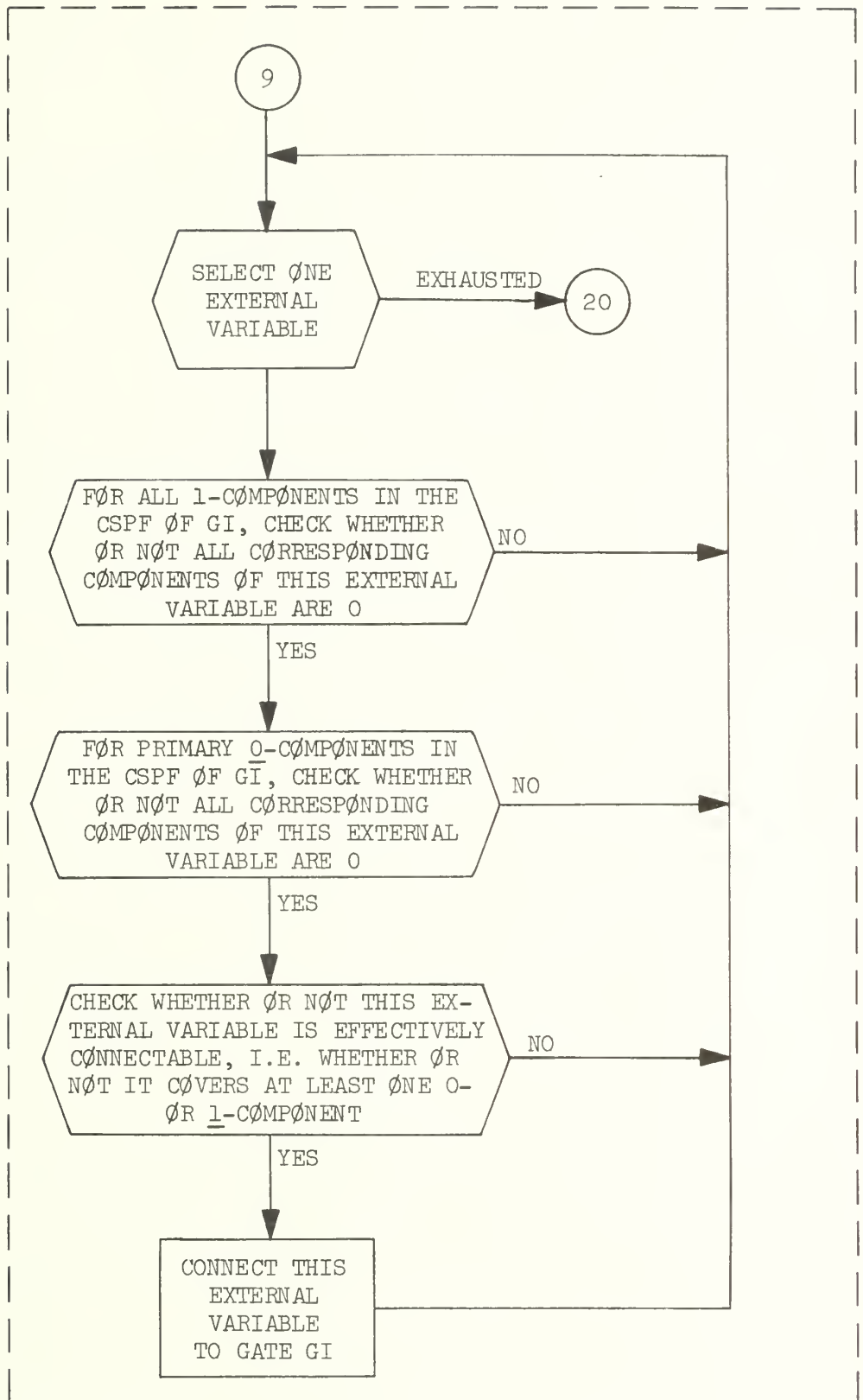
Step 7

Step 8

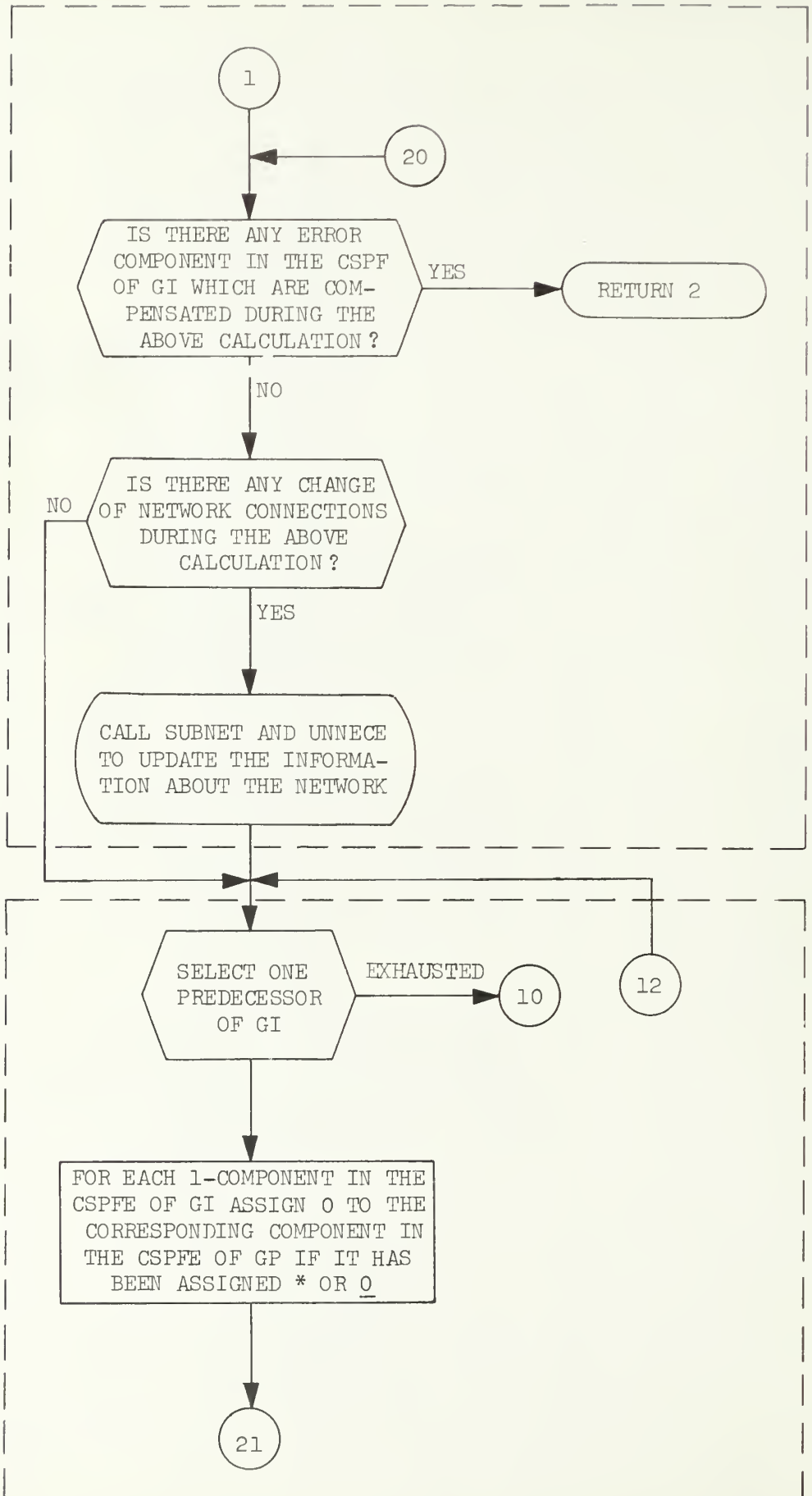




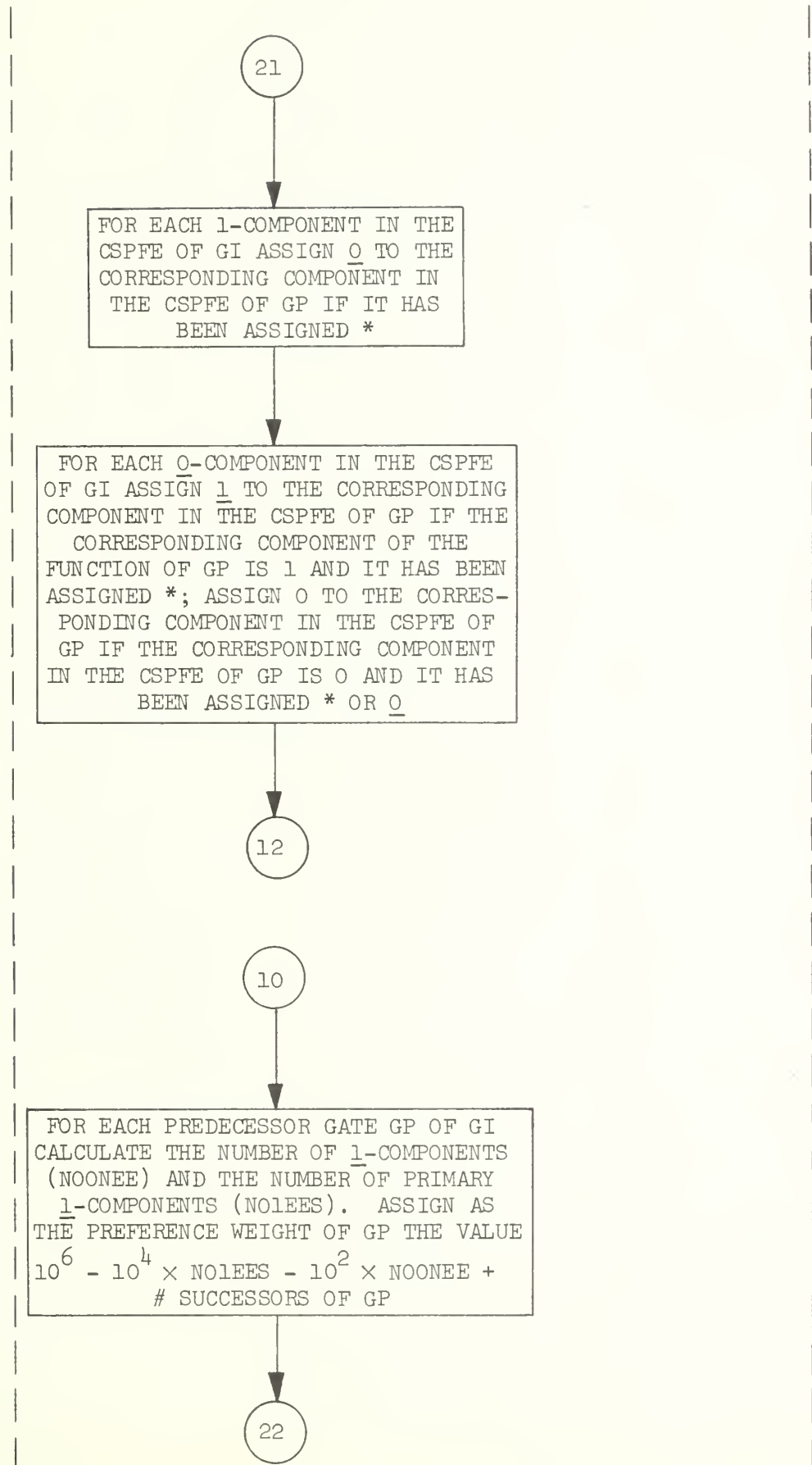
Step 9

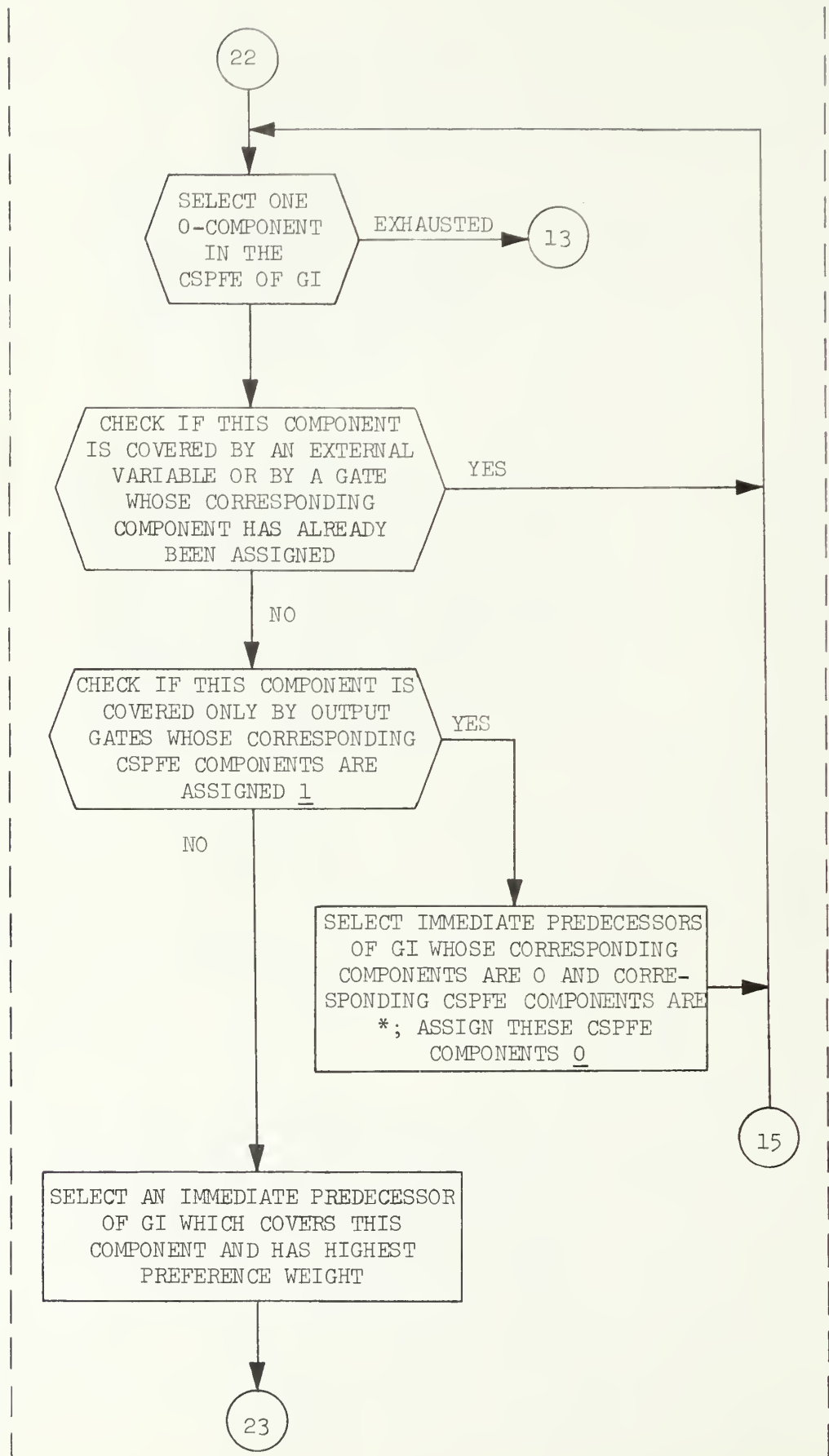


Step 10



Step 11







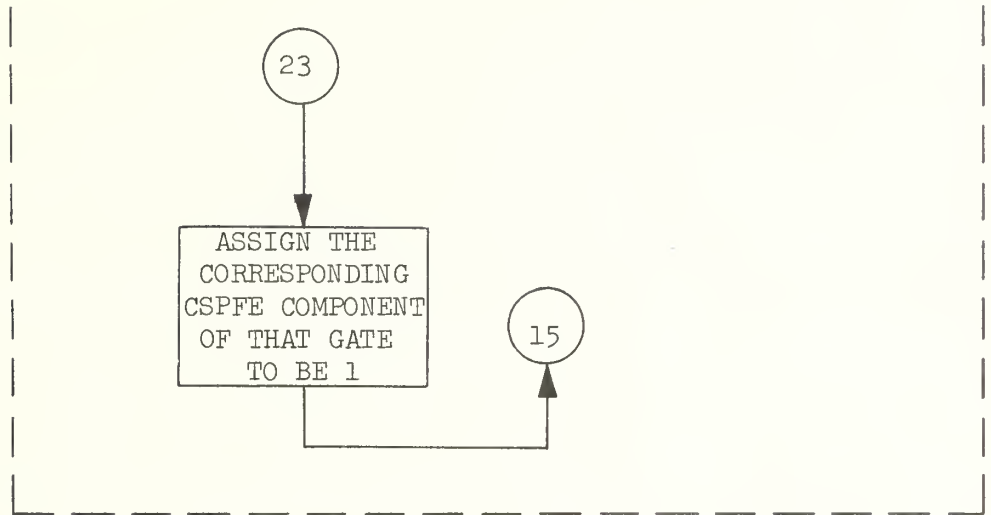


Figure 2.3.4.1 Flowchart of Subroutine RCEC (Dashed blocks are the steps 1, 2, ..., 11 corresponding to those of the error compensation procedure).

connect these functions to compensate for 1-error-components.

#### Step 9. Adding redundant external variables

Connect connectable redundant external variables to GI without reducing the number of primary error-components in the CSPFE of GI.

#### Step 10. Return to calling subroutine

If an error-component in the CSPFE of GI has been corrected or a substitution is performed during steps 2, 5, 6, 7 or 8, return (RETURN 2) to the calling subroutine (PROCCE in NETTRA-E1, -E2, and -E3) to check the outputs of the network.

#### Step 11. Propagation of CSPFE

Update the intermediate CSPFE for each immediate predecessor of gate GI. Go to Step 1.

The relation between subroutine RCEC and its supporting subroutines is shown in Figure 2.1 (in Section 2).

### 2.4 Example for NETTRA-E1

The printout obtained during the solution of a typical problem by NETTRA-E1 is shown in Figure 2.4.1. The original network, as specified in the beginning of the printout (Figure 2.4.1(a)), consists of 19 gates and 76 connections and realizes a single 5-variable output function. Only uncomplemented variables are available as inputs to the network.

Following this information is printed a complete truth table (b) showing the output of every gate in the original network for every possible input combination. Note that it is gate 1 which realizes the output function of the network.

\*\*\* 3-LEVEL NETWORK \*\*\*

\*\*\*\*\* 5 VARIABLE, 1 OUTPUT TEST NETWORK NUMBER 25

NUMBER OF VARIABLES = 5

NUMBER OF FUNCTIONS = 1

COST COEFFICIENT A = 100

B = 1

--- UNCOMPLEMENTED VARIABLES X ---

FUNCTION NO. 1.

10101000101010100101110011101000

ORIGINAL NETWORK COST=19076

(a) Heading information and network parameters.

Figure 2.4.1 Printout obtained from NETTRA-El for a sample problem.

## TRUTH TABLE

```

X1 = 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
X2 = 0 0 0 0 0 0 0 0 0 1 1 1 1 1 1 1 1 0 0 0 0 0 0 0 0 0 1 1 1 1 1 1
X3 = 0 0 0 0 1 1 1 1 0 0 0 0 0 1 1 1 1 0 0 0 0 0 1 1 1 1 0 0 0 0 1 1
X4 = 0 0 1 1 0 0 1 1 0 0 1 1 0 0 1 1 0 0 1 1 0 0 1 1 0 0 1 1 0 0 1 1
X5 = 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1
  1 = 1 0 1 0 1 0 0 0 0 1 0 1 0 1 0 1 0 0 1 0 1 1 1 0 0 1 1 1 0 1 0 0
  2 = 1 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
  3 = 0 1 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
  4 = 1 0 1 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
  5 = 0 1 0 1 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
  6 = 1 0 0 0 1 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
  7 = 0 1 0 0 0 1 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
  8 = 0 0 0 0 0 0 1 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
  9 = 1 0 0 0 0 0 0 0 0 1 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
 10 = 0 1 0 0 0 0 0 0 0 0 1 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
 11 = 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 1 0 0 0 0 0 0 0 0 0 0 0
 12 = 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 1 0 1 0 0 0 0 0 0 0 0
 13 = 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 0 0 0 0 0 0 0 0 0 0 0 0 0
 14 = 1 1 1 1 0 0 0 0 0 0 0 0 0 0 0 0 0 0 1 1 1 1 0 0 0 0 0 0 0 0 0
 15 = 1 1 0 0 1 1 0 0 0 0 0 0 0 0 0 0 0 0 1 1 0 0 1 1 0 0 0 0 0 0 0
 16 = 1 1 0 0 0 0 0 0 0 1 1 0 0 0 0 0 0 0 1 1 0 0 0 0 0 0 1 1 0 0 0
 17 = 1 0 1 0 0 0 0 0 0 1 0 1 0 0 0 0 0 0 1 0 1 0 0 0 0 0 1 0 1 0 0
 18 = 1 0 0 0 1 0 0 0 0 1 0 0 0 0 1 0 0 0 1 0 0 0 1 0 0 0 1 0 0 0 0
 19 = 0 0 0 0 0 0 0 0 0 1 0 0 0 0 1 0 1 0 1 0 0 0 0 0 0 1 1 0 0 0 1

```

(b) Truth table for original network.

GATE	..	LEVEL	FED BY									
1	/	1/	3	5	7	8	10	11	12	19		
2	/	3/	X1	X2	X3	X4	X5					
3	/	2/	X1	X2	X3	X4	2					
4	/	3/	X1	X2	X3	X5						
5	/	2/	X1	X2	X3	4						
6	/	3/	X1	X2	X4	X5						
7	/	2/	X1	X2	X4	6						
8	/	2/	X1	X2	X5	2	4	6				
9	/	3/	X1	X3	X4	X5						
10	/	2/	X1	X3	X4	9						
11	/	2/	X2	X3	X4	X5	2					
12	/	2/	X2	X3	X5	2	4					
13	/	3/	X1	X5								
14	/	3/	X2	X3								
15	/	3/	X2	X4								
16	/	3/	X3	X4								
17	/	3/	X3	X5								
18	/	3/	X4	X5								
19	/	2/	13	14	15	16	17	18				

(c) Configuration of original network.

\*\*\*\* BEGIN 1-TH APPLICATION OF PROCCE : \*\*\*\*\*

NETWORK DERIVED BY PROCCE  
TIME ELAPSED = 14 CENTISECONDS

# TRUTH TABLE

```

X1 = 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
X2 = 0 0 0 0 0 0 0 0 0 1 1 1 1 1 1 1 1 0 0 0 0 0 0 0 0 0 1 1 1 1 1 1
X3 = 0 0 0 0 1 1 1 1 0 0 0 0 1 1 1 1 0 0 0 0 1 1 1 1 0 0 0 0 1 1 1 1
X4 = 0 0 1 1 0 0 1 1 0 0 1 1 0 0 1 1 0 0 1 1 0 0 1 1 0 0 1 1 0 0 1 1
X5 = 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1
 1 = 1 0 1 0 1 0 0 0 1 0 1 0 1 0 1 0 1 0 0 1 0 1 1 1 0 0 1 1 1 0 1 0 0 0
 2 = 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
 3 = 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
 4 = 1 0 1 0 0 0 0 0 1 0 1 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
 5 = 0 1 0 1 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
 6 = 1 0 0 0 1 0 0 0 1 0 0 0 1 0 0 0 1 0 0 0 1 0 0 0 1 0 0 0 1 0 0 0 0
 7 = 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
 8 = 0 1 0 1 0 1 1 1 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
 9 = 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0
10 = 0 1 0 0 0 0 0 0 0 1 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
11 = 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
12 = 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 1 0 1 0 0 0 0 0 0 0 0 0 0 0 0
13 = 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
14 = 1 1 1 1 0 0 0 0 0 0 0 0 0 0 0 0 0 1 1 1 1 0 0 0 0 0 0 0 0 0 0 0
15 = 1 1 0 0 1 1 0 0 0 0 0 0 0 0 0 0 0 1 1 0 0 1 1 0 0 0 0 0 0 0 0 0
16 = 1 1 0 0 0 0 0 0 1 1 0 0 0 0 0 0 1 1 0 0 0 0 0 0 1 1 0 0 0 0 0 0
17 = 1 0 1 0 0 0 0 0 1 0 1 0 0 0 0 0 1 0 1 0 0 0 0 0 1 0 1 0 0 0 0 0
18 = 1 0 0 0 1 0 0 0 1 0 0 0 1 0 0 0 1 0 0 0 1 0 0 0 1 0 0 0 1 0 0 0
19 = 0 0 0 0 0 0 0 1 0 0 0 1 0 1 0 1 0 0 0 0 0 0 1 1 0 0 0 1 0 1 1 1

```

(d) Truth table for transformed, reduced network.



GATE .. LEVEL	FED BY
1 / 1/	5 8 10 12 19
2 / 1/	
3 / 1/	
4 / 3/	X1 X3 X5
5 / 2/	X1 X2 X3 4
6 / 3/	X4 X5
7 / 1/	
8 / 2/	X1 X2 4 6
9 / 3/	X5
10 / 2/	X1 X3 X4 9
11 / 1/	
12 / 2/	X2 X3 X5 4
13 / 3/	X1 X5
14 / 3/	X2 X3
15 / 3/	X2 X4
16 / 3/	X3 X4
17 / 3/	X3 X5
18 / 3/	X4 X5
19 / 2/	13 14 15 16 17 18

\* A NETWORK DERIVED BY PROCCE  
COST=15045.

(e) Configuration of transformed, reduced network.

Next appears a description of the configuration of the network (c). Each gate is listed along with the gates and/or external variables which are its inputs. The level numbers, which can also be seen in (c), will be discussed later in Section 5.3.

The truth table (note that the outputs for disconnected gates are shown as all 1's) and network configuration for the transformed, reduced network resulting from the action of NETTRA-E1 are shown in (d) and (e), respectively. The derived network, found in .14 seconds, consists of 15 gates and 45 connections. If NETTRA-E1 were applied to this new network, a network with even fewer gates and connections would be obtained.

### 3. REPETITIVE APPLICATION OF ERROR-COMPENSATION PROCEDURE

Now that it has been explained how the error-compensation procedure (as used in NETTRA-E1) is often able to remove a gate from a network (containing an excessive number of gates) and alter the remaining network configuration so as to again produce the desired functions, the most obvious extension can be made to produce an even better program to reduce the number of gates in a network. If the procedure can remove one gate, it may be able to remove a second, or a third. Thus the programs NETTRA-E2 and -E3 are introduced in the next sections: 3.1 and 3.2.

#### 3.1 Single-Path Application

The storage requirements of NETTRA-E2 are almost identical to those given for NETTRA-E1 since the programs only differ by a few FORTRAN statements in the subroutine MAIN. NETTRA-E2 requires 163K bytes of core storage, about 78K for the actual program instructions and about 85K for the stored data.

The following subroutines, written in FORTRAN IV for the IBM 360/75, constitute the program NETTRA-E2: CALS1, CONECT, FORC, MAIN (this differs from the MAIN in NETTRA-E1), MINI2, ORDRQ2, OUTPUT, POT, PROCCE, RCEC, RPLCF, and SUBNET. Two system-supplied timing subroutines, STIMEZ and KTIMEZ are also assumed to be available, but if they are not, their use can be omitted from the program, or another suitable timing routine substituted, without

harming the procedure itself.

The general organization of the program is identical to that of NETTRA-E1. It is shown in Figure 2.1. In this figure, an arrow from block i to block j represents the fact that the subroutine in block i calls the subroutine in block j.

It is subroutine MAIN which actually performs the repeated calls to subroutine PROCCE, removing one gate after another until PROCCE can no longer produce a reduced network.(i.e., a network with fewer gates).

The setup of the input data for NETTRA-E2 is the same as the setup for NETTRA-E1 and -E3. The details can be found in section 5. Listings of all of the subroutines used in NETTRA-E2 can be located in the appendix.

### 3.1.2 Example of NETTRA-E2

Beginning with the same initial network described by Figure 2.4.1 (a), (b), and (c), NETTRA-E2 is applied. Since NETTRA-E2 is essentially just repetitive applications of NETTRA-E1, the first portion of the printout obtained by NETTRA-E2 is identical to what appears in Figure 2.4.1 (a) - (e).

Following that, PROCCE is called a second, third, fourth and fifth time, each time finding a transformed, reduced network of fewer gates than before (except that the fifth result is not improved over the fourth) and printing out the new truth table and network configuration (as in Figure 2.4.1 (d) and (e)).

For this example, NETTRA-E2 could not find a better network than that obtained in the fourth application of PROCCE. The truth table for this new network is given in Figure 3.1.2.1 (a) and the transformed, reduced network configuration description is shown in Figure 3.1.2.1 (b) as they appear in the

\*\*\*\* BEGIN 4-TH APPLICATION OF PROCCE :

\*\*\*\*\*

NETWORK DERIVED BY PROCCE

TIME ELAPSED = 94 CENTISECONDS

# TRUTH TABLE

```

X1 = 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
X2 = 0 0 0 0 0 0 0 0 1 1 1 1 1 1 1 1 0 0 0 0 0 0 0 0 1 1 1 1 1 1 1 1
X3 = 0 0 0 0 1 1 1 1 0 0 0 0 1 1 1 1 0 0 0 0 1 1 1 1 0 0 0 0 1 1 1 1
X4 = 0 0 1 1 0 0 1 1 0 0 1 1 0 0 1 1 0 0 1 1 0 0 1 1 0 0 1 1 0 0 1 1
X5 = 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1
 1 = 1 0 1 0 1 0 0 0 1 0 1 0 1 0 1 0 1 0 1 1 1 0 0 1 1 1 0 1 0 0 0
 2 = 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
 3 = 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
 4 = 1 1 1 1 0 0 0 0 1 1 1 1 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
 5 = 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
 6 = 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
 7 = 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
 8 = 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
 9 = 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 0 0 0 0 0 0 0 0 0 0 0 0
10 = 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 0 0 0 0 0 0 0 0 0 0 0
11 = 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
12 = 0 0 0 0 0 0 1 0 0 0 0 0 0 0 0 0 1 0 1 0 0 0 1 0 0 0 0 0 0 0
13 = 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
14 = 1 1 1 1 0 0 0 0 0 0 0 0 0 0 0 0 1 1 1 1 0 0 0 0 0 0 0 0 0 0
15 = 1 1 0 0 1 1 0 0 0 0 0 0 0 0 0 0 1 1 0 0 1 1 0 0 0 0 0 0 0 0
16 = 1 1 0 0 0 0 0 0 1 1 0 0 0 0 0 0 1 1 0 0 0 0 0 0 1 1 0 0 0 0
17 = 1 0 1 0 0 0 0 0 1 0 1 0 0 0 0 0 1 0 1 0 0 0 0 0 1 0 1 0 0 0
18 = 0 0 0 0 1 0 0 0 1 0 0 0 1 0 0 0 0 0 0 0 1 0 0 0 1 0 0 0 1 0
19 = 0 0 0 0 0 0 0 1 0 0 0 1 0 1 0 1 0 0 0 0 0 0 1 1 0 0 0 1 0 1 1 1

```

(a) Truth table for final network.

Figure 3.1.2.1 Printout of final transformed, reduced network obtained by NETTRA-E2.

GATE .. LEVEL	FED BY
1 / 1/	10 12 19
2 / 1/	
3 / 1/	
4 / 3/	X1 X3
5 / 1/	
6 / 1/	
7 / 1/	
8 / 1/	
9 / 3/	X1 X5
10 / 2/	X1 9
11 / 1/	
12 / 2/	X2 X5 4 18
13 / 1/	
14 / 4/	X2 X3
15 / 3/	X2 X4
16 / 3/	X3 X4
17 / 3/	X3 X5
18 / 3/	X4 X5 14
19 / 2/	9 14 15 16 17 18

\* A NETWORK DERIVED BY PROCCE  
COST=11030.

(b) Configuration of final network.



printout from NETTRA-E2. This final transformed, reduced network consists of only 11 gates and 30 connections.

### 3.2 Multi-Path Application

The program NETTRA-E3 represents what is called a multi-path application of the error-compensation procedure. It is only slightly more complicated than NETTRA-E2 (due to the necessity of storing intermediately produced networks in a stack) and employs only one additional subroutine.

NETTRA-E2 only produces a single sequence of networks beginning with the original network specified by the input data and ending with a network from which the error-compensation routine can successfully remove no more gates. Let this sequence of networks be labeled  $W_1 \rightarrow W_2 \rightarrow W_3 \rightarrow \dots \rightarrow W_E$ , where  $W_1$  is the original network and  $W_E$  is the final network. Each  $W_{i+1}$  is derived from  $W_i$  by the removal of a single gate (although other gates may consequently be removed) and the successful compensation of the resulting errors by the error-compensation routine. Actually, (in general) there are several different networks that can be obtained from each  $W_i$  by the removal of and successful compensation for several different gates. NETTRA-E2 settles for the first network,  $W_{i+1}$ , that can be obtained from  $W_i$  by the removal of a certain gate and the successful compensation of errors, but NETTRA-E3 stores that improved network ( $W_{i+1}$ ) and continues to search for additional improved networks that can be derived from  $W_i$  by removing different gates. These are also stored, in a stack, and NETTRA-E3 will then search for all of the possible improved networks obtainable from those in the stack, storing the new networks in the stack as they are obtained, and so on. Thus NETTRA-E3 will produce a "tree" of solutions

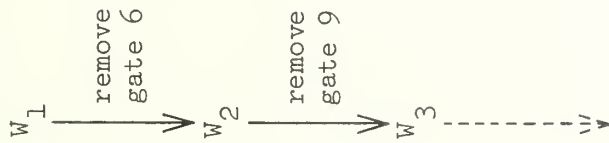
(intermediate and final) while NETTRA-E2 only produces a single "string" of solutions (which, incidentally, will be identical to a certain path through the tree produced by NETTRA-E3).

For example, Figure 3.2.1 shows the difference between the (intermediate and final) solutions obtained by NETTRA-E2 and the solutions obtained by NETTRA-E3.

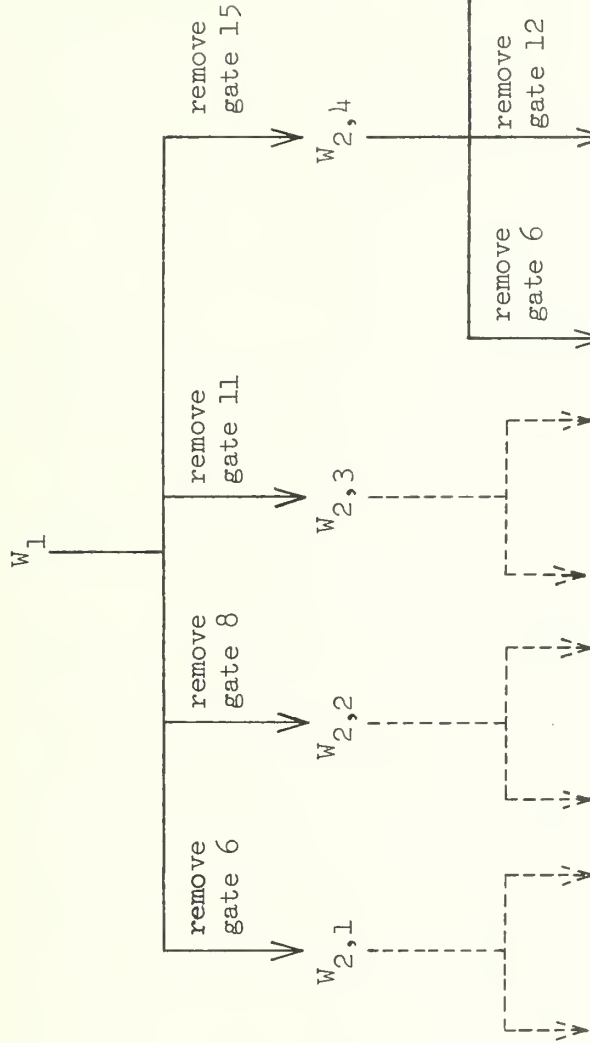
Beginning with  $W_1$  (assumed to be a network of 16 gates) NETTRA-E2 discovers that gate 6 (for example) can successfully be removed from  $W_1$  to produce the 15-gate network  $W_2$ . It immediately prints out this result (i.e., the configuration of  $W_2$ ) and proceeds to "operate" on  $W_2$ , trying to remove some gate in order to obtain a network  $W_3$  of 14 gates.

However, beginning with the same original network,  $W_1$ , NETTRA-E3 also finds that gate 6 can be successfully removed to produce a 15-gate network (let it be called  $W_{2,1}$  this time), but it does not "forget about"  $W_1$  and immediately attempt to transform  $W_{2,1}$ . Instead, it prints the result,  $W_{2,1}$  (along with a message identifying the "parent" [i.e.,  $W_1$ ] of the network), stores it in a stack memory, and continues to search for other 15-gate networks obtainable from  $W_1$  by the removal of different gates (these are represented by  $W_{2,2}$  [obtained by removing gate 8 from  $W_1$ ],  $W_{2,3}$  [removing gate 11], and  $W_{2,4}$  [gate 15]). These networks,  $W_{2,2}$ ,  $W_{2,3}$ , and  $W_{2,4}$  are also put into the stack. NETTRA-E3 then selects the top network of the stack ( $W_{2,4}$  in this example) and attempts to remove individual gates from it in the same manner as it treated  $W_1$ . This process continues until the stack is empty. The terminal nodes of the "solution tree" must then be searched to identify the best solution (i.e., the network with the fewest number of gates).

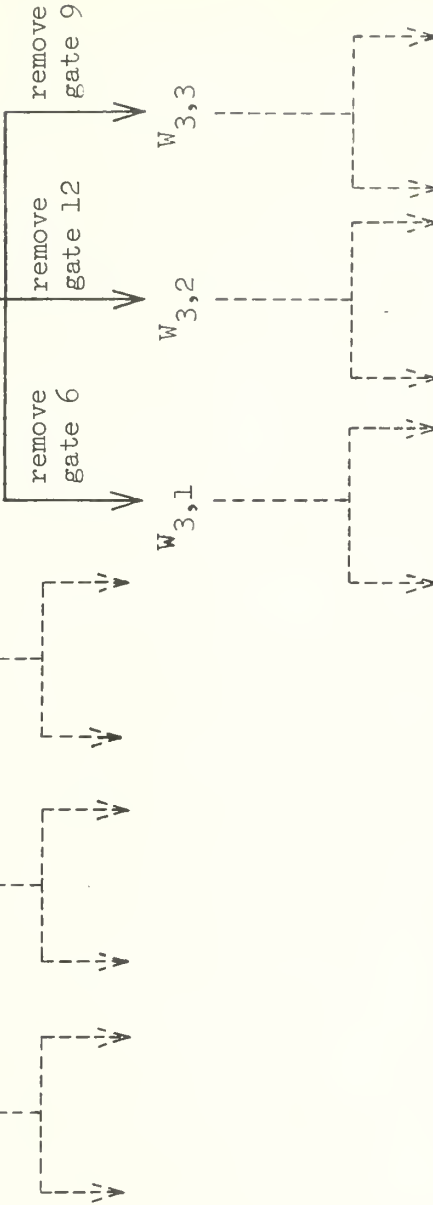
16-gate networks →



15-gate networks →



14-gate networks →



13-gate networks →

•  
•  
•

Figure 3.2.1 Typical intermediate solutions produced by NETTRA-E2 and -E3.

While NETTRA-E3 will obviously produce a solution as good as the best solution produced by NETTRA-E2, it is quite possible that it might produce an even better solution. Of course NETTRA-E3 can require much more computation time than NETTRA-E2, so each program has an advantage over the other, depending on the intended application.

### 3.2.1 Program organization of NETTRA-E3

NETTRA-E3 needs more storage space than NETTRA-E1 or -E2 due mainly to the addition of the previously mentioned stack. The program requires 207K bytes of core storage, about 82K for the actual program instructions and about 125K for the stored data.

The following subroutines, written in FORTRAN IV for the IBM 360/75, constitute the program NETTRA-E3: ALPATH, CALS1, CONECT, FORC, MAIN (this differs from the MAIN in NETTRA-E1 or -E2), MINI2, ORDRQ2, OUTPUT, POT, PROCCE (this differs from the PROCCE found in both NETTRA-E1 and -E2), RCEC, RPLCF, and SUBNET. Two system-supplied timing subroutines, STIMEZ AND KTIMEZ are also assumed to be available, but if they are not, their use can be omitted from the program, or another suitable timing routine substituted, without harming the procedure itself.

Figure 3.2.1.1 illustrates the general organization of the program. In this figure, an arrow from block i to block j represents the fact that the subroutine in block i calls the subroutine in block j.

In NETTRA-E3, subroutine MAIN calls the subroutine ALPATH only once. After that, it is ALPATH that controls the application of the error-compensation procedure to the networks stored in the stack. The modified version of PROCCE stores the intermediate solutions (networks) in the stack in the order

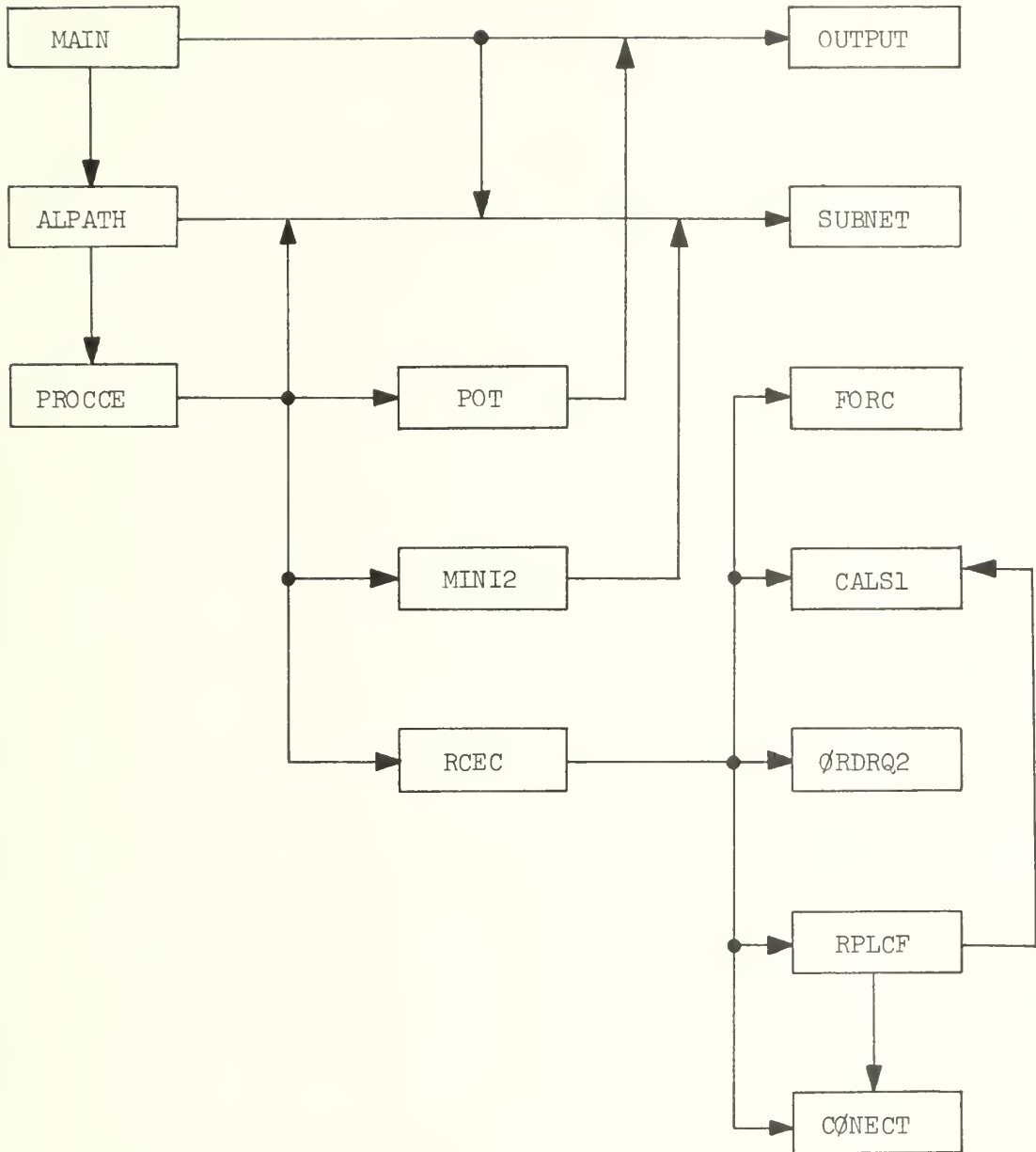


Figure 3.2.1.1 General organization of the program NETTRA-E3.

in which it obtains them. NETTRA-E3 terminates when the stack empties.

Once again, the input data setup is the same as that for NETTRA-E1 and -E2. The details can be found in Section 5. Listings of all of the subroutines used in NETTRA-E3 can be seen in the appendix.

### 3.2.2 Example of NETTRA-E3

As in the example for NETTRA-E2, this example also begins with the same initial network described in Figure 2.4.1 (a), (b), and (c). However, NETTRA-E3 produces a whole "tree" of 126 transformed, reduced networks (note that many of these can be and probably are identical) compared with the single "string" of three transformed, reduced networks derived by NETTRA-E2 from the same initial network. This tree is partially pictured in Figure 3.2.2.1 where the numbers represent the number of gates and connections in the corresponding network. (Each number is determined by multiplying the number of gates by 1000 and adding the number of connections. For example, 22087 represents a network of 22 gates and 87 connections.)

It can be seen that seven reduced networks are obtained from the initial network alone. And from each of these "sons" of the initial network, two or more additional networks are produced, etc. It requires 117 seconds of computation time to completely generate this tree of solutions (i.e., transformed, reduced networks).

Only one "path" of solutions is completely represented in Figure 3.2.2.1. It terminates with a network consisting of ten gates and 34 connections. This is the smallest size of any of the 126 transformed networks (although it is not the only network of such size in the tree). The printout obtained from NETTRA-E3 for this network is shown in Figure



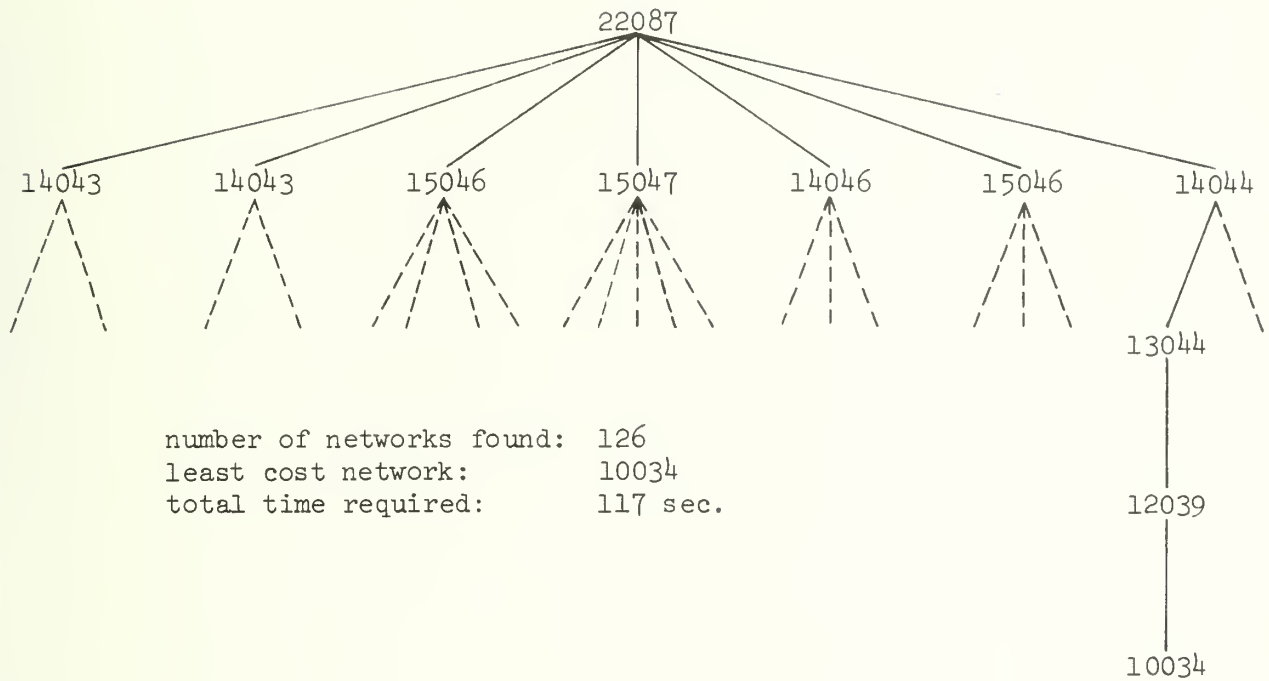


Figure 3.2.2.1 Solution tree found by NETTRA-E3.

NETWORK NUMBER 32 DERIVED BY PROCCE.

THE PARENT OF THIS NETWORK IS NUMBER 31

TRUTH TABLE

```

X1 = 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
X2 = 0 0 0 0 0 0 0 0 0 1 1 1 1 1 1 1 1 0 0 0 0 0 0 0 0 1 1 1 1 1 1 1
X3 = 0 0 0 0 1 1 1 1 0 0 0 0 1 1 1 1 0 0 0 0 1 1 1 1 0 0 0 0 1 1 1 1
X4 = 0 0 1 1 0 0 1 1 0 0 1 1 0 0 1 1 0 0 1 1 0 0 1 1 0 0 1 1 0 0 1 1
X5 = 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1
 1 = 1 0 1 0 1 0 0 0 1 0 1 0 1 0 1 0 0 1 0 1 1 1 0 0 1 1 1 0 1 0 0 0
 2 = 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
 3 = 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
 4 = 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
 5 = 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
 6 = 1 0 0 0 1 0 0 0 1 0 0 0 1 0 0 0 0 0 0 0 1 0 0 0 1 0 0 0 1 0 0 0
 7 = 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
 8 = 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
 9 = 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
10 = 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 0 0 0 0 0 0 0 0 0 0 0 0 0
11 = 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
12 = 0 0 0 0 0 0 1 0 0 0 0 0 0 0 0 0 0 1 0 1 0 0 0 1 0 0 0 0 0 0 0 0
13 = 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
14 = 0 1 0 1 0 0 0 0 0 0 0 0 0 0 0 0 0 1 1 1 1 0 0 0 0 0 0 0 0 0 0 0
15 = 1 1 0 0 1 1 0 0 0 0 0 0 0 0 0 0 0 1 1 0 0 1 1 0 0 0 0 0 0 0 0 0
16 = 1 1 0 0 0 0 0 0 1 1 0 0 0 0 0 0 1 1 0 0 0 0 0 0 1 1 0 0 0 0 0 0
17 = 1 0 1 0 0 0 0 0 1 0 1 0 0 0 0 0 0 0 0 0 0 0 0 0 1 0 1 0 0 0 0 0
18 = 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
19 = 0 0 0 0 0 0 0 1 0 0 0 1 0 1 0 1 0 0 0 0 0 0 1 1 0 0 0 1 0 1 1 1

```

(a) Truth table for transformed, reduced network.

GATE ..	LEVEL	FED BY
1	/ 1/	10 12 19
2	/ 1/	
3	/ 1/	
4	/ 1/	
5	/ 1/	
6	/ 3/	X4 X5 14
7	/ 1/	
8	/ 1/	
9	/ 1/	
10	/ 2/	X1 13
11	/ 1/	
12	/ 2/	X2 X5 6 17
13	/ 5/	X1 X5
14	/ 4/	X2 X3 13
15	/ 3/	X2 X4
16	/ 3/	X3 X4
17	/ 3/	X3 X5 14
18	/ 1/	
19	/ 2/	6 13 14 15 16 17

THIS NETWORK HAS A COST OF: 10030

(b) Configuration of final network.

3.2.2.2. A similar printout is made for each of the 126 networks found by NETTRA-E3.

This particular network is assigned the name "Network #32". It is derived from "Network #31" which is the network corresponding to the network of cost 12039 appearing in Figure 3.2.2.1. The truth table for Network #32 appears in Figure 3.2.2.2 (a), and the network's configuration description is in Figure 3.2.2.2 (b).

Thus, for the particular initial network in Figure 2.4.1 (c), NETTRA-E3 is able to obtain a better solution (i.e., a solution of less cost) than that obtained by NETTRA-E2 (see Figure 3.1.2.1).

#### 4. MAJOR FUNCTIONS OF SUBROUTINES

Thirteen different subroutines are used in the programs described in this manual: ALPATH, CALS1, CONECT, FORC, MAIN, MINI2, ORDRQ2, OUTPUT, POT, PROCCE, RCEC, RPLCF, and SUBNET. Complete program listings of these subroutines (a couple of subroutines have more than one listing since they appear in slightly different forms in different programs) can be found in the appendix.

The main functions of these thirteen subroutines are as follows:

ALPATH: This subroutine is used only in NETTRA-E3. It controls the calling of subroutine PROCCE for the various networks produced in the multi-path mode which are stored in a stack. Instead of the sequence of solutions (networks) of decreasing cost produced by NETTRA-E2, this subroutine causes the production of a "tree" of solutions.

CALS1: Subroutine RCEC calls subroutine CALS1 in several places. When CALS1 is entered, the following sets of functions are given: (1) set S which is a set of to-be-replaced input functions of the gate under considerations, (2) set S2 which contains candidate functions for the replacement of functions in set S. CALS1, based on these two sets, calculates a subset S1 of S consisting of the functions that may be replaced by functions in set S2. Since set S2 may contain mutually

incompatible functions, it is possible that set S2 cannot actually replace the calculated set S1. In such a case, CALS1 will be reentered from subroutine RPLCF with a new restricted S2.

CONNECT: This subroutine has one argument which is to specify a function in the possible output table. When this subroutine is called from RCEC, it actually constructs (make necessary connections) this function and connect it to the gate under consideration.

FORC: This subroutine is called by RCEC before other error-compensation procedures are applied. It removes redundant connections among input connections of the gate under consideration.

MAIN: This subroutine repeatedly reads in groups of input data which include information about the given networks, such as the number of external variables, the availability of the complements of variables as input variables, the number of output functions, the number of NOR gates, the list of connections, and the truth table of the output functions (see Section 5 for details). Using this information, MAIN constructs the incidence matrix, INC\$MX, for the network. INC\$MX is a two-dimensional array whose arguments represent gates or external variables. An array element  $\text{INC\$MX}(\text{GI}, \text{GJ}) \geq 1$  indicates a connection from GI to GJ; an array element  $\text{INC\$MX}(\text{GI}, \text{GJ}) \leq 0$  indicates the absence of a connection from GI to GJ. Next, subroutine SUBNET is called to calculate the level of each gate and to make lists of predecessors and successors (i.e., which gates precede which and which gates succeed which). MAIN then prints out

the truth table and the constructed incidence matrix of the original network by calling the subroutine OUTPUT. Finally the desired transduction procedure is applied to the network by calling the subroutine(s) realizing that procedure. The transformed, reduced network is stored in INC\$MX, replacing the original network. Then MAIN prints the results of the transduction procedure, i.e., the new incidence matrix and the new truth table.

MINI2: Subroutine PROCCE, when initially entered, calls MINI2 to eliminate quickly some easily removable gates in the given network. MINI2 is a subroutine which realizes a pruning procedure (i.e., it transforms a network strictly by removing connections), and it is described in some detail in [7]. Calling the entry point FORMGO simply forms an ordering of the gates of the network and stores it in the array GORDER for future use. A call to INITGS, another entry point of MINI2, initializes the CSPF vectors of the gates of a network in preparation for calling RCEC to compensate for errors in that network.

ORDRQ2: This subroutine arranges the input connections of the selected gate according to an ordering based on the number of components covering 0-error-components. Two arrays are used to store two groups of input connections from gates. Group 1 contains functions which have at least one primary 1-error-component, whereas group 2 contains functions which have 1-error-components other than primary ones. Both arrays are sorted according to the number of 1-error-components in each function.

OUTPUT: This subroutine may be entered at five different points by a call



to either OUTPUT, PAGE, LINE, TRUTH or CKT.

OUTPUT assigns mnemonic names to external variables and gates for the purpose of achieving a readable printout.

PAGE ejects one page on the printer.

LINE skips a specified number of lines on the printout sheet. The number is specified by the argument in the call (e.g., "CALL LINE(5)" skips 5 lines).

TRUTH prints out the truth table of the network currently stored in INC\$MX.

CKT prints out the network itself.

POT: This is the subroutine which constructs the potential output table. The procedure and the representation of the table are discussed in great detail in Section 2.2.2.

PROCCE: This subroutine directs the execution of the error-compensation procedure. It does not itself perform the detailed logic of the procedure, but it controls the sequence of calling the subroutines which do (primarily, RCEC, POT, and MINI2).

RCEC: This procedure realizes the essential part of the entire error-compensation procedure. It has two returns: RETURN1 will be executed if no error-compensation can be performed, whereas RETURN2 will be executed if some error-compensations are performed. In the later case, PROCCE will compare the outputs of the new network with the specified outputs. Depending on the result of comparison, RCEC will be reentered unless a

network realizing the given functions has been obtained.

RPLCF: This subroutine is called by RCEC immediately after it calls CALS1.

From set S1 calculated by subroutine CALS1 and set S2, it calculates a subset S3 of S2 which is necessary for replacing set S1. Because of the incompatibility of functions in S2, this subset may not actually exist. In this case, some incompatible functions in set S2 will be prohibited and a new set S1 will be calculated by calling CALS1. RPLCF will then repeat the procedure from the beginning.

SUBNET: This subroutine may be entered at any of three different points by a call to either SUBNET, UNNECE, or PVALUE.

SUBNET generates detailed information on the network configuration stored in INC\$MX: (1) it calculates the level of each gate in the network. Level 1 is assigned to gates having no output connections (thus all gates which have been removed from the network will be assigned level 1). (2) It lists all immediate successors and immediate predecessors for each gate. (3) It calculates the successor matrix which is stored in a two-dimensional array, SUC\$MX. The value of SUC\$MX(GI, GJ) indicates the existence or non-existence of a path from gate (or external variable) GI to gate GJ.

UNNECE disconnects certain types of obviously unnecessary connections in the network and updates the above information (discussed in (1), (2), and (3)). The connections removed from the given network are those existing in no paths from the external

variables to the output gates.

PVALUE calculates the actual truth table for the entire network stored in INC\$MX.

## 5. INPUT DATA SETUP

In order to fully understand the description of the setup of the input data cards, certain preliminary explanations are necessary.

The purpose of network transductions is to reduce the cost of a network which realizes a certain function (or functions) or to alter the network in such a way as to allow another transduction to eventually accomplish such a reduction. This cost,  $C$ , is formally defined by the weighted sum of the number of gates,  $R$ , and the number of connections<sup>†</sup>,  $I$ , of a particular network, i.e.,

$$C = A \times R + B \times I$$

where weights  $A$  and  $B$  are arbitrary non-negative numbers.

Suppose the original network which is to be transformed produces  $m$  output functions of  $n$  variables. Let  $x_\ell$ ,  $\ell = 1, \dots, n$ , be the external variables and  $f_h$ ,  $h = 1, \dots, m$ , be the output functions. Before a transformation can be performed on a network by a program, a description of that network must be input to the program. In the case when all of the output functions are completely specified (i.e., no "don't cares"), specifying only the interconnection pattern of the network is sufficient. But if one or more of the output functions is not completely specified, then the user must also provide the truth table (truth tables for all output functions are condensed into a single table) of the problem. Providing the truth table to the program consists of two steps,

---

† A "connection" refers to either a connection from an external variable or an interconnection between two gates.

namely the specification of external variables, and the specification of output functions.

The method of specifying the output functions depends directly upon the method chosen to specify the external variables. External variables may be specified in either of two ways, (a) an implicit specification of external variables, or (b) an explicit specification of external variables.

(a) In the case of implicit specification of external variables, the user specifies the number  $n$  of external variables along with a parameter which indicates whether or not the uncomplemented variables are available. Reading the value  $n$  along with the parameter, the program internally generates the entries of external variables of an ordinary truth table, that is, a truth table which consists of  $2^n$  input vectors, as shown in Fig. 5.1. In this truth table, the input vectors are arranged according to the order such that an integer  $j$  expressed in a binary representation  $(x_1 \dots x_n)$  increases, where  $x_1$  is the most significant digit and  $x_n$  is the least significant digit. For example, the truth table for a function of three variables is shown in Fig. 5.2.

The implicit specification of external variables is used for logical design problems in which the output functions have relatively few don't-care terms.

The uncomplemented variables	$x_1$	$x_1^0 \dots x_1^j \dots x_1^{2^n-1}$	
	$\vdots$	$\vdots$	
The complemented variables	$x_n$	$x_n^0 \dots x_n^j \dots x_n^{2^n-1}$	These entries exist only in the case of logical design problems where the complemented variables are available as external inputs.
	$\bar{x}_1$	$\bar{x}_1^0 \dots \bar{x}_1^j \dots \bar{x}_1^{2^n-1}$	
	$\vdots$	$\vdots$	
	$\bar{x}_n$	$\bar{x}_n^0 \dots \bar{x}_n^j \dots \bar{x}_n^{2^n-1}$	
The output functions	$f_1$	$f_1^0 \dots f_1^j \dots f_1^{2^n-1}$	
	$f_2$	$\vdots$	
	$\vdots$	$\vdots$	
	$f_m$	$f_m^0 \dots f_m^j \dots f_m^{2^n-1}$	

Fig. 5.1 The truth table of output functions of  $n$  variables

$x_1$	0	0	0	0	1	1	1	1	These entries exist only in the case of logical design problems where complemented variables are available as input variables.
$x_2$	0	0	1	1	0	0	1	1	
$x_3$	0	1	0	1	0	1	0	1	
$\bar{x}_1$	1	1	1	1	0	0	0	0	
$\bar{x}_2$	1	1	0	0	1	1	0	0	
$\bar{x}_3$	1	0	1	0	1	0	1	0	
$f_1$	$f_1^0$	.	.	.	.	.	.	$f_1^7$	

Fig. 5.2 The truth table of a function of three variables.

(b) In the case of explicit specification of external variables, the user specifies the entries of external variables of the truth table using additional cards called < external-variable-card > s. The explicit specification of external variables is used in the case of logical design problems where output functions have many don't-care terms. Suppose that the output functions are defined for a subset of input vectors of the entire truth table of Fig. 5.1. Let  $\vec{x}^j$ ,  $j = j_1, j_2, \dots, j_\mu$  denote these input vectors. The user can 'condense' the truth table of Fig. 5.1 into another table shown in Fig. 5.3.

		only $\mu$ input vectors			
The uncomplemented variables	$x_1$	$x_1^{j_1}$	$x_1^{j_2}$	$\dots$	$x_1^{j_\mu}$
	$\vdots$	$\vdots$	$\vdots$	$\dots$	$\vdots$
	$x_n$	$x_n^{j_1}$	$x_n^{j_2}$	$\dots$	$x_n^{j_\mu}$
The complemented variables	$\bar{x}_1$	$\bar{x}_1^{j_1}$	$\bar{x}_1^{j_2}$	$\dots$	$\bar{x}_1^{j_\mu}$
	$\vdots$	$\vdots$	$\vdots$	$\dots$	$\vdots$
	$\bar{x}_n$	$\bar{x}_n^{j_1}$	$\bar{x}_n^{j_2}$	$\dots$	$\bar{x}_n^{j_\mu}$
$f_1$	$f_1^{j_1}$	$f_1^{j_2}$	$\dots$	$f_1^{j_\mu}$	
$\vdots$	$\vdots$	$\vdots$	$\dots$	$\vdots$	
$f_m$	$f_m^{j_1}$	$f_m^{j_2}$	$\dots$	$f_m^{j_\mu}$	

These entries exist only in the case of logical design problems where the complemented variables are available as external inputs.

Fig. 5.3 A 'condensed' truth table having only the input vectors  $\vec{x}^j$ ,  $j = j_1, \dots, j_\mu$ , for which the output functions are defined.



Using < external-variable-card > s, the user can set up internally the table of Fig. 5.3 in place of Fig. 5.1.

### 5.1 Input Data Card Format

For each separate problem, a set of input data cards must be submitted which consists of the following<sup>†</sup>:

- |       |                                   |                                |
|-------|-----------------------------------|--------------------------------|
| (i)   | < heading-card >                  |                                |
| (ii)  | < problem-parameter-card >        |                                |
| (iii) | < external-variable-card > s      | } omitted for<br>certain cases |
| (iv)  | < output-function-card > s        |                                |
| (v)   | < connection-description-card > s |                                |

Both (i) and (ii) will always consist of only a single card, but (iii), (iv), and (v) may each consist of several cards. Furthermore, types (iii) and (iv) are omitted if all output functions are completely specified, and (iii) need only be prepared in the case of the explicit specification of external variables for the truth table. Following is a description of the formats for each type of input card, (i), (ii), (iii), (iv) and (v):

#### (i) < Heading-card >

This is the first card of the input deck for a problem. This card may contain any alphanumeric information, in columns 1~80, which may be used for the identification of the problem, but none of the information on this card will be used in the actual computation.

This information will be printed on the first page of the output.

---

<sup>†</sup> The current implementations of the NETTRA programs accept only heading, problem-parameter, and connection-description cards. Eventually it is hoped that these programs will be modified to accept all of the options described in this section.

## (ii) &lt; Problem-parameter-card &gt;

This card specifies the nature of the problem the user wants to solve. There are 7 fields in which to specify the parameters with characters and numerals. These fields are as follows:

Cols. 1~4: An integer, N, which is right-justified.

This number, N, represents the number of external variables, n, of the output functions. Be sure to punch n (rather than 2n) for N in the case of both complemented and uncomplemented variables available.

Cols. 5~8: An integer, M, which is right-justified.

This number, M, is the number of output functions, m, to be realized simultaneously. Therefore, of course, M will also be the number of output gates in the network.

Cols. 9~12: An integer, R, which is right-justified.

This number, R, specifies the number of gates which are included in the network. For various reasons, the user may wish to input networks in which one or more of the gates are "isolated" (i.e., are not connected to any other gates). This is permissible as long as these "isolated" gates are also included in the total number of gates, R.

Cols. 13~16: An integer, A, which is right-justified.

The number A is the value of the non-negative weight for the number of gates in the cost function. (See Table 5.1.1, 'Typical combinations of values A and B for different network reduction problems'.)

Cols. 17~20: An integer, B , which is right-justified.

The number B is the value of the non-negative weight for the number of connections in the cost function. (See Table 5.1.1.)

Col. 24: A blank 'b'<sup>†</sup>, or one of the characters, 'C', 'X', 'Y', 'U' or 'V'.

The 'b' or 'C' parameter represents an implicit specification of both the external variables and an implicit specification of the output functions (in this case, the output functions will be calculated from the connection pattern of the network). The 'X' or 'Y' parameter indicates an implicit specification of external variables only. The 'U' or 'V' parameter indicates an explicit specification of external variables. (See summary of these symbols in Table 5.1.2)

The 'b' or 'X' parameter specifies that only uncomplemented external variables are available for the network. The 'C' or 'Y' parameter specifies that both uncomplemented and complemented variables are available for the network. If the user specifies the 'b', 'X', 'C', or 'Y' parameter, the program sets up the truth table by generating a set of  $2^n$  input vectors  $(x_1^j, \dots, x_n^j)$ , for  $j = 0, \dots, 2^n - 1$ , in the case of a 'b' or 'X' parameter, or  $(x_1^j, \dots, x_n^j, \bar{x}_1^j, \dots, \bar{x}_n^j)$  for  $j = 0, \dots, 2^n - 1$ , in the case of a 'C' or 'Y' parameter.

The 'b' or 'C' parameters should be used for problems in which the output functions contain no don't-care terms. For such problems, the preparation of the < external-variable-card > s and the < output-function-card > s can be dispensed with since the program can calculate completely all output functions using only a description of the

---

<sup>†</sup> A 'b' stands for a blank (i.e., no character punched).

Network Reduction Problem	Values of A and B
reducing the number of gates only.	A = 1 and B = 0
reducing the number of gates primarily, then reducing the number of connections secondarily. <sup>†</sup>	A = 100 and B = 1
reducing the number of connections only.	A = 0 and B = 1
reducing the number of connections primarily, then reducing the number of gates secondarily.	A = 1 and B = 100
reducing the sum of the number of gates and the number of connections.	A = B = 1

Table 5.1.1 Typical combinations of values A and B for different network reduction problems.

<sup>†</sup> Most of the programs in the NETTRA system are oriented toward this reduction problem, so the user will probably find this combination of A and B the most useful.

uncomplemented variables only available	both complemented and uncomplemented variables available	
'b'	'c'	} implicit specification of external variables and output functions
'X'	'Y'	} implicit specification of external variables
'U'	'V'	} explicit specification of external variables

Table 5.1.2 Possible symbols for column 24 of < problem-parameter-card >.

connection pattern of the network (provided by the <connection-description-card>s).

Similarly, the 'X' or 'Y' parameter implies the use of a complete truth table (i.e.,  $2^n$  input vectors for  $n$  external variables) inside the program. Since from this information the program can easily generate the truth table entries for the external variables, as just mentioned, the < external-variable-card > s are unnecessary. The  $m$  < output-function-card > s, however, must still be prepared.

The 'U' parameter specifies that only uncomplemented external variables are available for the network. The 'V' parameter specifies that both uncomplemented and complemented variables are available for the network. In either case, the 'U' or the 'V' parameter, the user must prepare  $n$  < external-variable-card > s and  $m$  < output-function-card > s. The program sets up the truth table by reading these < external-variable-card > s and < output-function-card > s.

Cols. 25~28: An integer, NEPMAX, which is right-justified.

This parameter is omitted for all NETTRA programs except those involving "error-compensation" routines. In the cases where NEPMAX is required, a further discussion of this parameter can be found elsewhere in the manual. The abbreviation NEPMAX is a mnemonic for "maximum number of error positions", and the default is  $NEPMAX = 2^{(n-1)}$ , where  $n$  is the number of external variables.

(iii) < External-variable-card > s

In combination with the 'U' or 'V' parameter in column 24 of the < problem-parameter-card >, the  $n$  < external-variable-card > s specify the entries of external variables of the truth table of

Fig. 5.3. Each card contains the binary representation of external variable  $x_\ell$ , i.e.,  $(x_\ell^{j1}, x_\ell^{j2}, \dots, x_\ell^{j\mu})$ , starting from column 1 of the card. The maximum number of bits in a binary representation is limited to 32. (This means the maximum number of input vectors is 32.) If the actual number of bits is less than 32, then a termination symbol '/' (slash) is put on the right of the right-most bit of the binary representation on the first < external-variable-card >. The remaining columns after the termination symbol '/' in the first card, as well as the same columns in the following cards, may contain any alphanumeric information which may be used for identification. This information will not be printed on the output pages.

In the case of the 'V' parameter, the program generates the binary representations corresponding to complemented variables by taking negations of the entries of the < external-variable-card > s. Therefore the user must not provide < external-variable-card > s representing the complemented variables,  $\bar{x}_\ell$ .

If one of the parameters 'b', 'C', 'X', or 'Y' appears in column 24 of the < problem-parameter-card >, the user does not prepare < external-variable-card > s.

(iv) < Output-function-card > s

The m < output-function-card > s specify the set of m output functions to be realized simultaneously. Each card contains the binary representation of one output function  $f_h$ , starting from column 1 of the card. A symbol '\*' is used to denote don't-care terms, if any. The maximum number of bits in a binary representation is limited to 32.



The actual number of bits must be  $2^n$  in the case of an implicit specification of external variables, or must be the same as defined by the < external-variable-card > s in the case of an explicit specification of external variables. The remaining columns, up to column 72 (inclusive), may contain any alphanumeric information which may be used for identification. This information will not be printed on the output pages.

If either the 'b' or 'C' parameter appears in column 24 of the < problem-parameter-card >, the < output-function-card > s must be omitted.

(v). < Connection-description-card > s

In the present version of the program, 9 cards (some of which may be just blank cards) are required.<sup>†</sup> Each of these cards is divided into 16 fields of 5 columns each (i.e., columns 1~5, 6~10, 11~15, ..., 71~75, 76~80). Beginning with the first field of the first card, continuing through the succeeding fields of that card and through the fields of as many additional cards as necessary (up to a maximum of 9, total), the expressions (explained in the next paragraph)  $C_1, C_2, C_3, \dots$ , are punched right-justified in their respective fields.

Each gate of the network is labeled uniquely by assigning it one of the integers 1, 2, ..., R, such that the output gates receive

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† For many uses, the user will probably find that these 9 cards far exceed his needs, and may thus be inconvenient. In such a case, the number of required cards may be easily adjusted by making the obvious changes in two statements (A READ statement and a DO statement) following the comment card "C\*\*\*\* READ IN NETWORK INFORMATION AND SET UP INC\$MX \*\*\*\*\*" in subroutine MAIN.

the labels 1, 2, ..., m. The names  $X_1, X_2, \dots, X_n$  are assigned to the external variables  $x_1, x_2, \dots, x_n$  (and the names  $Y_1, Y_2, \dots, Y_n$  to the complemented external variables  $\bar{x}_1, \bar{x}_2, \dots, \bar{x}_n$ , if appropriate).<sup>†</sup> Now, for each connection of the network (i.e., including both the connections from external variables to gates and connections from gates to other gates), a 4 character expression,  $C_i$ , is formed, to represent that connection as follows: to represent a connection from gate GI to gate GJ, the numeric label GI is inserted into the first two character positions of  $C_i$  and the numeric label GJ is inserted into the second two positions (e.g., the  $C_i$  for a connection from gate 9 to gate 5 would be "0905"); to represent a connection from external variable XI to gate GJ, the alphanumeric label XI is inserted into the first two character positions of  $C_i$  and the numeric label GJ into the second two positions (e.g., the  $C_i$  for a connection from external variable  $x_3$  to gate 10 would be "X310").

Every connection of the network must be represented by a  $C_i$ , although there are no restrictions on the order in which the connections (i.e.,  $C_i$ 's) are punched onto the input cards.

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<sup>†</sup> At the time of writing, the programs have not yet been changed to recognize this new labeling system. The old labels are simply:

1, 2, ..., n, for external variables  $x_1, x_2, \dots, x_n$  (and  $n+1, n+2, \dots, 2n$  for the complemented variables  $\bar{x}_1, \bar{x}_2, \dots, \bar{x}_n$ , if they are permitted in the problem);  $n+1, n+2, \dots, n+m$  for the m output gates of the network ( $2n+1, 2n+2, \dots, 2n+m$  if complemented variables are included); and finally  $n+m+1, n+m+2, \dots, n+R$  ( $2n+m+1, 2n+m+2, \dots, 2n+R$ ) for the non-output gates of the network.

These five groups of cards, (i), (ii), (iii), (iv) and (v) in the correct order constitute the necessary description for a single problem. In order to solve several problems during the same computer run, the descriptions for the desired problems are just arranged serially. See Fig. 5.1.1 for an example of the sequencing of all cards for the execution of a NETTRA program using typical JCL statements for the IBM 360/75.

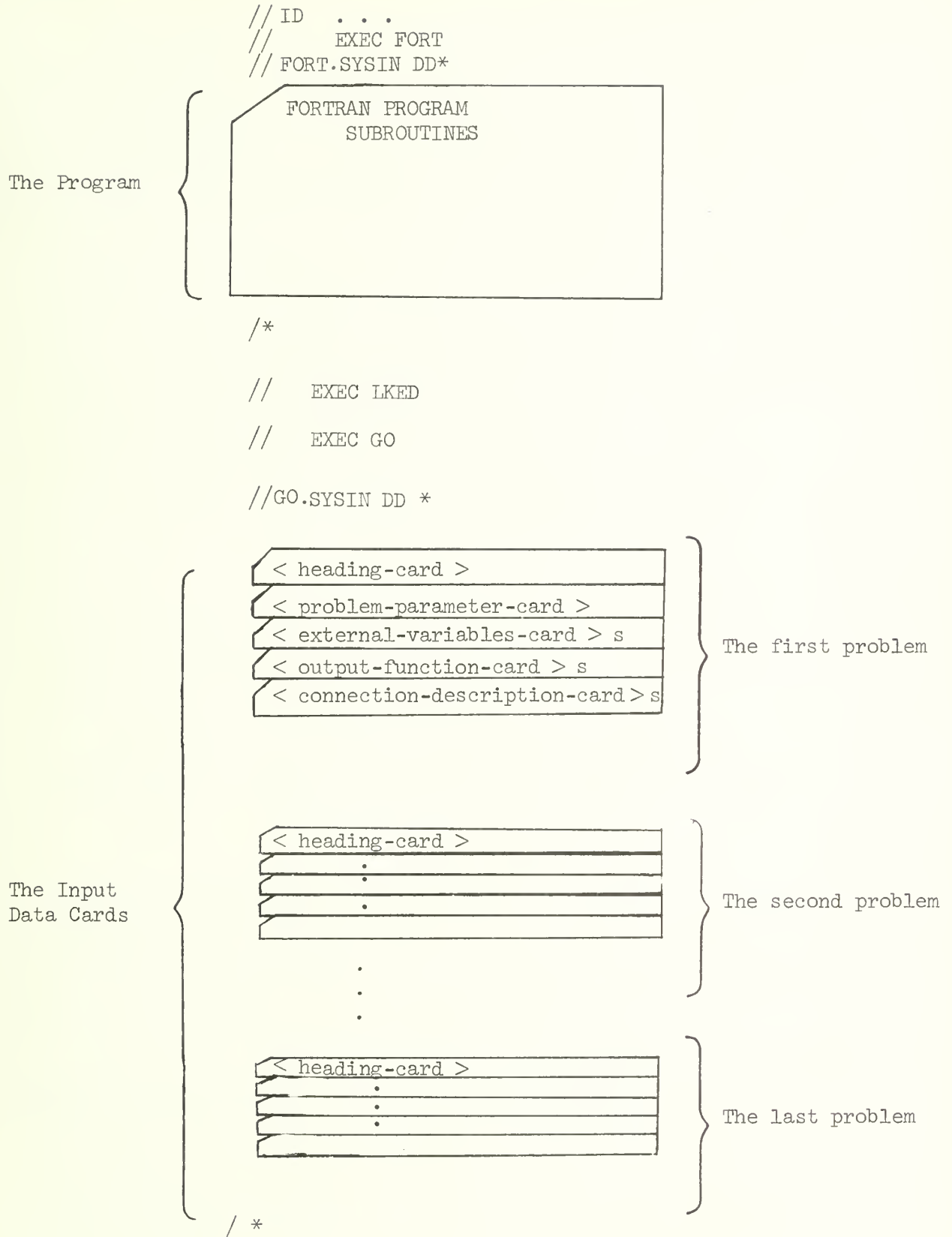


Fig. 5.1.1 Input card sequence for the execution of a typical NETTRA program.

## 5.2 Restrictions on Problem Size

In order to fit the programs into a finite amount of space, some restrictions on the size of an acceptable problem are required:

1. The number  $t$  of input vectors in the truth table is 32 or less.
2. The number  $n$  of external variables.

Because of  $t \leq 32$ ,  $n$  is 5 or less in the case of completely specified functions. In the case of incompletely specified functions, however, any  $n \leq 20$  is acceptable if only uncomplemented variables are available, or  $n \leq 10$  if both uncomplemented and complemented variables are available, provided that the truth table is defined by the `< external-variable-card > s`.

3. The number  $R$  of gates.

The number of gates,  $R$ , may not exceed  $40-n$  in the case of only uncomplemented variables available (a 'b', 'X', or 'U' parameter). In the case of both uncomplemented and complemented variables available (a 'C', 'Y' or 'V' parameter), the limit is lowered to  $40-2n$ .

All of these limitations are essentially imposed by the array sizes in the programs as presently written. To loosen the restrictions is mainly a task of increasing array dimensions appropriately.

## 5.3 Examples of Input Data Setup

The following examples will illustrate, for the general program in the NETTRA system, various possible input data card setups complying with the directions given in Section 5.1.

Example 1: A two output network of four variables shown in Fig. 5.3.1. Assume the two output functions are  $f_1 = CCEF^{\dagger}$  and  $f_2 = 3BBB$  and only uncomplemented variables are available. Furthermore, assume it is desired to reduce the number of gates primarily and the number of connections secondarily (see Table 5.1.1).<sup>††</sup>

From this description, the < problem-parameter-card > must contain the following values:

Cols. 1~4	4, the number of external variables
Cols. 5~8	2, the number of output functions
Cols. 9~12	8, the number of gates in the original network
Cols. 13~16	100, the value of A
Cols. 17~20	1, the value of B
Cols. 24	'b', uncomplemented variables only available and implicit specification of both the external variables and the output functions
Cols. 25~28	'b', since the NEPMAX parameter is unrelated to the program to be used

Fig. 5.3.2 shows the setup of data cards used to specify the network in Fig. 5.3.1 as input for the program. Notice that in forming the  $C_i$ , the four uncomplemented variables are represented by the labels  $X_1, X_2, X_3, X_4$ ; the two output gates by the numbers 1, 2; and the remaining gates, by the numbers 3, 4, 5, 6, 7, 8. This manner of labeling is

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<sup>†</sup> For convenience, the output functions are expressed in hexadecimal notation. When the numbers in this notation are expanded into binary, they represent the output vectors as they appear (i.e., in the same left-to-right order) in the complete truth table described earlier and pictured in Fig. 5.1.

<sup>††</sup> This assumption is implicit in most of the transduction procedures and their implementations which comprise the NETTRA system.

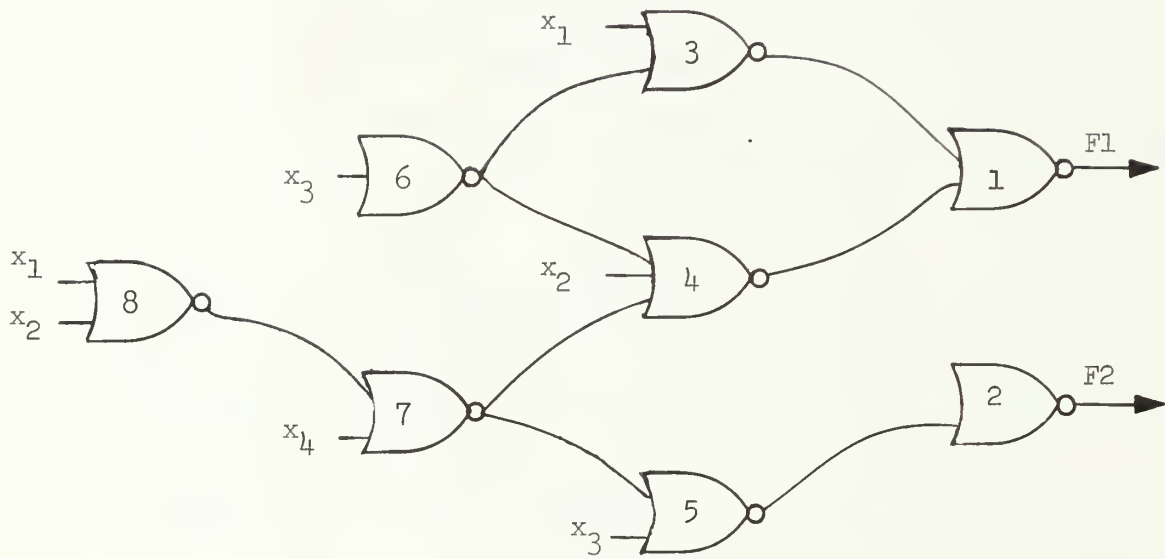


Fig. 5.3.1 Network to be transformed in Examples 1 and 2.



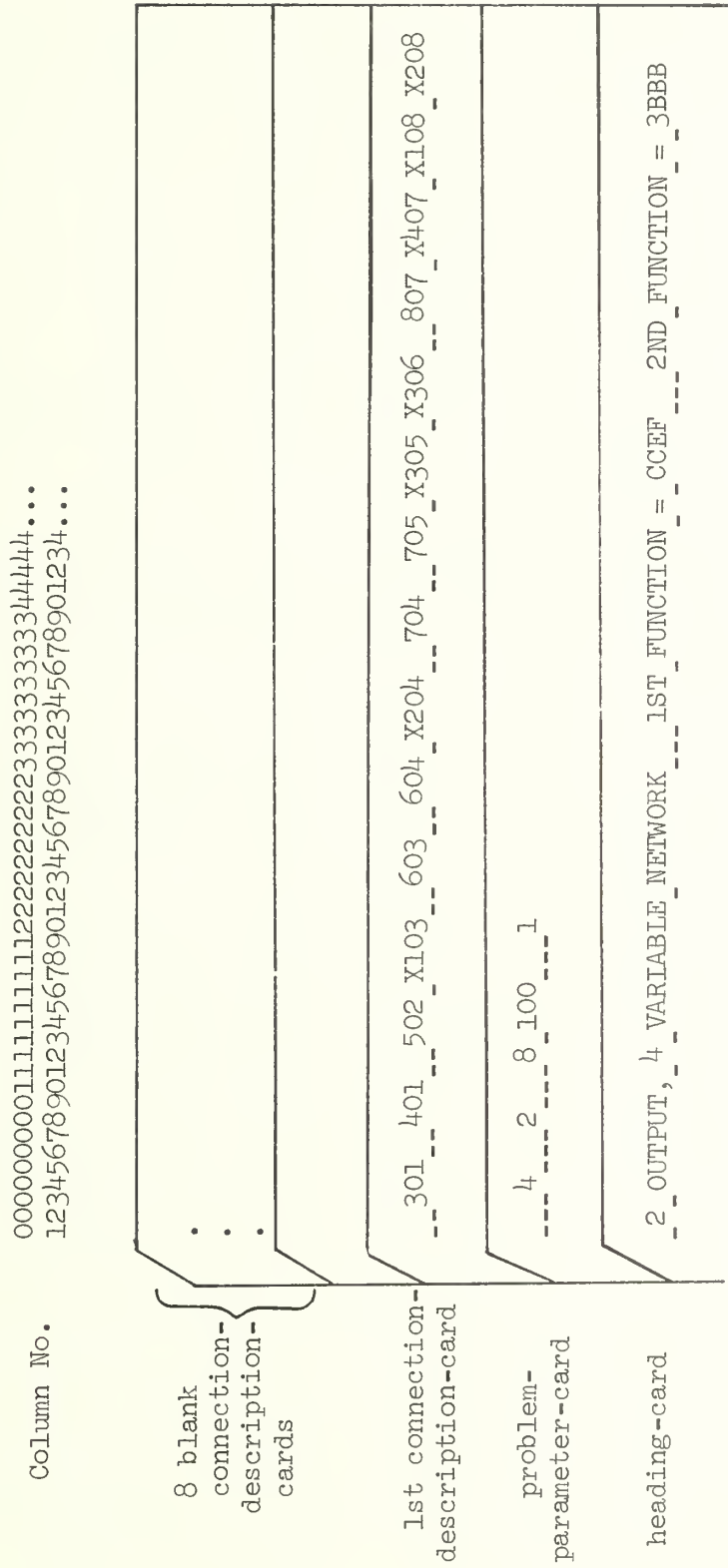


Fig. 5.3.2 Possible setup of data cards to specify the problem given in Example 1.

strictly required by the instructions for preparing the < connection-description-card > s (see Section 5.1).

The heading card in Fig. 5.3.2 will simply be read by the program and printed character for character onto the output page as an identification of the particular problem. Below that, the number of variables, number of functions, and the cost coefficients, A and B, will be printed (all with appropriate labels). Also, immediately following will be a statement of what types of external variables are permitted (i.e., either just uncomplemented variables or both complemented and uncomplemented) along with their generic names:

X - for uncomplemented variables	}	if external variables were implicitly specified
Y - for complemented variables		
<u>or</u>		
U - for uncomplemented variables	}	if external variables were explicitly specified
V - for complemented variables		

For example, if both X and Y appear as generic names (as would occur in the case of an implicit specification of external variables with both complemented and uncomplemented variables available) then the external variable names which appear on subsequent output pages will be X1, X2, ..., Xn and Y1, Y2, ..., Yn. Or, if both U and V appear as generic names (as would occur in the case of an explicit specification of external variables with both complemented and uncomplemented variables available) the external variable names which appear in the output will be U1, U2, ..., Un (for the uncomplemented variables) and V1, V2, ..., Vn (for the complemented variables). It should be noted, however, that the letters U and V, as used as replacements for X and Y (respectively) in the

naming of external variables (e.g. U1, V1 instead of X1, Y1), appear strictly on the output pages of the program - they are not used internally in the program and they must not appear in the variable names punched on the < connection-description-card > s by the user. They are intended only as an aid to the user so that, at a glance at the transformed network in the output, he can easily distinguish whether the external variables were implicitly or explicitly specified for that particular problem.

Following the statement of whether only uncomplemented or both complemented and uncomplemented external variables are employed, the user will find next on the output page the cost of the original network which was input to the program. This is the cost which was defined in the beginning of Section 5.

The cost will be followed by a truth table (generally in the same form as Fig. 5.1) showing the outputs (0 or 1) of all of the gates in the network for every external variable input combination (i.e., combinations of 0's and 1's) of interest.

Finally, below the truth table will be printed a description of the network submitted as input. This is for documentation purposes, and it is also much more readable than the network description which appeared on the < connection-description-card > s. In this description, each gate is listed along with the names of the gates and external variables which feed it. Also, to assist the user in sketching the network from its description, the level of each gate in the network is included (gates which do not feed other gates are assigned to level 1, all other gates are assigned level numbers such that each gate is in a level one

higher than the highest level gate directly fed by it).

All of the information just described will be printed before the execution of the transduction actually begins. This will be followed, beginning at the top of a new output page, by the network(s) actually obtained as a result of the computation. First the complete truth table of the transformed network will be printed, followed by a network connection description of the form just described above. Finally, the cost of the new network will be calculated and printed.

In this example, it was assumed that there were no "don't-cares" in the output functions implicitly specified by the input, thus no `< external-variable-card > s` or `< output-function-card > s` were included. In the next example, however, `< output-function-card > s` will be required in order to specify some of the components of the output functions as "don't-cares".

Example 2: The two output network of four variables shown in Fig. 5.3.1. This is the same network used in Example 1, but this time the output functions are not assumed to be completely specified. Let  $f_1 = '11001**01*10*111'$  and  $f_2 = '0**110111*111011'$  be the required functions. Also, suppose that both complemented and uncomplemented variables are desired to be available during the transduction. Again the problem is to reduce the number of gates primarily and the number of connections secondarily.

For this problem, the following values must appear on the `< problem-parameter-card >`:

Cols. 1~4	4, the number of external variables
Cols. 5~8	2, the number of output functions

Cols.	9~12	8, the number of gates in the original network
Cols.	13~16	100, the value of A
Cols.	17~20	1, the value of B
Col.	24	Y, indicative of an implicit specification of external variables and the availability of both complemented and uncomplemented variables

Fig. 5.3.3 shows the setup of the data cards corresponding to this problem. Notice the differences and similarities to the data cards shown in Fig. 5.3.2. The < problem-parameter-card > differs only in column 24. The < external-variable-card > s are missing in both Fig. 5.3.2 and Fig. 5.3.3 since the external variables are implicitly specified for both problems. The < output- function-card > s, however, appear in Fig. 5.3.3 but not in 5.3.2 since they are necessary to specify "don't-care" components which do not occur in the completely specified output functions of Example 1. In both cases, though, the < connection-description-card > s are identical since the original networks are identical.

By allowing "don't-care" terms in the output functions, and by allowing the use of both complemented and uncomplemented variables<sup>†</sup> (even though the original network employed only uncomplemented variables), the restrictions during the transduction process are loosened (compared to what they were for Example 1), perhaps permitting a network of

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<sup>†</sup> In the case of NETTRA-PG1, -P1, and -P2, it is useless to specify Y rather than X in column 24 for this example. Since the original network uses only uncomplemented variables, to these programs which perform "pruning" procedures (i.e., procedures which are incapable of adding new connections) the availability of complemented variable is not meaningful.

Column No.:

00000000011  
12345678901...

...78  
...90

<div> 8 blank connection- description-cards </div>	.
	.
	.
1st connection- description-card	_301_ 401_ 502 X103_ 603_ 604 X204_ 704_ 705 X305 X306_ 807 X407 X108 X208
2nd output- function-card	0**110111*111011_-----REQUIRED OUTPUT FOR GATE 2_
1st output function-card	11001**01*10*111_-----REQUIRED OUTPUT FOR GATE 1_
problem- parameter-card	4_ 2_ 8 100 1 Y_
heading-card	_2 OUTPUT, 4 VARIABLE NETWORK, F1=11001**01*10*111, F2=0**110111*111011

Fig. 5.3.3 Possible setup of data cards to specify the problem given in Example 2.

less cost to be obtained.

Notice that the first < output-function-card > corresponds to the output of gate 1 and the second < output-function-card > corresponds to the output of gate 2. This must hold true for every problem in which < output-function-card > s are included; the gates labeled 1, 2, ..., m must correspond to the output functions specified on < output-function-card > s 1, 2, ..., m, respectively.

Of course, the printed output of the program will be in the same format described in Example 1.

Example 3: The three output network of six variables shown in Fig. 5.3.4. The outputs are again assumed to be incompletely specified. In fact, only the following 11 input combinations are specified out of a possible 64 ( $= 2^6$ ):

$x_1$	0 0 0 0 0 0 0 0 0 0 1
$x_2$	0 0 0 0 0 0 0 1 1 1 0
$x_3$	0 0 0 0 0 0 0 0 0 1 1
$x_4$	0 0 0 0 1 1 1 0 0 0 1
$x_5$	0 0 1 1 0 0 1 0 1 1 0
$x_6$	0 1 0 1 0 1 0 1 1 0 0
<hr/>	
$F_1$	0 0 1 1 0 0 * 0 0 0 0
$F_2$	1 1 * 1 1 1 0 1 1 0 *
$F_3$	1 1 0 0 0 0 0 1 0 0 0



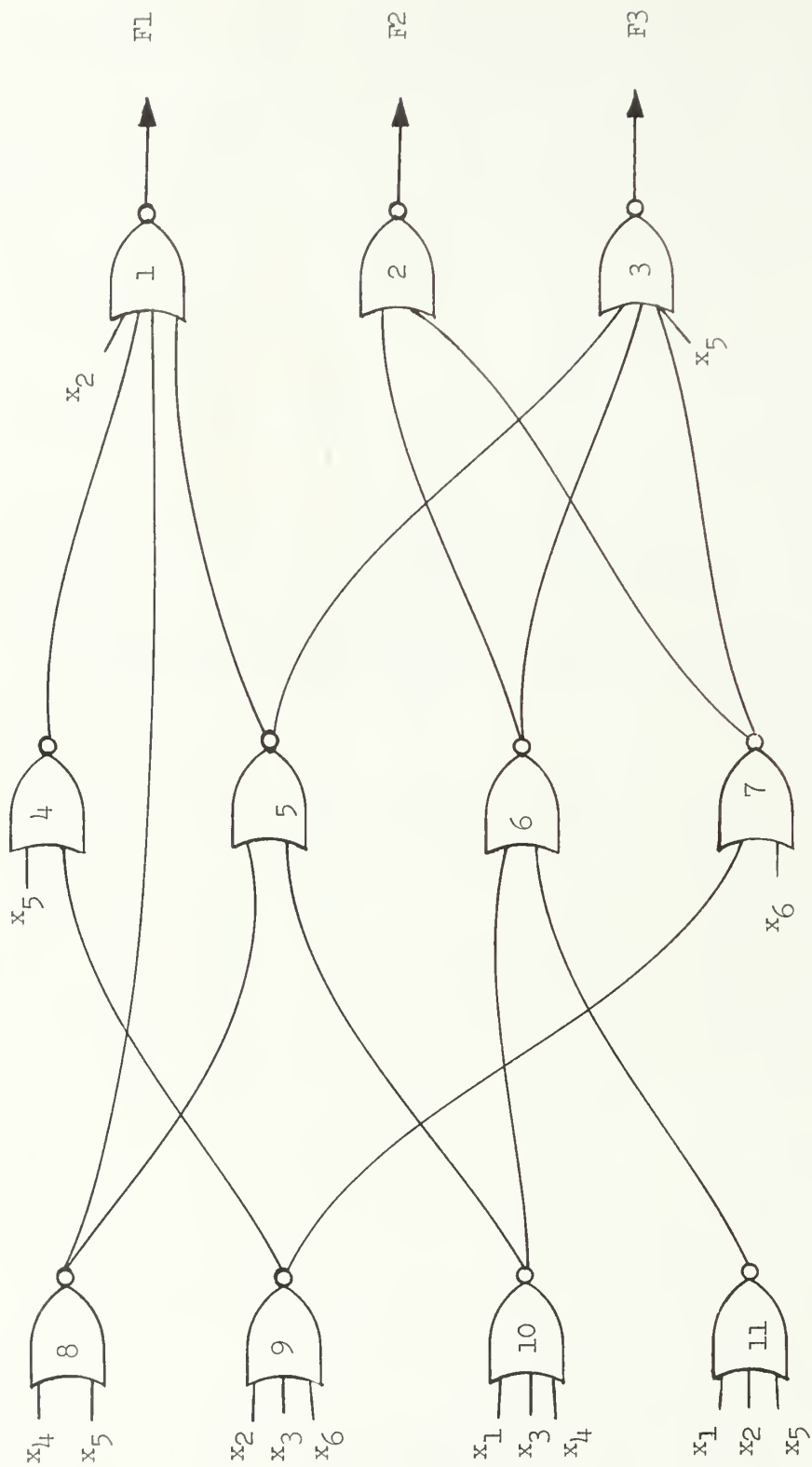


Fig. 5.3.4 Network to be transformed in Example 3.

Additionally, only uncomplemented variables are assumed to be available, and the problem is to reduce the number of gates primarily and the number of connections secondarily.

For this example, the following parameters appear on the < problem-parameter-card >:

Cols.	1~4	6, the number of external variables
Cols.	5~8	3, the number of output functions
Cols.	9~12	11, the number of gates in the original network
Cols.	13~16	100, the value of A
Cols.	17~20	1, the value of B
Col.	24	U, indicative of an explicit specification of external variables and the availability of only uncomplemented variables

Fig. 5.3.5 shows a possible setup of the data cards corresponding to this example. Notice that in this example, the <external-variable-card> s are included, whereas in the two previous examples they were omitted. Although this problem is not too realistic (none of the 3 functions is actually a 6-variable function), it demonstrates the input data preparation to be used in cases where many external variables are present and a high percentage of "don't care" terms exist.

Again, the printed output from the program will follow the same format described in Example 1.



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APPENDIX:  
Program Listings

Following are the listings of the FORTRAN programs NETTRA-E1, NETTRA-E2, and NETTRA-E3. These programs realize, respectively, the transduction procedures discussed in Section 2, Section 3.1, and Section 3.2.

Since NETTRA-E2 and NETTRA-E3 only differ in a few subroutines from NETTRA-E1, only these subroutines are listed for NETTRA-E2 and -E3: MAIN (differs slightly in each of the three programs), ALPATH (used only in NETTRA-E3, and PROCCE (slightly different for NETTRA-E3).

Explanations of variables used in the programs can be found in the listings themselves.

\*\*\*\*\*

```

      PPPP      RRRR      OOO      GGG      RRRR      A      M      M
      P  P      R  R      O  O      G  G      R  R      A  A      MM  MM
      P  P      R  R      O  O      G  G      R  R      A  A      MM  MM
      PPPP      RRRR      OOO      GGG      RRRR      AAAAA      M  M  M
      P          R  R      O  O      G  G      R  R      A  A      M  M
      P          R  R      OOO      GGG      R  R      A  A      M  M
  
```

```

      N  N      EEEEE      TTTT      TTTT      RRRR      A      EEEEE      1
      NN  N      F          T          T          R  R      A  A      E      11
      N  N  N      F          T          T          R  R      A  A      E      1
      N  NN      FFE          T          T          RRRR      AAAAA      XXXXX      EEE      1
      N  N      E          T          T          R  R      A  A      E      1
      N  N      EEEEE      T          T          R  R      A  A      EEEEE      111
  
```

\*\*\*\*\*

```

SURROUTINE MAIN                                     E1 00010
EDITION BBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBB E1 00020
NOTE: ALL COMMON VARIABLES MIGHT NOT BE USED IN THIS PROGRAM. E1 00030
COMMON VARIABLES:                                   E1 00040
  $GT: POINTS TO A 'COLUMN' OF POTAB. FOR EACH 'ROW' THE ENTRY E1 00050
        IN THIS COL. TELLS GATE WHERE FN. IS REALIZED.          E1 00060
  $LTH: POINTS TO A 'COLUMN' OF POTAB. FOR EACH 'ROW' THE ENTRY E1 00070
        IN THIS COL. TELLS HOW MANY CONNECTIONS MUST BE ADDED. E1 00080
  $NDE: POINTS TO A 'COLUMN' OF POTAB. FOR EACH 'ROW' THE ENTRY E1 00090
        IN THIS COL. TELLS THE NUMBER OF 1-ERRORS CREATED IF THIS E1 00100
        ROW IS USED.                                             E1 00110
  $PW: POINTS TO A 'COLUMN' OF POTAB. FOR EACH 'ROW' THE ENTRY E1 00120
        IN THIS COLUMN TELLS THE PREFERENCE WEIGHT.             E1 00130
  A: WEIGHT FOR NO. OF GATES IN COMPUTING COST FUNCTION.        E1 00140
  B: WEIGHT FOR NO. OF CONNECTIONS IN COMPUTING COST FUNCTION. E1 00150
  COST: COST OF NETWORK - A MEASURE OF NETWORK SIZE.            E1 00160
  ESS1S: RECORDS NO. OF ESSENTIAL 1'S IN EVERY INPUT TO CURRENT GCOE1 00170
        (POSITIONS IN ESS1S CORRES. TO GATES NOT FEEDING GCO ARE E1 00180
        IGNORED).                                              E1 00190
  F$UB1: POINTS TO LAST ELEMENT IN F$1.                         E1 00200
  F$1: LISTS (CONSECUTIVELY) POSITIONS OF DESIRABLE 1'S (FOR E1 00210
        COVERING) IN A CONNECTIBLE FUNCTION.                    E1 00220
  GI: LABEL OF A PARTICULAR GATE.                                E1 00230
  GLEVEL: GLEVEL(GI) TELLS WHICH LEVEL OF THE NETWORK GI IS IN. E1 00240
  G$SMALL: STORES INTERMEDIATE AND FINAL CALCULATED CSPF'S.     E1 00250
  HLIST: HLIST(I,J) GIVES NAME OF I-TH GATE (OR EX. VAR.) IN NET- E1 00260
        WORK LEVEL J.                                           E1 00270
  IDX0: LIST OF 0-COORDINATES IN CSPFE OF THE GATE UNDER E1 00280
        CONSIDERATION.                                          E1 00290
  IDX0E: LIST OF 0-ERROR-COORDINATES IN CSPFE OF THE GATE UNDER E1 00300
        CONSIDERATION.                                          E1 00310
  IDX1: LIST OF 1-COORDINATES IN CSPFE OF THE GATE UNDER E1 00320
        CONSIDERATION.                                          E1 00330
  IDX1E: LIST OF 1-ERROR-COORDINATES IN CSPFE OF THE GATE UNDER E1 00340
        CONSIDERATION.                                          E1 00350
  IDX1E: LIST OF 1-ERROR-COORDINATES IN CSPFE OF THE GATE UNDER E1 00360
  
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C      CONSIDERATION. E1 00370
C      TFLAG: SAME AS FYEFLG IN SUBROUTINE PROCII. E1 00380
C      INC$MX: INC$MX(GI,GJ)>0 MEANS THERE EXISTS A CONVECTION FROM GATE E1 00390
C      (OR EX. VAR.) GJ TO GATE GI. INC$MX(GI,GJ)=0 IF NOT. E1 00400
C      INPTTV: LISTS FOR EACH CORRESPONDING ENTRY OF F$1, HOW MANY INPUTS E1 00410
C      HAVE A '1' IN THE POSITION INDICATED BY F$1. E1 00420
C      IPATH: IPATH(GI)=1 MEANS GATE GI IS ON A PATH FROM A CERTAIN GATE E1 00430
C      TO AN OUTPUT GATE. OTHERWISE IPATH(GI) = 0. E1 00440
C      IPRED: IPRED(I,GJ) GIVES THE NAME OF THE I-TH GATE OR EX. VAR. IVE1 00450
C      A LIST OF GATES AND EX. VAR. FEEDING GJ. E1 00460
C      ISUCC: ISUCC(I,GJ) GIVES THE NAME OF THE I-TH GATE FED BY GJ. E1 00470
C      JFLAG: SAME AS JAYFLG IN SUBROUTINE PROCII. E1 00480
C      KEYA: A FLAG INDICATING IF ANY ERROR COMPENSATION HAS BEEN E1 00490
C      PERFORMED. E1 00500
C      KEYR: A FLAG INDICATING IF ANY PRIMARY O-ERROR-COORDINATES HAS E1 00510
C      BEEN COMPENSATED. E1 00520
C      KFLAG: SAME AS KEIFLG IN PROCII. E1 00530
C      LEVM: NUMBER OF LEVELS IN THE NETWORK (NOTE EX. VAR. ARE ALSO E1 00540
C      ASSIGNED LEVELS JUST LIKE GATES). E1 00550
C      LGLIST: LGLIST(J) TELLS NO. OF GATES AND EX. VAR. IN LEVEL J OF E1 00560
C      NETWORK. E1 00570
C      LIP: NUMBER OF PREDECESSORS FOR THE GATE UNDER CONSIDERATION. E1 00580
C      LIPPED: LIPPED(GI) TELLS NO. OF IMMEDIATE PREDECESSORS OF GATE GI. E1 00590
C      LISTC: ORDERED LIST OF CONNECTIBLE INPUTS TO GCC. ORDERED BY E1 00600
C      DECREASING NO. OF O'S IN GCC COVERED. E1 00610
C      LISTL: ORDERED LIST OF GATES AND EX. VAR. WHICH ORIGINALLY FED E1 00620
C      GCC AND WHICH HAVE NOT YET BEEN DISCONNECTED. ORDERED BY E1 00630
C      DECREASING NO. OF ESSENTIAL 1'S. E1 00640
C      LISUCC: LISUCC(GI) TELLS NO. OF IMMEDIATE SUCCESSORS OF GATE (OR E1 00650
C      EX. VAR.) GI. E1 00660
C      LMTS2: UPPER LIMIT OF THE NUMBER OF ELEMENTS IN SET S2. E1 00670
C      LPOTAB: FOR GATE GI, LPOTAB(GI) POINTS TO LAST ROW OF POTAB E1 00680
C      CONCERNING GI. E1 00690
C      M: NUMBER OF NETWORK OUTPUT GATES. E1 00700
C      N: NUMBER OF EXTERNAL VARIABLES (OR INPUT FNC.) AVAILABLE. E1 00710
C      NEPMAX: FOR ERROR COMPENSATION PROGRAMS. IF MORE THAN NEPMAX E1 00720
C      ERROR POSITIONS OCCUR WHEN A PARTICULAR GATE IS REMOVED, E1 00730
C      PROGRAM SKIPS ATTEMPT TO COMPENSATE FOR THAT GATE'S E1 00740
C      REMOVAL. VALUE CAN BE SPECIFIED BY USER, OTHERWISE EQUAL E1 00750
C      TO ONE HALF OF N2 BY DEFAULT. E1 00760
C      NM: SUM OF N PLUS M E1 00770
C      NM1: SUM OF NM PLUS 1. E1 00780
C      N2: PRODUCT OF N AND N2. E1 00790
C      NOS: NUMBER OF ELEMENTS IN SET S. E1 00800
C      NOS1: NUMBER OF ELEMENTS IN SET S1. E1 00810
C      NOS1SV: NUMBER OF ELEMENTS IN SET S1 BEFORE ENTERING SUBROUTINE E1 00820
C      RPLCF. E1 00830
C      NOS2: NUMBER OF ELEMENTS IN SET S2. E1 00840
C      NOT1: NUMBER OF ELEMENTS IN SET T1. E1 00850
C      NOT1SV: NUMBER OF ELEMENTS IN SET T1 BEFORE ENTERING SUBROUTINE E1 00860
C      RPLCF. E1 00870
C      NDO: NUMBER OF ACTIVE ELEMENTS IN ARRAY IDX0. E1 00880
C      NDOF: NUMBER OF ACTIVE ELEMENTS IN ARRAY IDX0E. E1 00890
C      NDI: NUMBER OF ACTIVE ELEMENTS IN ARRAY IDX1. E1 00900
C      NDIE: NUMBER OF ACTIVE ELEMENTS IN ARRAY IDX1E. E1 00910
C      NR: SUM OF N PLUS R. E1 00920
C      NRN2: PRODUCT OF NR AND N2. E1 00930
C      NRPLC: NRPLC(I) STORES THE NUMBER OF ELEMENTS IN RPLC(I,*) E1 00940
C      FOR I=1,2. E1 00950
C      N1: SUM OF N PLUS 1. E1 00960
C      N2: NUMBER OF DIFFERENT INPUT COMBINATIONS TO BE CONSIDERED E1 00970

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C          (USUALLY 2 TO THE POWER N).
C ORIGIN: ORIGIN(GI)=1 MEANS GI ORIGINALLY CONNECTED TO GCO.
C          ORIGIN(GI)=0 MEANS GI DID NOT FEED GCO ORIGINALLY.
C          P$: P$(1,-) CONSECUTIVELY LISTS OUTPUTS OF EVERY EX. VAR. AND
C              EVERY GATE (FOR EVERY INPUT COMBINATION): P$(1,1),...,
C              P$(1,N2) FOR FIRST EX VAR; P$(1,N2+1),...,P$(1,2*N2) FOR
C              SECOND EX VAR; ... ; P$(1,N*N2+1),..., P$(1,N*N2+N2) FOR
C              FIRST GATE; ETC. P$(2,-) IS USED AS WORK SPACE FOR
C              CALCULATIONS ASSOCIATED WITH P$(1,-).
C          PCO: FOR ERROR COMPENSATION PROCEDURES. PCO IS THE GATE
C              REMOVED FROM ORIGINAL NETWORK TO OBTAIN CURRENT ALTERED
C              NETWORK.
C          POINTA: NOT USED.
C          POINTC: POINTS TO LAST ELEMENT IN LISTC.
C          POINTL: POINTS TO LAST ELEMENT IN LISTL.
C          POINTR: POINTS TO LAST ELEMENT IN RNEC1 (IN SUBROUTINE SUBST1).
C          POTAB: POTENTIAL OUTPUT TABLE. HOLDS INFORMATION ABOUT ALL
C              COMBINATIONS OF CONNECTIONS TO FORM NEW (AND HOPEFULLY
C              USEFUL) FUNCTIONS.
C          PPOTAB: FOR GATE GI, PPOTAB(GI) POINTS TO FIRST OF A SEQUENCE OF
C              ROWS OF POTAB CONCERNING GI.
C          R: NUMBER OF GATES IN THE NETWORK (EXCLUDES EX VAR, ALSO
C              NOTE SOME OF R GATES MAY BE ISOLATED).
C          RPLC: RPLC(1,*) STORES THE SELECTED GATE'S IP GATES WHICH HAVE
C              ERROR-COORDINATES OF WEIGHT 2 OR ABOVE.
C              RPLC(2,*) STORES THE SELECTED GATE'S IP GATES WHICH HAVE
C              AT LEAST ONE ERROR-COORDINATE OF WEIGHT 1.
C          RSCONN: LIST OF CONNECTIONS ADDED TO A NETWORK (IN CODED FORM).
C          RTCONN: LIST OF CONNECTIONS REMOVED FROM A NETWORK (CODED FORM).
C              S: NO. OF CONNECTIONS ADDED TO A NETWORK. POINTS TO LAST
C              ENTRY IN RSCONN.
C          SETS: SET S CONSISTING OF INPUTS OF THE GATE UNDER CONSIDERATION
C              WHICH ARE TO BE REPLACED IF POSSIBLE.
C          SETS1: SET S1 CONSISTING OF ELEMENTS OF SET S WHICH CAN BE
C              REPLACED BY ELEMENTS IN SET S2.
C          SETS2: SET S2 CONSISTING OF FUNCTIONS WHICH ARE CANDIDATES FOR
C              REPLACING ELEMENTS IN SET S.
C          SETT1: SET T1 CONSISTING OF ESSENTIAL ONES COVERED BY ELEMENTS IN
C              SET S1.
C          STS: STARTING ELEMENT OF SET S.
C          SUC$MX: SUC$MX(GI,GJ)>0 MEANS GATE GJ IS A SUCCESSOR OF GATE GI.
C              SUC$MX(GI,GJ)=0 IF NOT.
C          SUMP: SUM OF ALL ACTIVE INPUTS OF THE GATE UNDER CONSIDERATION.
C          SUMS2: SUM OF ALL ACTIVE ELEMENTS OF SET S2.
C          T: NUMBER OF CONNECTIONS REMOVED FROM A NETWORK. POINTS TO
C              LAST ENTRY IN RTCONN.
C          TIME: USED TO STORE AMOUNT OF ELAPSED COMPUTATION TIME.
C          JNAME: MNEMONIC NAMES FOR EXTERNAL VARIABLES AND GATES.
C          VF$UR1: POINTS TO LAST ELEMENT IN VF$1.
C          VF$1: SIMILAR TO F$1, EXCEPT THIS LISTS JUST COMPONENT POSITIONS
C              (OF J'S IN CSPF VECTOR OF GCO) COVERED ONLY BY REMAINING
C              ORIGINALLY CONNECTED INPUTS TO GCO.
C
C          IMPLICIT INTEGER*4(A-T,V-Z,$), REAL(U)
C          COMMON NEPMAX
C          COMMON      N          , M          , A          , B
C          1          , R          , N2         , N1         , NR
C          2          , NM         , KFLAG      , JFLAG      , COST
C          3          , LEVM       , NRN2       , NM1        , NN2

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COMMON  ISJCC(40,40) , LISUCC(40) , IPRED(40,40) , LIPRED(40) E1 01590
1      , INC$MX(40,40) , SUP$MX(40,40) , P$(2,1280) , UNAME(40) E1 01600
2      , GLEVEL(40) , LGLIST(40) , HLIST(40,40) , TIME E1 01610
COMMON  T , RTCONN(100) , S , RSCONN(100) E1 01620
COMMON  IFLAG , POINTA , ESSIS(40) , F$1(32) E1 01630
1      , F$URL , INPTCV(32) , LISTC(40) , POINTC E1 01640
2      , LISTL(40) , POINTL , ORIGIN(40) , IPATH(40) E1 01650
3      , POINTP , VF$1(32) , VF$UB1 , GSMALL(40,32) E1 01660
COMMON  POTAB(200,42) , PPOTAB(40) , LPOTAB(40) , NRPLC(2) E1 01670
1      , RPLC(2,40) , IDXC(32) , IDXC1(32) , IDXC1(32) E1 01680
2      , IDXC1(32) , SUMP(32) , SETT1(32) , NOT1 E1 01690
3      , SETS1(40) , NOS1 , SETS(40) , NOS E1 01700
4      , STS , SUMS2(32) , SETS2(200) , NOS2 E1 01710
5      , LIF , NOOF , KEYA , KEYB E1 01720
6      , NOC , NOI , NOIE , $GT E1 01730
7      , $LTH , $PW , $NOE , GI E1 01740
COMMON  NOTISV , NOS1SV , LMTS2 E1 01750
DIMENSION INTLIS(144),UGATE(40),UHEAD(20) E1 01760
DATA KOUNT5 /0/, URLANK/' ' E1 01770
990 READ(5,1000,END=500) UHEAD, N, M, R, A, B, UC, NEPMAX E1 01780
NEPMAX IS THE MAXIMUM ALLOWABLE NUMBER OF ERROR POSITIONS E1 01790
1000 FORMAT(20A4/5I4,A4,I4) E1 01800
KEYXC=0 E1 01810
IF(UC.NE.URLANK) KEYXC=1 E1 01820
CALL PAGE E1 01830
CALL LINE(10) E1 01840
KOUNT5=KOUNT5+1 E1 01850
PRINT 2, KOUNT5 E1 01860
2 FORMAT(20X,'*** OPTIMAL NDR NETWORK ***',50X,'PROBLEM NO.= ',I4 ) E1 01870
CALL LINE(4) E1 01880
PRINT 1005, UHEAD E1 01890
1005 FORMAT(25X,20A4) E1 01900
CALL LINE(4) E1 01910
PRINT 10, N,M,A,B E1 01920
10 FORMAT(30X,'NUMBER OF VARIABLES =',I4 // E1 01930
1 30X,'NUMBER OF FUNCTIONS =',I4 // E1 01940
2 30X,'COST COEFFICIENT A =',I4// E1 01950
3 47X, 'B =',I4) E1 01960
CALL LINE(1) E1 01970
IF(KEYXC.NE.0) GO TO 25 E1 01980
PRINT 21 E1 01990
21 FORMAT(1H0,29X,'--- UNCOMPLEMENTED VARIABLES X ---') E1 02000
GO TO 30 E1 02010
25 CONTINUE E1 02020
PRINT 28 E1 02030
28 FORMAT(1H0,29X,'--- BOTH COMPLEMENTED AND UNCOMPLEMENTED VARIABLES E1 02040
1 X, Y ---') E1 02050
30 CONTINUE E1 02060
CALL LINE(5) E1 02070
***** SET UP EXTERNAL VARIABLES ***** E1 02080
N2=2**N E1 02090
IF(NEPMAX.EQ.0) NEPMAX = N2/2 E1 02100
H=N*N2 E1 02110
J=N2 E1 02120
L= 1 E1 02130
I=0 E1 02140
DO 1011 II=1,N E1 02150
J=J/2 E1 02160
L=L*2 E1 02170
SN= 1 E1 02180
DO 1010 LL=1,L E1 02190

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SN=-SN	E1 02200
V=(1+SN)/2	E1 02210
DO 1009 JJ=1,J	E1 02220
I=I+1	E1 02230
P\$(1,I)=V	E1 02240
IF(KEYXC.NE.0)P\$(1,I+H)=1-V	E1 02250
1009 CONTINUE	E1 02260
1010 CONTINUE	E1 02270
1011 CONTINUE	E1 02280
IF(KEYXC.NE.0) N=N+N	E1 02290
N1=N+1	E1 02300
NM=N+M	E1 02310
NM1=N+1	E1 02320
NV2=V*N2+1	E1 02330
NR=N+R	E1 02340
NRN2=NR*N2	E1 02350
CALL OUTPUT(INC\$MX,KEYXC)	E1 02360
***** READ IN NETWORK INFORMATION AND SET UP INC\$MX *****	E1 02370
READ 1001, CNTLIS	E1 02380
1001 FORMAT(16I5)	E1 02390
DO 1115 GI=1,NR	E1 02400
DO 1115 GJ=1,NR	E1 02410
1115 INC\$MX(GI,GJ)=0	E1 02420
DO 1120 I=1,144	E1 02430
ITEM=CNTLIS(I)	E1 02440
IF(ITEM.EQ.0) GO TO 1119	E1 02450
GI=ITEM/100	E1 02460
GJ=ITEM-100*GI	E1 02470
INC\$MX(GI,GJ)=1	E1 02480
GO TO 1120	E1 02490
1119 COST=A*R+B*(I-1)	E1 02500
GO TO 1130	E1 02510
1120 CONTINUE	E1 02520
1130 CONTINUE	E1 02530
CALL SUBNET	E1 02540
CALL PVALUE	E1 02550
CALL LINE(4)	E1 02560
PRINT 1140, COST	E1 02570
1140 FORMAT(20X,' ORIGINAL NETWORK COST=', 15)	E1 02580
CALL LINE(4)	E1 02590
CALL TRUTH(P\$,1)	E1 02600
CALL LINE(4)	E1 02610
CALL CKT(INC\$MX,GLEVEL)	E1 02620
	E1 02630
***** ENTRY REDUNDANCY CHECK *****	E1 02640
S = 0	E1 02650
T = 0	E1 02660
CALL UNNECE	E1 02670
GATES = M	E1 02680
C = 0	E1 02690
DO 4 GI = 1,NR	E1 02700
C = C + LISUCC(GI)	E1 02710
IF(GI.LE.NM)GOTO4	E1 02720
IF(LISUCC(GI).GT.0)GATES=GATES+1	E1 02730
4 CONTINUE	E1 02740
OLDGST = A*GATES + B*(C)	E1 02750
T=0	E1 02760
S=0	E1 02770
INITIALIZE TIMER TO 10 MINUTES	E1 02780
CALL STIMEZ(60000)	E1 02790
TIME = KTIMEZ(0)	E1 02800

C****	PROCEDURE PROCCE	E1	02810
	CALL PROCCE(WDEKED)	E1	02820
C	CALL FOR ELAPSED TIME	E1	02830
	TIME = KTIMEZ(0) - TIME	E1	02840
	CALL LINE(4)	E1	02850
	PRINT 3915	E1	02860
3916	FORMAT(20X,'TIME ELAPSED =',I8,' CENTISECONDS')	E1	02870
3915	FORMAT(20X,'NETWORK DERIVED BY PROCCE')	E1	02880
	PRINT 3916,TIME	E1	02890
	CALL LINE(4)	E1	02900
	CALL TRUTH(P\$,1)	E1	02910
	CALL LINE(4)	E1	02920
	CALL CKT(INC\$MX,GLEVEL)	E1	02930
	GATES = M	E1	02940
	C = 0	E1	02950
	DO 36 GT = 1,NR	E1	02960
	C = C + LISUCC(GI)	E1	02970
	IF(GI.LE.NM) GO TO 36	E1	02980
	IF(LISUCC(GI).GT.0) GATES = GATES + 1	E1	02990
36	CONTINUE	E1	03000
	NEWCST = A*GATES + B*C	E1	03010
	IF(NEWCST.LT.OLDCST)GO TO 37	E1	03020
	PRINT 105	E1	03030
105	FORMAT(1H ,10X,'NO REDUNDANCY FOUND.')	E1	03040
	GO TO 990	E1	03050
37	CALL LINE(3)	E1	03060
	PRINT 320,NEWCST	E1	03070
320	FORMAT(9X,'* A NETWORK DERIVED BY PROCCE'/9X,' COST=',I5,'.')	E1	03080
	GO TO 990	E1	03090
500	STOP	E1	03100
	END	E1	03110
	SUBROUTINE CALS1	E1	03120
C		E1	03130
C****	THIS SUBROUTINE CALCULATES A MAXIMUM SUBSET, S1, OF S	E1	03140
C	WHICH CONTAINS FUNCTIONS REPLACEABLE BY S2. THE NUMBER	E1	03150
C	OF INITIAL ELEMENTS IN S1 (NOS1) IS SET BY CALLING PROGRAM.	E1	03160
C	THE NUMBER OF ESSENTIAL ONES COVERED BY INITIAL S1 IS ALSO	E1	03170
C	SET BY CALLING PROGRAM (NCT1 AND SETT1(*))	E1	03180
C		E1	03190
C		E1	03200
C	DEFINITIONS OF 'COMMON' VARIABLES CAN BE FOUND IN MAIN PROGRAM.	E1	03210
C		E1	03220
C	IMPLICIT INTEGER*4(A-T,V-Z,\$), REAL(U)	E1	03230
C	COMMON NEPMAX	E1	03240
	COMMON N , M , A , B	E1	03250
1	, R , N2 , N1 , NR	E1	03260
2	, NM , KFLAG , JFLAG , COST	E1	03270
3	, LEVM , NRN2 , NM1 , NN2	E1	03280
	COMMON ISUCC(40,40) , LISUCC(40) , IPRED(40,40) , LIPRED(40)	E1	03290
1	, INC\$MX(40,40) , SUC\$MX(40,40) , P\$(2,1280) , UNAME(40)	E1	03300
2	, GLEVEL(40) , LGLIST(40) , HLIST(40,40) , TIME	E1	03310
	COMMON T , RTCONN(100) , S , RSCONN(100)	E1	03320
	COMMON IFLAG , POINTA , ESSIS(40) , F\$1(32)	E1	03330
1	, F\$UB1 , INPTCV(32) , LISTC(40) , POINTC	E1	03340
2	, LISTL(40) , POINTL , ORIGIN(40) , IPATH(40)	E1	03350
3	, POTNTR , VF\$1(32) , VF\$UB1 , GSMALL(40,32)	E1	03360
	COMMON POTAB(200,42) , PPOTAB(40) , LPOTAB(40) , NRPLC(2)	E1	03370
1	, RPLC(2,40) , IDX0(32) , IDX0E(32) , IDX1(32)	E1	03380
2	, IDX1E(32) , SUMP(32) , SETT1(32) , NOT1	E1	03390



3	,SETS1(40)	,NOS1	,SETS(40)	,NOS	E1 03400
4	,STS	,SUMS2(32)	,SETS2(200)	,NOS2	E1 03410
5	,LIP	,NOOE	,KEYA	,KEYB	E1 03420
6	,NOO	,NOI	,NOIE	,\$GT	E1 03430
7	,\$LTH	,\$PW	,\$NOE	,GI	E1 03440
	COMMON	NOT1SV	,NOS1SV	,LMTS2	E1 03450
	NOT1=NOT1SV				E1 03460
	NOS1=NOS1SV				E1 03470
C	STS IS THE STARTING ELEMENT OF SETS				E1 03480
	DO 7800 NO = STS,NOS				E1 03490
	GP = SETS(NO)				E1 03500
C****	GP > 1000 : ALREADY REMOVED ****				E1 03510
	IF(GP.GT.1000)GO TO 7800				E1 03520
	BSGP = (GP-1)*N2				E1 03530
C****	CALCULATE ESSENTIAL ONES IN GP AND CHECK WHETHER SETS2 COVERS				E1 03540
C****	THEM OR NOT *****				E1 03550
	NOTO = NOT1				E1 03560
	DO 7100 NOZ=1,NOO				E1 03570
	TH = IDX0(NOZ)				E1 03580
	IF(SUMP(TH).NE.1.OR.P\$(1,BSGP+TH).NE.1)GO TO 7100				E1 03590
	IF(SUMS2(TH).LE.0)GO TO 7800				E1 03600
	NOTO = NOTO + 1				E1 03610
	SETT1(NOTO) = TH				E1 03620
7100	CONTINUE				E1 03630
	NOT1 = NOTO				E1 03640
	NOS1 = NOS1 + 1				E1 03650
	SETS1(NOS1) = GP				E1 03660
C****	UPDATE SUMP *****				E1 03670
	DO 7300 TH=1,N2				E1 03680
	SUMP(TH) = SUMP(TH) - P\$(1,BSGP+TH)				E1 03690
7300	CONTINUE				E1 03700
C****	UPDATE SET S (MAKE GP INACTIVE) *****				E1 03710
	SETS(NO) = 1000 + GP				E1 03720
7800	CONTINUE				E1 03730
	RETURN				E1 03740
	END				E1 03750
	SUBROUTINE CONNECT(PTR)				E1 03760
C					E1 03770
C****	THIS SUBROUTINE CONNECTS THE FUNCTION IN POTAB SPECIFIED BY PTR				E1 03780
C	TO GATE GI AND MAKES ALL OTHER NECESSARY CONNECTIONS FOR				E1 03790
C	REALIZING THIS FUNCTION.				E1 03800
C					E1 03810
C					E1 03820
C	DEFINITIONS OF 'COMMON' VARIABLES CAN BE FOUND IN MAIN PROGRAM.				E1 03830
C					E1 03840
	IMPLICIT INTEGER*4(A-T,V-Z,\$), REAL(U)				E1 03850
	COMMON NEPMAX				E1 03860
	COMMON N	, M	, A	, B	E1 03870
1	, R	, N2	, N1	, NR	E1 03880
2	, NM	, KFLAG	, JFLAG	, COST	E1 03890
3	, LEVM	, NRN2	, NM1	, NN2	E1 03900
	COMMON ISUCC(40,40)	, LISUCC(40)	, IPRED(40,40)	, LIPRED(40)	E1 03910
1	, INC\$MX(40,40)	, SUC\$MX(40,40)	, P\$(2,1280)	, UNAME(40)	E1 03920
2	, GLEVEL(40)	, LGLIST(40)	, HLIST(40,40)	, TIME	E1 03930
	COMMON T	, RTCONN(100)	, S	, RSCONN(100)	E1 03940
	COMMON IFLAG	, POINTA	, ESSIS(40)	, F\$1(32)	E1 03950
1	, F\$UB1	, INPTCV(32)	, LISTC(40)	, POINTC	E1 03960
2	, LISTL(40)	, POINTL	, ORIGIN(40)	, IPATH(40)	E1 03970
3	, POINTR	, VF\$1(32)	, VF\$UB1	, G\$SMALL(40,32)	E1 03980

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COMMON POTAB(200,42),PPOTAB(40) ,LPOTAB(40) ,NRPLC(2) E1 03990
1 ,RPLC(2,40) ,IDX0(32) ,IDX0E(32) ,IDX1(32) E1 04000
2 ,IDX1E(32) ,SUMP(32) ,SETT1(32) ,NOT1 E1 04010
3 ,SETS1(40) ,NOS1 ,SETS(40) ,NOS E1 04020
4 ,STS ,SUMS2(32) ,SETS2(200) ,NOS2 E1 04030
5 ,LIP ,NOOE ,KEYA ,KEYB E1 04040
6 ,NOO ,NO1E ,NO1E ,GT E1 04050
7 ,LTH ,PW ,NOE ,GI E1 04060
COMMON NOT1SV ,NOS1SV ,LMTS2 E1 04070
C**** CONNECT THIS FUNCTION E1 04080
GP = POTAB(PTR,$GT) E1 04090
INC$MX(GP,GI) = 1 E1 04100
S = S + 1 E1 04110
RSCONN(S) = 100*GP + GI E1 04120
C**** CONNECT OTHER NECESSARY CONNECTIONS E1 04130
IF(POTAB(PTR,$LTH).EQ.0)GO TO 7200 E1 04140
LTH = POTAB(PTR,$LTH) E1 04150
DO 7100 TH=1,LTH E1 04160
GO = POTAB(PTR,$LTH+TH) E1 04170
INC$MX(GO,GP) = 1 E1 04180
S = S + 1 E1 04190
RSCONN(S) = 100*GO + GP E1 04200
7100 CONTINUE E1 04210
7200 KEYA = 1 E1 04220
RETURN E1 04230
END E1 04240

SUBROUTINE FORC(GJ) E1 04250
C**** REMOVE FIRST ORDER REDUNDANT CONNECTION ***** E1 04260
SUBROUTINE TO REMOVE FIRST ORDER REDUNDANT CONNECTION E1 04270
GJ TO GI IS THE CONNECTION TO BE CHECKED. INPUT SUM OF GATE GI E1 04280
IS STORED AT SUMP. E1 04290
DEFINITIONS OF 'COMMON' VARIABLES CAN BE FOUND IN MAIN PROGRAM. E1 04300
IMPLICIT INTEGER*4(A-T,V-Z,$), REAL(U) E1 04310
COMMON NEPMAX E1 04320
COMMON N , M , A , B E1 04330
1 , R , N2 , N1 , NR E1 04340
2 , NM , KFLAG , JFLAG , COST E1 04350
3 , LEVM , NRN2 , NM1 , NN2 E1 04360
COMMON ISUCC(40,40) , LISUCC(40) , IPRED(40,40) , LIPRED(40) E1 04370
1 , INC$MX(40,40) , SUC$MX(40,40) , P$(2,1280) , UNAME(40) E1 04380
2 , GLEVEL(40) , LGLIST(40) , HLIST(40,40) , TIME E1 04390
COMMON T , RSCONN(100) , S , RSCONN(100) E1 04400
COMMON IFLAG , POINTA , ESS1S(40) , F$1(32) E1 04410
1 , F$UB1 , INPTCV(32) , LISTC(40) , POINTC E1 04420
2 , LISTL(40) , POINTL , ORIGIN(40) , IPATH(40) E1 04430
3 , POINTR , VF$1(32) , VF$UB1 , GSMALL(40,32) E1 04440
COMMON POTAB(200,42),PPOTAB(40) ,LPOTAB(40) ,NRPLC(2) E1 04450
1 ,RPLC(2,40) ,IDX0(32) ,IDX0E(32) ,IDX1(32) E1 04460
2 ,IDX1E(32) ,SUMP(32) ,SETT1(32) ,NOT1 E1 04470
3 ,SETS1(40) ,NOS1 ,SETS(40) ,NOS E1 04480
4 ,STS ,SUMS2(32) ,SETS2(200) ,NOS2 E1 04490
5 ,LIP ,NOOE ,KEYA ,KEYB E1 04500
6 ,NOO ,NO1E ,NO1E ,GT E1 04510
7 ,LTH ,PW ,NOE ,GI E1 04520
COMMON NOT1SV ,NOS1SV ,LMTS2 E1 04530

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C		E1 04580
C****	LIST ESSENTIAL ONES OF GJ TO GI	E1 04590
	KEYAA = 0	E1 04600
	BSGJ = (GJ-1)*N2	E1 04610
	DO 310 TH=1,N2	E1 04620
	IF(SUMP(TH).NE.1.OR.P\$(1,BSGJ+TH).NE.1)GO TO 310	E1 04630
	IF(GSMALL(GI,TH).GT.-1000)GO TO 300	E1 04640
	KEYAA = 1	E1 04650
	GO TO 310	E1 04660
300	IF(GSMALL(GI,TH).LT.0) RETURN	E1 04670
310	CONTINUE	E1 04680
C****	DISCONNECT GJ TO GI, UPDATE SUMP *****	E1 04690
	INC\$MX(GJ,GI) = 0	E1 04700
	T = T + 1	E1 04710
	RTCCNN(T) = 100*GJ + GI	E1 04720
	DO 320 TH=1,N2	E1 04730
320	SUMP(TH) = SUMP(TH) - P\$(1,BSGJ+TH)	E1 04740
	IF(KEYAA.EQ.1) KEYA = 1	E1 04750
	KEYR = 1	E1 04760
	RETURN	E1 04770
	END	E1 04780
	SUBROUTINE MINIZ2(IMPROV)	E1 04790
C	EDITION AA	E1 04800
C	THE NAME ATTEMPTS TO INDICATE THAT THIS SUBROUTINE IS A MINIATURE	E1 04810
C	VERSION OF PROCEDURE II (PROCII) - ACTUALLY, THIS ROUTINE ONLY	E1 04820
C	REMOVES CONNECTIONS, NONE ARE ADDED	E1 04830
C		E1 04840
C	DEFINITIONS OF 'COMMON' VARIABLES CAN BE FOUND IN MAIN PROGRAM.	E1 04850
C		E1 04860
C	VARIABLE DEFINITIONS:	E1 04870
C	BESTSL: NAME OF A PRIORITY CANDIDATE TO DISCONNECT FROM GATE GCO.	E1 04880
C	CHOICE: NAME OF A GATE CHOSEN TO BECOME A COVER.	E1 04890
C	COMPNT: A COMPONENT OF AN INTERMEDIATE CSPF VECTOR.	E1 04900
C	EFLAG: SIGNALS WHICH ENTRY POINT USED.	E1 04910
C	FEEDGT: A GATE FEEDING GATE 'GATE'.	E1 04920
C	F\$UBO: NUMBER OF 'NECESSARY ZEROS' LISTED IN F\$O.	E1 04930
C	F\$O: LISTS (CONSECUTIVELY) POSITIONS OF NECESSARY ZEROS IN A	E1 04940
C	CONNECTABLE FUNCTION VECTOR.	E1 04950
C	GATE: NAME OF A GATE.	E1 04960
C	G\$COUNT: A COUNTER.	E1 04970
C	G\$ORDER: A SPECIAL ORDERING OF GATES AND EXTERNAL VARIABLES SUCH	E1 04980
C	THAT NO GATE SUCCEEDS A PREDECESSOR IN THE ORDERING.	E1 04990
C	MARKED: MARKED(GI)=1 FOR GI FEEDING 'GATE' INDICATES THAT GI HAS	E1 05000
C	ALREADY BEEN ASSIGNED NECESSARY ZEROS CORRESPONDING TO	E1 05010
C	'1' COMPONENTS IN THE CSPF VECTOR FOR 'GATE'.	E1 05020
C	NMINLV: NUMBER OF GATES IN A CERTAIN LEVEL OF THE NETWORK.	E1 05030
C	SELECT: NAME OF AN INPUT SELECTED AS A CANDIDATE FOR DISCONNECTION	E1 05040
C	FROM GATE 'GCO'.	E1 05050
C	T: COUNTS REMOVED CONNECTIONS.	E1 05060
C	T\$ORDER: A SPECIAL ORDERING OF GATES AND EXTERNAL VARIABLES SUCH	E1 05070
C	THAT VARIABLES COME FIRST FOLLOWED BY GATES WITH DECREASED	E1 05080
C	NUMBERS OF OUTPUTS (TIES ARE BROKEN BY G\$ORDER).	E1 05090
C	T\$POINT: POINTER TO T\$ORDER.	E1 05100
C	T1\$PRED: LIST OF GCO'S PREDECESSORS AT ONE STAGE OF COMPUTATION.	E1 05110
C	T2\$PRED: LIST OF GCO'S PREDECESSORS AT ONE STAGE OF COMPUTATION.	E1 05120
C	T1\$SUB: A POINTER TO T1\$PRED.	E1 05130
C	T2\$SUB: A POINTER TO T2\$PRED.	E1 05140
C	USED: USED(GI)=1 MEANS GI IS AN OUTPUT GATE, OR IS A COVER FOR	E1 05150
C	SOME 0-COMPONENT OF 'GATE'. (IT ALSO HAS A TEMPORARY USE	E1 05160



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C          IN BEGINNING OF PROGRAM.) E1 05170
C E1 05180
C COUNT,I,II,J,K,L,MOST,Q,TCOUNT,X,XX,Y ARE USED AS JUST TEMPORARY E1 05190
C VARIABLES. E1 05200
C HOW TO INCREASE CAPACITY OF SUBROUTINE. E1 05210
C DIMENSION: T1PRED(X),T2PRED(X),GORDER(X), E1 05220
C MARKED(X),USED(X),TORDER(X) - X EQUAL TO MAX NUMBER E1 05230
C OF GATES PLUS EXTERNALE E1 05240
C VARIABLES. E1 05250
C F$O(Y) - Y EQUAL TO: 2** (MAX ALLOWED NO. OF EX. VAR.) E1 05260
C E1 05270
C IMPLICIT INTEGER*4(A-T,V-Z,$), REAL(U) E1 05280
C COMMON NEPMAX E1 05290
C COMMON N , M , A , B E1 05300
1 , R , N2 , N1 , NR E1 05310
2 , NM , KFLAG , JFLAG , COST E1 05320
3 , LEVM , NRN2 , NM1 , NN2 E1 05330
COMMON ISUCC(40,40) , LISUCC(40) , IPRED(40,40) , LIPRED(40) E1 05340
1 , INC$MX(40,40) , SUP$MX(40,40) , P$(2,1280) , UNAME(40) E1 05350
2 , GLEVEL(40) , LGLIST(40) , HLIST(40,40) , TIME E1 05360
COMMON T , RTCONN(100) , S , RSCONN(100) E1 05370
COMMON IFLAG , POINTA , ESSIS(40) , F$1(32) E1 05380
1 , F$UB1 , INPTCV(32) , LISTC(40) , POINTC E1 05390
2 , LISTL(40) , POINTL , ORIGIN(40) , IPATH(40) E1 05400
3 , POINTR , VF$1(32) , VF$UB1 , GSMALL(40,32) E1 05410
COMMON POTAB(200,42) , PPOTAB(40) , LPOTAB(40) , NRPLC(2) E1 05420
1 , RPLC(2,40) , IDX(32) , IDXOE(32) , IDX1(32) E1 05430
2 , IDX1E(32) , SUMP(32) , SETT1(32) , NOT1 E1 05440
3 , SETS1(40) , NOS1 , SETS(40) , NOS E1 05450
4 , STS , SUMS2(32) , SETS2(200) , NOS2 E1 05460
5 , LIP , NOOE , KEYA , KEYB E1 05470
6 , NCO , NO1 , NO1E , $GT E1 05480
7 , $LTH , $PW , $NCE , GI E1 05490
COMMON NOT1SV , NOS1SV , LMTS2 E1 05500
DIMENSION T1PRED(40),T2PRED(40),GORDER(40),F$O(32),MARKED(40) E1 05510
DIMENSION USED(40),TORDER(40) E1 05520
IMPROV = 0 E1 05530
T = 0 E1 05540
C ORDER GATES IN GORDER E1 05550
EFLAG = 0 E1 05560
GO TO 63 E1 05570
C THIS ENTRY POINT FOR CALCULATION OF GORDER ONLY E1 05580
ENTRY FORMGO E1 05590
EFLAG = 1 E1 05600
63 CONTINUE E1 05610
COUNT = 0 E1 05620
DO 1 I=1,LEVM E1 05630
NMINLV = LGLIST(I) E1 05640
IF(NMINLV.EQ.0)GOTO1 E1 05650
DO 2 J=1,NMINLV E1 05660
COUNT = COUNT + 1 E1 05670
GORDER(COUNT) = HLIST(J,I) E1 05680
2 CONTINUE E1 05690
1 CONTINUE E1 05700
IF(EFLAG.EQ.1)RETURN E1 05710
C CALCULATE NUMBER OF OUTPUTS OF EACH GATE E1 05720
C (THE ARRAY 'USED' IS USED HERE JUST TEMPORARILY) E1 05730
DO 51 I=N1,NR E1 05740
TCOUNT = 0 E1 05750
DO 52 J=1,NR E1 05760
IF(INC$MX(I,J).EQ.1)TCOUNT = TCOUNT + 1 E1 05770

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52	CONTINUE	E1	05780
	TCOUNT NOW CONTAINS THE NUMBER OF OUTPUTS OF GATE I	E1	05790
	USED(I) = TCOUNT	E1	05800
51	CONTINUE	E1	05810
	MOST = 0	E1	05820
	DO 53 I = N1, NR	E1	05830
	IF(USED(I).GT.MOST) MOST = USED(I)	E1	05840
53	CONTINUE	E1	05850
	DO 56 I = 1, N	E1	05860
56	TORDER(I) = I	E1	05870
	TPOINT = N1	E1	05880
	MOST = MOST + 1	E1	05890
50	MOST = MOST - 1	E1	05900
	IF(MOST.LT.0) GO TO 54	E1	05910
	DO 55 I = 1, NR	E1	05920
	II = GORDER(I)	E1	05930
	IF(II.LE.N) GO TO 55	E1	05940
	IF(USED(II).NE.MOST) GO TO 55	E1	05950
	TORDER(TPOINT) = II	E1	05960
	TPOINT = TPOINT + 1	E1	05970
55	CONTINUE	E1	05980
	GO TO 50	E1	05990
54	CONTINUE	E1	06000
	INITIALIZE GSMALL	E1	06010
	DO 4 I = N1, NM	E1	06020
	X = (I-1)*N2	E1	06030
	DO 4 J = 1, N2	E1	06040
	Y = P\$(1,X+J)	E1	06050
	IF(Y.EQ.0) GSMALL(I,J) = -100	E1	06060
	IF(Y.EQ.1) GSMALL(I,J) = 1	E1	06070
	IF(Y.EQ.-1) GSMALL(I,J) = 0	E1	06080
4	CONTINUE	E1	06090
	EFLAG = 0	E1	06100
	GO TO 57	E1	06110
	ENTRY INITGS	E1	06120
	EFLAG = 1	E1	06130
57	DO 3 I = 1, NR	E1	06140
	USED(I) = 0	E1	06150
	IF(I.LT.N1) GO TO 58	E1	06160
	IF(I.GT.NM) GO TO 58	E1	06170
	GO TO 3	E1	06180
58	DO 59 J = 1, N2	E1	06190
59	GSMALL(I,J) = 0	E1	06200
3	CONTINUE	E1	06210
	DO 62 I = N1, NM	E1	06220
	USED(I) = 1	E1	06230
62	CONTINUE	E1	06240
	INITIALIZATION	E1	06250
	DO 34 I = 1, NR	E1	06260
	GATE = GORDER(I)	E1	06270
	IF(GATE.LT.N1) GO TO 34	E1	06280
	XX = LIPRED(GATE)	E1	06290
	IF(XX.EQ.0) GO TO 34	E1	06300
	F\$UP1 = 0	E1	06310
	F\$UB0 = 0	E1	06320
	DO 35 J = 1, N2	E1	06330
	COMPNT = GSMALL(GATE,J)	E1	06340
	IF(COMPNT.EQ.0) GO TO 35	E1	06350
	IF(COMPNT.LT.0) GO TO 36	E1	06360
	IF(COMPNT.GE.1000) GO TO 35	E1	06370
	F\$JB0 = F\$UB0 + 1	E1	06380

	F\$0(F\$UR0) = J	E1 06390
	GO TO 35	E1 06400
36	IF(COMPNT.LE.-1000) GO TO 35	E1 06410
	F\$UR1 = F\$UR1 + 1	E1 06420
	F\$1(F\$UR1) = J	E1 06430
35	CONTINUE	E1 06440
	IF(F\$UR1.EQ.0)GO TO 34	E1 06450
	DO 38 K=1,XX	E1 06460
	FEEDGT = IPRED(K,GATE)	E1 06470
	X = (FEEDGT-1)*N2	E1 06480
	DO 39 L=1,F\$UR1	E1 06490
	Y = F\$1(L)	E1 06500
	IF(P\$1(X+Y).LE.0)GO TO 39	E1 06510
	IF(GSMALL(FEEDGT,Y).GT.1000)GOTO39	E1 06520
	IF(GSMALL(GATE,Y).EQ.-200)GOTO39	E1 06530
	IF(GSMALL(GATE,Y).EQ.-100)GO TO 40	E1 06540
	GSMALL(GATE,Y) = -200	E1 06550
	GO TO 39	E1 06560
40	GSMALL(GATE,Y) = -FEEDGT	E1 06570
39	CONTINUE	E1 06580
38	CONTINUE	E1 06590
	DO 60 K=1,XX	E1 06600
60	MARKED(IPRED(K,GATE)) = 0	E1 06610
	DO 41 K=1,F\$UR1	E1 06620
	X = GSMALL(GATE,F\$1(K))	E1 06630
	IF(X.EQ.-100)GO TO 41	E1 06640
	IF(X.EQ.-200)GOTO41	E1 06650
	X = -X	E1 06660
	GSMALL(+X,F\$1(K))=1	E1 06670
	USED(X) = 1	E1 06680
	IF(MARKED(X).EQ.1)GOTO41	E1 06690
	MARKED(X) = 1	E1 06700
	DO 42 L=1,F\$UR0	E1 06710
	Y = GSMALL(X,F\$0(L))	E1 06720
	IF(Y.GT.1000.OR.Y.LT.-1000)GO TO 42	E1 06730
	GSMALL(+X,F\$0(L))=-100	E1 06740
42	CONTINUE	E1 06750
41	CONTINUE	E1 06760
34	CONTINUE	E1 06770
	IF(FFLAG.EQ.1)RETURN	E1 06780
C	INITIALIZE COUNTER TO LOOP ONCE FOR EACH GATE	E1 06790
	GDCOUNT = 0	E1 06800
C	INCREMENT GDCOUNT	E1 06810
5	GDCOUNT = GDCOUNT + 1	E1 06820
C	ARE ALL GATES EXHAUSTED?	E1 06830
	IF(GDCOUNT.LE.NR)GO TO 6	E1 06840
	IF(T.GT.0) IMPROV = 1	E1 06850
	IF(IMPROV.EQ.0)RETURN	E1 06860
C	IF HERE, NETWORK WAS ALTERED, SO UPDATE ARRAYS	E1 06870
	CALL SUBNET	E1 06880
	CALL PVALUE	E1 06890
	RETURN	E1 06900
6	GCD = GORDER(GDCOUNT)	E1 06910
C	IS GCD AN ISOLATED GATE OR EXTERNAL VARIABLE?	E1 06920
	IF(GCD.LE.V)GOTO5	E1 06930
	DO 8 I=1,N2	E1 06940
	IF(GSMALL(GCD,I).GE.1)GOTO7	E1 06950
8	CONTINUE	E1 06960
C	IF HERE, GATE IS ISOLATED - REMOVE INPUTS	E1 06970
	X = LIPRED(GCD)	E1 06980
	IF(X.EQ.0)GOTO5	E1 06990

	DO 9 I=1,X	E1 07000
	Y = IPRED(I,GCC)	E1 07010
	INC\$MX(Y,GCC) = 0	E1 07020
7	RECORD THE DISCONNECTION	E1 07030
	T = T + 1	E1 07040
9	CONTINUE	E1 07050
	GOTO 5	E1 07060
C	REMOVE UNNECESSARY CONNECTIONS TO GCC IN THE NEXT FEW SECTIONS	E1 07070
C		E1 07080
C	CALCULATE F(GCC)	E1 07090
7	F\$JBI = 0	E1 07100
	DO 10 I=1,N2	E1 07110
	IF(GSMALL(GCC,I).GE.0)GOTO10	E1 07120
	F\$UB1 = F\$JBI + 1	E1 07130
	F\$(F\$UB1) = I	E1 07140
10	CONTINUE	E1 07150
	DO 11 I=1,F\$UB1	E1 07160
11	INPTCV(F\$(I)) = 0	E1 07170
	X = LIPRED(GCC)	E1 07180
	DO 222 I=1,X	E1 07190
	ESSIS(IPRED(I,GCC)) = 0	E1 07200
222	CONTINUE	E1 07210
	T1SUB = 0	E1 07220
	T2SUB = 0	E1 07230
	DO 48 I = 1,NR	E1 07240
	IF(INC\$MX(I,GCC).EQ.0)GOTO48	E1 07250
	T1SUB = T1SUB + 1	E1 07260
	T1PRED(T1SUB) = I	E1 07270
48	CONTINUE	E1 07280
17	DO 18 I=1,X	E1 07290
	Y = (T1PRED(I)-1)*N2	E1 07300
	DO 19 J=1,F\$UB1	E1 07310
	Q = F\$(J)	E1 07320
	IF(P\$(1,Y+Q).NE.1)GO TO 19	E1 07330
	IF(INPTCV(Q).LE.0) GO TO 20	E1 07340
	INPTCV(Q) = INPTCV(Q) + 1	E1 07350
	GO TO 19	E1 07360
20	IF(INPTCV(Q).LT.0)GO TO 21	E1 07370
	INPTCV(Q) = -T1PRED(I)	E1 07380
	GO TO 19	E1 07390
21	INPTCV(Q) = 2	E1 07400
19	CONTINUE	E1 07410
18	CONTINUE	E1 07420
C	MARK ESSENTIAL 1'S	E1 07430
	DO 22 I=1,F\$UB1	E1 07440
	Q = INPTCV(F\$(I))	E1 07450
	IF(Q.GE.0)GO TO 22	E1 07460
	ESSIS(-Q) = ESSIS(-Q) + 1	E1 07470
22	CONTINUE	E1 07480
46	SELECT = 0	E1 07490
	BESTSL = 0	E1 07500
	DO 45 L=1,X	E1 07510
	Q = T1PRED(L)	E1 07520
	IF(INC\$MX(Q,GCC).EQ.0)GOTO45	E1 07530
	IF(ESSIS(Q).GT.0)GOTO45	E1 07540
	IF(SELECT.EQ.0)SELECT = Q	E1 07550
	IF(USED(Q).EQ.1)GOTO45	E1 07560
	IF(BESTSL.NE.0)GOTO45	E1 07570
	BESTSL = Q	E1 07580
45	CONTINUE	E1 07590
	IF(SELECT.EQ.0)GO TO 47	E1 07600

	Q = SELECT	E1 07610
	IF(PESTSL.NE.0)Q = RFSTSL	E1 07620
C	IF HERE, GATE HAS NO ESSENTIAL 1'S - REMOVE IT	E1 07630
	INC\$MX(Q,GCD) = 0	E1 07640
	T = T + 1	E1 07650
C	UPDATE ESSIS	E1 07660
	Y = (Q - 1)*N2	E1 07670
	DO 24 J=1,F\$JUB1	E1 07680
	V = F\$1(J)	E1 07690
	IF(P\$(1,Y+V).NE.1)GO TO 24	E1 07700
C	UPDATE INPTCV FOR COMPONENT V	E1 07710
	INPTCV(V) = INPTCV(V) - 1	E1 07720
	IF(INPTCV(V).GT.1)GO TO 24	E1 07730
C	CASE WHEN NEW ESSEN 1 CREATED	E1 07740
	DO 27 K = 1,X	E1 07750
	W = T1PRED(K)	E1 07760
	IF(INC\$MX(W,GCD).EQ.0) GO TO 27	E1 07770
	Z = (W - 1) * N2	E1 07780
	IF(P\$(1,Z+V).EQ.0)GO TO 27	E1 07790
	ESSIS(W) = ESSIS(W) + 1	E1 07800
C	IN THIS CASE, NO NEED TO UPDATE INPTCV FURTHER	E1 07810
	GSMALL(GCD,V) = -W	E1 07820
	GO TO 24	E1 07830
	27 CONTINUE	E1 07840
	24 CONTINUE	E1 07850
	GO TO 46	E1 07860
	47 DO 49 I = 1,NR	E1 07870
	IF(INC\$MX(I,GCD).EQ.0)GO TO 49	E1 07880
	T2SUB = T2SUB + 1	E1 07890
	T2PRED(T2SUB) = I	E1 07900
	49 CONTINUE	E1 07910
C	NOW ALL CURRENT INPUTS HAVE ESSENTIAL 1'S	E1 07920
C	INPUTS STILL CONNECTED TO GCD ARE LISTED IN T2PRED IN REVERSE	E1 07930
C	ORDER	E1 07940
C		E1 07950
C	UPDATE G(I)'S OF THOSE GATES STILL CONNECTED TO GATE GCD	E1 07960
		E1 07970
	DO 29 II=1,F\$JUB1	E1 07980
	T = F\$1(II)	E1 07990
	CHOICE = -GSMALL(GCD,I)	E1 08000
	IF(CHOICE.LT.100)GO TO 61	E1 08010
	CHOICE = 0	E1 08020
	DO 30 JJJ=1,NP	E1 08030
	JJ = TORDER(JJJ)	E1 08040
	IF(INC\$MX(JJ,GCD).EQ.0)GO TO 30	E1 08050
	IF(P\$(1,(JJ-1)*N2+I).NE.1)GO TO 30	E1 08060
	IF(JJ.LE.N)GO TO 29	E1 08070
	IF(CHOICE.EQ.0)CHOICE=JJ	E1 08080
	IF(GSMALL(JJ,I).GE.1)GO TO 29	E1 08090
30	CONTINUE	E1 08100
61	GSMALL(CHOICE,I) = 1	E1 08110
	USED(CHOICE) = 1	E1 08120
29	CONTINUE	E1 08130
	DO 32 I=1,N2	E1 08140
	IF(GSMALL(GCD,I).LT.1)GO TO 32	E1 08150
	DO 33 J=1,T2SUB	E1 08160
	IF(GSMALL(T2PRED(J),I).EQ.0)GSMALL(T2PRED(J),I)=-100	E1 08170
33	CONTINUE	E1 08180
32	CONTINUE	E1 08190
	GO TO 5	E1 08200
	END	E1 08210



```

SUBROUTINE ORDRQ2                                     E1 08220
C**** THIS SUBROUTINE MAKES A LIST OF PREDECESSORS OF GI   E1 08230
C ACCORDING TO ORDERING Q2                                E1 08240
C LIP IS STORED AT COMMON STORAGE                        ****E1 08250
C                                                         E1 08260
C DEFINITIONS OF 'COMMON' VARIABLES CAN BE FOUND IN MAIN PROGRAM. E1 08270
C                                                         E1 08280
C IMPLICIT INTEGER*4(A-T,V-Z,$), REAL(U)                E1 08290
COMMON NEPMAX                                           E1 08300
COMMON N , M , A , B                                     E1 08310
1 , R , N2 , N1 , NR                                     E1 08320
2 , NM , KFLAG , JFLAG , COST                           E1 08330
3 , LEVM , NRN2 , NM1 , NN2                             E1 08340
COMMON ISJCC(40,40) , LISUCC(40) , IPRED(40,40) , LIPRED(40) E1 08350
1 , INC$MX(40,40) , SUC$MX(40,40) , P$(2,1280) , UNAME(40) E1 08360
2 , GLEVFL(40) , LGLIST(40) , HLIST(40,40) , TIME       E1 08370
COMMON T , RTCONN(100) , S , RSCONN(100)               E1 08380
COMMON IFLAG , POINTA , ESSIS(40) , F$1(32)             E1 08390
1 , F$UB1 , INPTCV(32) , LISTC(40) , POINTC             E1 08400
2 , LISTL(40) , POINTL , ORIGIN(40) , IPATH(40)         E1 08410
3 , POINTR , VF$1(32) , VF$UB1 , GSMALL(40,32)          E1 08420
COMMON POTAB(200,42) , PPOTAB(40) , LPOTAB(40) , NRPLC(2) E1 08430
1 , RPLC(2,40) , IDX0(32) , IDX0E(32) , IDX1(32)        E1 08440
2 , IDX1E(32) , SUMP(32) , SETT1(32) , NOT1            E1 08450
3 , SETS1(40) , NOS1 , SETS(40) , NOS                  E1 08460
4 , STS , SUMS2(32) , SETS2(200) , NOS2                E1 08470
5 , LIP , NOOE , KEYA , KEYB                            E1 08480
6 , NOC , NO1 , NO1E , $GT                              E1 08490
7 , $LTH , $PW , $NOE , GI                              E1 08500
COMMON NOT1SV , NOS1SV , LMTS2                         E1 08510
DIMENSION WRPLC(2,40)                                  E1 08520
WRPLC(1) = 0                                             E1 08530
NRPLC(2) = 0                                             E1 08540
DO 8100 LI=1,LIP                                         E1 08550
GP=IPRED(LI,GI)                                          E1 08560
IF(GP.LE.N) GO TO 8100                                  E1 08570
IF(INC$MX(GP,GI).EQ.0) GO TO 8100                       E1 08580
BSGP=(GP-1)*N2                                          E1 08590
NOONEE = 0                                              E1 08600
ESSN=1                                                  E1 08610
DO 8050 NO=1,NOOE                                       E1 08620
TH=IDX0E(NO)                                            E1 08630
IF(P$(1,BSGP+TH).LE.0) GO TO 8050                      E1 08640
NOONEE=NOONEE+1                                         E1 08650
IF(SUMP(TH).EQ.1) ESSN=2                               E1 08660
8050 CONTINUE                                           E1 08670
C**** PUT GP INTO RPLC(1,*) OR RPLC(2,*) DEPENDING ON ESSN E1 08680
C ESSN=1 : NO ESSENTIAL ERROR                           E1 08690
C ESSN=2 : WITH ESSENTIAL ERRORS                        E1 08700
C RPLC TABLES ARE STORED ACCORDING TO ORDERING Q1(NOONEE) E1 08710
C NRPLC(ESSN) :NUMBER OF ELEMENTS IN RPLC(ESSN)**      E1 08720
IF(NOONEE.EQ.0) GO TO 8100                              E1 08730
P2=NRPLC(ESSN)                                          E1 08740
IF(P2.EQ.0) GO TO 8070                                  E1 08750
DO 8060 PR=1,P2                                         E1 08760
RP=P2-PR+1                                              E1 08770
IF(WRPLC(ESSN,RP).LE.NOONEE) GO TO 8080               E1 08780
WRPLC(ESSN,RP+1)=WRPLC(ESSN,RP)                      E1 08790
RPLC(ESSN,RP+1)=RPLC(ESSN,RP)                        E1 08800

```

8060	CONTINUE	E1 08810
8070	RP=0	E1 08820
8080	WRPLC(ESSN,RP+1)=NOONFF	E1 08830
	RPLC(ESSN,RP+1)=GP	E1 08840
	NRPLC(ESSN)=NRPLC(ESSN)+1	E1 08850
8100	CONTINUE	E1 08860
	RETURN	E1 08870
	END	E1 08880

SUBROUTINE OUTPUT(MATRIX,ARRAY)

E1 08890

DEFINITIONS OF 'COMMON' VARIABLES CAN BE FOUND IN MAIN PROGRAM.

E1 08900

E1 08910

E1 08920

IMPLICIT INTEGER\*4(A-T,V-Z,\$), REAL(U)

E1 08930

COMMON NEPMAX

E1 08940

COMMON	N	, M	, A	, B	E1 08950
--------	---	-----	-----	-----	----------

1	, P	, N2	, N1	, NR	E1 08960
---	-----	------	------	------	----------

2	, NM	, KFLAG	, JFLAG	, COST	E1 08970
---	------	---------	---------	--------	----------

3	, LEVM	, NPM2	, NM1	, NN2	E1 08980
---	--------	--------	-------	-------	----------

COMMON	ISJCC(40,40)	, LISJCC(40)	, IPRED(40,40)	, LIPRED(40)	E1 08990
--------	--------------	--------------	----------------	--------------	----------

1	, INC\$MX(40,40)	, SUC\$MX(40,40)	, P\$(2,1280)	, UNAME(40)	E1 09000
---	------------------	------------------	---------------	-------------	----------

2	, GLEVEL(40)	, LGLIST(40)	, HLIST(40,40)	, TIME	E1 09010
---	--------------	--------------	----------------	--------	----------

COMMON	T	, RTCONN(100)	, S	, RSCONN(100)	E1 09020
--------	---	---------------	-----	---------------	----------

COMMON	IFLAG	, POINTA	, ESSIS(40)	, F\$1(32)	E1 09030
--------	-------	----------	-------------	------------	----------

1	, F\$UB1	, INPTCV(32)	, LISTC(40)	, POINTC	E1 09040
---	----------	--------------	-------------	----------	----------

2	, LISTL(40)	, POINTL	, ORIGIN(40)	, IPATH(40)	E1 09050
---	-------------	----------	--------------	-------------	----------

3	, POINTP	, VF\$1(32)	, VF\$UB1	, GSMALL(40,32)	E1 09060
---	----------	-------------	-----------	-----------------	----------

COMMON	POTAB(200,42)	, PPOTAB(40)	, LPOTAB(40)	, NRPLC(2)	E1 09070
--------	---------------	--------------	--------------	------------	----------

1	, RPLC(2,40)	, IDX2(32)	, IDXOE(32)	, IDX1(32)	E1 09080
---	--------------	------------	-------------	------------	----------

2	, IDXIF(32)	, SUMP(32)	, SETT1(32)	, NOT1	E1 09090
---	-------------	------------	-------------	--------	----------

3	, SETS1(40)	, NOS1	, SETS(40)	, NOS	E1 09100
---	-------------	--------	------------	-------	----------

4	, STS	, SUMS2(32)	, SETS2(200)	, NOS2	E1 09110
---	-------	-------------	--------------	--------	----------

5	, LIP	, NOOF	, KEYA	, KEYB	E1 09120
---	-------	--------	--------	--------	----------

6	, NOO	, NO1	, NO1E	, \$GT	E1 09130
---	-------	-------	--------	--------	----------

7	, \$LTH	, \$PW	, \$NOE	, GI	E1 09140
---	---------	--------	---------	------	----------

COMMON		, NDT1SV	, NOS1SV	, LMTS2	E1 09150
--------	--	----------	----------	---------	----------

DIMENSION	JX(5), UY(5), UG(40), UF(40), ARRAY(40), ARRAY2(2,1280)	E1 09160
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DIMENSION	MATRIX(40,40)	E1 09170
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DATA	UX /' X1', ' X2', ' X3', ' X4', ' X5'/	E1 09180
------	--	----------

DATA	UY/' Y1', ' Y2', ' Y3', ' Y4', ' Y5'/	E1 09190
------	---------------------------------------	----------

DATA	JF /' 1', ' 2', ' 3', ' 4', ' 5', ' 6', ' 7', ' 8'	E1 09200
------	--	----------

1	, ' 9', ' 10', ' 11', ' 12', ' 13', ' 14', ' 15', ' 16'	E1 09210
---	---	----------

2	, ' 17', ' 18', ' 19', ' 20', ' 21', ' 22', ' 23', ' 24'	E1 09220
---	--	----------

3	, ' 25', ' 26', ' 27', ' 28', ' 29', ' 30', ' 31', ' 32'	E1 09230
---	--	----------

4	, ' 33', ' 34', ' 35', ' 36', ' 37', ' 38', ' 39', ' 40'	E1 09240
---	--	----------

DATA	GMAX/40/	E1 09250
------	----------	----------

E1 09260

E1 09270

E1 09280

E1 09290

E1 09300

E1 09310

E1 09320

E1 09330

E1 09340

E1 09350

E1 09360

E1 09370

E1 09380

E1 09390

KEYXC=ARRAY(1)

IF(KEYXC.NE.0) GO TO 50

DO 1 GI=1,N

UNAME(GI)=UX(GI)

1 CONTINUE

GO TO 100

50 CONTINUE

L=N/2

DO 4 GI=1,L

UNAME(GI)=UX(GI)

UNAME(GI+L)=UY(GI)

4 CONTINUE

100 CONTINUE



DO 2 GI=N1, GMAX	E1 09400
UNAME(GI)=UF(GI-N)	E1 09410
2 CONTINUE	E1 09420
RETURN	E1 09430
ENTRY LINE(L)	E1 09440
DO 6 LL=1, L	E1 09450
PRINT 5	E1 09460
5 FORMAT(1H )	E1 09470
6 CONTINUE	E1 09480
RETURN	E1 09490
ENTRY PAGE	E1 09500
PRINT 7	E1 09510
7 FORMAT(1H1)	E1 09520
RETURN	E1 09530
ENTRY CKT(MATRIX, ARRAY)	E1 09540
PRINT 10	E1 09550
10 FORMAT(1H , 8X, 'GATE .. LEVEL', 6X, 'FED BY'//)	E1 09560
DO 20 GJ=N1, NR	E1 09570
G=0	E1 09580
DO 15 GI=1, NR	E1 09590
IF(MATRIX(GI, GJ).EQ.0) GO TO 15	E1 09600
G=G+1	E1 09610
UG(G)=UNAME(GI)	E1 09620
15 CONTINUE	E1 09630
IF(G.EQ.0) GO TO 18	E1 09640
PRINT 17, UNAME(GJ), ARRAY(GJ), (UG(GG), GG=1, G)	E1 09650
17 FORMAT(1H0, 9X, A3, 5X, '/', I2, '/', 5X, 35( A3))	E1 09660
GO TO 20	E1 09670
18 PRINT 19, UNAME(GJ), ARRAY(GJ)	E1 09680
19 FORMAT(1H0, 9X, A3, 5X, '/', I2, '/')	E1 09690
20 CONTINUE	E1 09700
RETURN	E1 09710
ENTRY TRUTH(ARRAY2, J)	E1 09720
IF(J.EQ.2) GO TO 36	E1 09730
PRINT 35	E1 09740
35 FORMAT(11X, 'TRUTH TABLE'//)	E1 09750
GO TO 38	E1 09760
36 PRINT 37	E1 09770
37 FORMAT(11X, 'REQUIREMENT TABLE')	E1 09780
38 CONTINUE	E1 09790
DO 40 GI=1, NR	E1 09800
ILO=(GI-1)*N2+1	E1 09810
IHI=ILO+N2-1	E1 09820
PRINT 41, UNAME(GI), (ARRAY2(J, I), I=ILO, IHI)	E1 09830
40 CONTINUE	E1 09840
41 FORMAT(1H0, 9X, A3, ' = ', 32(I1, 1X))	E1 09850
RETURN	E1 09860
END	E1 09870
	E1 09880
	E1 09890
	E1 09900
	E1 09910
SUBROUTINE POT	E1 09920
	E1 09930
THIS SUBROUTINE GENERATES THE POTENTIAL OUTPUT TABLE	E1 09940
	E1 09950
POTAB(P1, P2) STORES POTENTIAL OUTPUT TABLE	E1 09960
P1=1, \$MXPTR: FUNCTION ENTRY NUMBER	E1 09970
P2=1, 32 : VALUE OF EACH COMPONENT OF THAT FUNCTION	E1 09980

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C      P2=33($GT) : GATE NUMBER WHERE THE FUNCTION IS REALIZED      E1 09990
C      P2=34($LTH): NUMBER OF CONNECTIONS TO BE ADDED              E1 10000
C      P2=35,40   : GATES WHICH ARE TO BE CONNECTED TO POTAB(*,$GT) E1 10010
C      P2=41($PW) : PREFERENCE WEIGHT                               E1 10020
C      P2=42($NOF): NUMBER OF ONE ERRORS                           E1 10030
C                                                                E1 10040
C      DEFINITIONS OF 'COMMON' VARIABLES CAN BE FOUND IN MAIN PROGRAM. E1 10050
C                                                                E1 10060
      IMPLICIT INTEGER*4(A-T,V-Z,$), REAL(U)                        E1 10070
      COMMON NEPMAX                                                  E1 10080
      COMMON      N          , M          , A          , B          E1 10090
1      ,      R          , N2         , N1         , NP         E1 10100
2      ,      NM         , KFLAG      , JFLAG      , COST       E1 10110
3      ,      LEVM       , NPN2      , NM1         , NN2        E1 10120
      COMMON      ISUCC(40,40) , LISUCC(40) , JPRED(40,40) , LIPRED(40) E1 10130
1      ,      INC$MX(40,40) , SUC$MX(40,40) , P$(2,1280) , UNAME(40) E1 10140
2      ,      GLEVFL(40) , LGLIST(40) , HLIST(40,40) , TIME      E1 10150
      COMMON      T          , RTCONN(100) , S          , RSCONN(100) E1 10160
      COMMON      IFLAG      , POINTA     , ESSIS(40) , F$1(32)      E1 10170
1      ,      F$UR1      , INPTCV(32) , LISTC(40) , POINTC      E1 10180
2      ,      LISTL(40) , POINTL      , ORIGIN(40) , IPATH(40)     E1 10190
3      ,      POINTR      , VF$1(32) , VF$UB1     , GSMALL(40,32) E1 10200
      COMMON      POTAB(200,42) , PPOTAB(40) , LPOTAB(40) , NRPLC(2) E1 10210
1      ,      RPLC(2,40) , IDX0(32) , IDX0E(32) , IDX1(32)      E1 10220
2      ,      IDX1E(32) , SUMP(32) , SETT1(32) , NOT1          E1 10230
3      ,      SETS1(40) , NOS1      , SETS(40) , NOS          E1 10240
4      ,      STS      , SUMS2(32) , SETS2(200) , NOS2         E1 10250
5      ,      LIP      , NOOF      , KEYA      , KEYB         E1 10260
6      ,      NOO      , NO1      , NO1E      , $GT          E1 10270
7      ,      $LTH      , $PW      , $NOE      , GI          E1 10280
      COMMON      NOT1SV      , NOS1SV      , LMTS2          E1 10290
      DIMENSION INDEX(32)                                           E1 10300
      DATA $MXPTR/200/                                             E1 10310
C**** INITIALIZE PPOTAB(*) ****                                     E1 10320
      DO 90 GI=1,NP                                                 E1 10330
90      PPOTAB(GI)=0                                               E1 10340
      POINTR = 1                                                    E1 10350
      DO 980 LEVV=1,LEVM                                           E1 10360
      LEV=LEVM-LEVV+1                                              E1 10370
      LGL=LGLIST(LEV)                                              E1 10380
      DO 960 LG=1,LGL                                              E1 10390
      GI=HLIST(LG,LEV)                                             E1 10400
      IF(LEV.GT.1) GO TO 100                                       E1 10410
      IF(GI.GT.NM.CR.GI.LE.N.CR.M.EQ.1) GO TO 960                E1 10420
100      LISI=LISUCC(GI)                                           E1 10430
      BSGI=(GI-1)*N2                                              E1 10440
      IF(POINTR.GT.$MXPTR) GO TO 990                               E1 10450
      PPOTAB(GI)=POINTR                                           E1 10460
C**** COPY PRESENT OUTPUT                                          E1 10470
      DO 110 TH=1,N2                                              E1 10480
      POTAB(POINTR,TH)=P$(1,BSGI+TH)                             E1 10490
110      CONTINUE                                                  E1 10500
      POTAB(POINTR,$GT)=GI                                         E1 10510
      POTAB(POINTR,$LTH)=0                                         E1 10520
      POINTR=POINTR+1                                              E1 10530
      IF(GI.LE.NM) GO TO 950                                       E1 10540
      DO 380 LEVJ=LEV,LEVM                                         E1 10550
      LGLJ=LGLIST(LEVJ)                                           E1 10560
      DO 360 LGJ=1,LGLJ                                           E1 10570
      GJ=HLIST(LGJ,LEVJ)                                           E1 10580
      IF(INC$MX(GJ,GI).GT.0.CR.GI.EQ.GJ) GO TO 360              E1 10590

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C**** CHECK IF GJ IS CONNECTED TO ALL SUCCESSORS OF GI          ****E1 10600
      DO 120 LIT=1,LIST                                          E1 10610
        IF(INC$MX(GJ,ISUCC(LII,GI)).LE.0) GO TO 360             E1 10620
      CONTINUE                                                    E1 10630
120 C**** CHECK IF GJ IS STRONGLY CONNECTIBLE TO GI          ****E1 10640
      BSGJ = (GJ-1)*N2                                           E1 10650
      NO = 0                                                       E1 10660
      DO 180 TH = 1,N2                                           E1 10670
        IF(P$(1,BSGI+TH).NE.1.OR.P$(1,BSGJ+TH).NE.1)GO TO 180 E1 10680
        NO=NO+1                                                    E1 10690
        INDEX(NO)=TH                                              E1 10700
      CONTINUE                                                    E1 10710
180 C**** NO=0 => NOT STRONGLY CONNECTABLE          ****E1 10720
      C**** NO>0 => STRONGLY CONNECTABLE          ****E1 10730
      IF(NC.EQ.0) GO TO 360                                       E1 10740
      IF(POINTR.GT.$MXPTR) GO TO 990                               E1 10750
      DO 200 TH=1,N2                                              E1 10760
        POTAB(POINTR,TH)=P$(1,BSGI+TH)                          E1 10770
      CONTINUE                                                    E1 10780
200 DO 210 NORUN=1,N2                                           E1 10790
      POTAB(POINTR,INDEX(NORUN))=0                                E1 10800
210 CONTINUE                                                    E1 10810
      SP=PPOTAB(GI)+1                                             E1 10820
      IF(POINTR.EQ.SP) GO TO 300                                   E1 10830
      EP = POINTR - 1                                             E1 10840
C**** CHECK IF THIS ENTRY IS SAME AS ONE OF THE PREVIOUS ENTRIES ****E1 10850
      DO 230 PTR=SP,EP                                           E1 10860
        DO 220 TH=1,N2                                           E1 10870
          IF(POTAB(POINTR,TH).NE.POTAB(PTR,TH)) GO TO 230      E1 10880
220 CONTINUE                                                    E1 10890
        GO TO 360                                                 E1 10900
230 CONTINUE                                                    E1 10910
300 POTAB(POINTR,$GT)=GI                                         E1 10920
      POTAB(POINTR,$LTH)=1                                       E1 10930
      POTAB(POINTR,$LTH+1)=GJ                                     E1 10940
      POINTR = POINTR + 1                                         E1 10950
360 CONTINUE                                                    E1 10960
380 CONTINUE                                                    E1 10970
C**** IF THE SET OF STRONGLY CONNECTABLE GATES CONTAINS MORE   E1 10980
      C THAN ONE GATE TRY THEIR COMBINATIONS          ****E1 10990
      C PTR1: STARTING POSITION OF THE LIST              E1 11000
      C PTR2: STARTING POSITION OF THE COMBINATIONS OF THE LIST E1 11010
      C PTR : THE ENTRY WHOSE COMBINATIONS WITH OTHERS ARE UNDER E1 11020
      C CONSIDERATION                                  E1 11030
      C PTRL: LAST ENTRY PRECEDING PTR                  E1 11040
      C PTRL2:LAST ENTRY OF COMBINATIONS OF ENTRIES PRECEDING PTR E1 11050
      C                                                  E1 11060
      PTR1=PPOTAB(GI)+1                                           E1 11070
      PTR2=POINTR                                                  E1 11080
      EP=PTR2-1                                                    E1 11090
      SP=PTR1+1                                                    E1 11100
      IF(SP.GT.EP) GO TO 950                                       E1 11110
      IF(PTR2-PTR1.GT.6) PRINT 1110                               E1 11120
1110 FORMAT(/////21X,'***** WARNING: NUMBER OF STRONGLY CONNECTABLE FUNE1 11130
      1CTIONS FOR A GATE EXCEEDS 6 *****'/
2      21X,'***** NOT ALL POSSIBLE OUTPUTS ARE AVAILABLE IN CE1 11150
3ALCULATION              *****'/////))                        E1 11160
      DO 560 PTR=SP,EP                                           E1 11170
        PTRL=PTR-1                                                E1 11180
        PTRL2=POINTR-1                                           E1 11190
C**** MAKE THE COMBINATIONS OF PTR AND ENTRIES PRECEDING IT AS E1 11200

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```

C          TH NEW ENTRIES                                     *****E1 11210
DO 450 PT=PTR1,PTPL                                          E1 11220
C***** MAKE NEW ENTRY AS THE COMBINATION OF ENTRIES PTR AND PT *****E1 11230
IF(POINTR.GT.$MXPTR) GO TO 990                                E1 11240
    DO 420 TH=1,N2                                           E1 11250
        POTAB(POINTR,TH)=1                                   E1 11260
        IF(POTAB(PTR,TH).EQ.0.OR.POTAB(PT,TH).EQ.0)         E1 11270
            POTAB(POINTR,TH)=0                               E1 11280
1          CONTINUE                                           E1 11290
420        POTAB(POINTR,$GT)=GI                               E1 11300
        POTAB(POINTR,$LTH)=2                                  E1 11310
        POTAB(POINTR,$LTH+1)=POTAB(PT,$LTH+1)               E1 11320
        POTAB(POINTR,$LTH+2)=POTAB(PTR,$LTH+1)              E1 11330
        POINTR=POINTR+1                                       E1 11340
450        CONTINUE                                           E1 11350
IF(PTR2.GT.PTRL2) GO TO 560                                  E1 11360
C***** MAKE THE COMBINATIONS OF PTR AND THE COMBINATIONS OF ENTRIES E1 11370
C PRECEDING PTR AS THE NEW ENTRIES *****E1 11380
DO 520 PT=PTR2,PTRL2                                         E1 11390
IF(POINTR.GT.$MXPTR) GO TO 990                                E1 11400
IF(POTAB(PT,$LTH).GE.6) GO TO 520                            E1 11410
    DO 480 TH=1,N2                                           E1 11420
        POTAB(POINTR,TH)=1                                   E1 11430
        IF(POTAB(PTR,TH).EQ.0.OR.POTAB(PT,TH).EQ.0)         E1 11440
            POTAB(POINTR,TH)=0                               E1 11450
7          CONTINUE                                           E1 11460
480        LTH=POTAB(PT,$LTH)+2                               E1 11470
        DO 500 TH=1,LTH                                       E1 11480
            POTAB(POINTR,$GT+TH-1)=POTAB(PT,$GT+TH-1)        E1 11490
500        CONTINUE                                           E1 11500
            POTAB(POINTR,$LTH)=POTAB(POINTR,$LTH)+1          E1 11510
            POTAB(POINTR,$GT+LTH)=POTAB(PTR,$LTH+1)          E1 11520
            POINTR=POINTR+1                                    E1 11530
520        CONTINUE                                           E1 11540
560        CONTINUE                                           E1 11550
950        LPOTAB(GI)=POINTR-1                                 E1 11560
960        CONTINUE                                           E1 11570
980        CONTINUE                                           E1 11580
        RETURN                                              E1 11590
C***** NUMBER OF POSSIBLE OUTPUT TABLE ENTRIES EXCEEDS THE LIMIT *****E1 11600
990 CALL LINE(5)                                              E1 11610
        PRINT 1000,$MXPTR                                     E1 11620
1000 FORMAT(21X,'***** WARNING: NUMBER OF POSSIBLE OUTPUT TABLE ENTRIES E1 11630
1 EXCEEDS THE LIMIT($MXPTR=',I3,') *****'/
2          21X,'***** NOT ALL POSSIBLE OUTPUTS ARE AVAILABLE IN CALCULAE1 11640
3TION                                     *****'///)
        LPOTAB(GI) = $MXPTP                                  E1 11670
        RETURN                                              E1 11680
        END                                                  E1 11690

SUBROUTINE PROCCE(WORKED)                                     E1 11700
C EDITION AAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAA E1 11710
C IF PROCCE SUCCESSFULLY COMPENSATES ERRORS, 'WORKED' IS SET TO 1, OE1 11720
C 'WORKED' IS SET TO 0                                       E1 11730
C                                                         E1 11740
C DEFS. OF MOST 'COMMON' VARIABLES CAN BE FOUND IN MAIN PROGRAM. E1 11750
C                                                         E1 11760
C VARIABLE DEFINITIONS:                                     E1 11770
C EP: EP(I)=1 MEANS AT LEAST ONE NETWORK OUTPUT GATE HAS AN E1 11780
C ERRONEOUS OUTPUT IN THE I-TH COMPONENT WHEN PCO IS REMOVED E1 11790

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C      FROM THE NETWORK.  FP(I)=0 OTHERWISE. E1 11900
C  ERRORS: TOTAL NO. OF ERRORS IN NETWORK OUTPUTS WHEN PCO REMOVED. E1 11810
C  GATES: NUMBER OF GATES REMOVED FROM NETWORK BY CALL TO MINI2. E1 11820
C  IMPROV: A PARAMETER RETURNED BY MINI2. '1' MEANS MINI2 WAS ABLE E1 11830
C          TO REDUCE COST OF NETWORK. E1 11840
C  MAX: MAXIMUM NUMBER OF REQUIRED 1'S IN A CSPF VECTOR (AFTER E1 11850
C        CALLING MINI2) PLUS 1. E1 11860
C  MIN: ORIGINALLY SET TO ZERO, MIN IS INCREMENTED EACH TIME BY 1 E1 11870
C        UNTIL ITS VALUE EQUALS MAX. E1 11880
C  NEP: NO. OF ERROR POSITIONS FOR A GIVEN NETWORK AFTER A SE- E1 11890
C        LECTED GATE HAS BEEN REMOVED. AN ERROR POSITION IS A E1 11900
C        COMPONENT POSITION WHICH IS IN ERROR FOR AT LEAST ONE E1 11910
C        OUTPUT. E1 11920
C  NEPMAX: READ FROM INPUT CARDS, THIS PARAMETER IS PASSED TO PROCCE E1 11930
C          WHEN IT IS CALLED BY MAIN. IT REPRESENTS THE MAXIMUM E1 11940
C          ALLOWABLE NUMBER OF ERROR POSITIONS. IF AN ALTERED (I.E., E1 11950
C          SOME PCO REMOVED) NETWORK EXCEEDS THIS MAXIMUM, ERROR E1 11960
C          COMPENSATION IS NOT ATTEMPTED FOR THAT NETWORK. E1 11970
C  NETCUT: STORES OUTPUTS OF GATES IN ALTERED (PCO REMOVED) NETWORK. E1 11980
C  ONECNT: USED IN COUNTING NO. OF 1'S IN CSPF VECTOR OF A GATE. E1 11990
C  ONES: AFTER THE INITIAL CALCULATION OF THE CSPF SETS IN THE E1 12000
C        BEGINNING, ONES(GI) GIVES THE NUMBER OF 1'S IN THE CSPF E1 12010
C        VECTOR OF GI. THIS INFORMATION IS REQUIRED FOR GENERATING E1 12020
C        PORDER. E1 12030
C  ORGOUT: USED TO STORE ORIGINAL (UNALTERED) NETWORK OUTPUTS IN E1 12040
C          CODED FORM (SAME CODE AS IN GSMALL) AND (40,32) FORMAT. E1 12050
C  PCO: CURRENT GATE REMOVED FROM ORIGINAL NETWORK TO OBTAIN E1 12060
C        CURRENT ALTERED NETWORK. PCC = PORDER(PCOUNT). E1 12070
C  PCOUNT: A POINTER TO PORDER. E1 12080
C  PORDER: ORDERING OF GATES ACCORDING TO NUMBER OF 1'S IN THEIR E1 12090
C          CSPF VECTORS. GATES ARE INDIVIDUALLY REMOVED FROM ORIGI- E1 12100
C          NAL NETWORK IN THIS ORDER E1 12110
C  PSUB: USED AS A POINTER TO PORDER DURING ITS INITIALIZATION. E1 12120
C  QINC$M: STORES A COPY OF INC$MX FOR THE ORIGINAL NETWORK. E1 12130
C  START: POINTS TO BEGINNING OF LIST OF NETWORK OUTPUTS IN P$. E1 12140
C  STOP: POINTS TO END OF LIST OF NETWORK OUTPUTS IN P$. E1 12150
C E1 12160
C  I,J,NI,X,Y ARE USED AS JUST TEMPORARY VARIABLES. E1 12170
C E1 12180
C  HOW TO INCREASE CAPACITY OF SUBROUTINE. E1 12190
C  DIMENSION: PORDER(X) E1 12200
C              ONES(X) E1 12210
C              QINC$M(X,X) - X EQUAL TO MAX NO. OF GATES PLUS EX. VAR. E1 12220
C              EP(Y) - Y EQUAL TO: 2** (MAX ALLOWED NO OF EX VAR) E1 12230
C              NETOUT(X,Y) E1 12240
C              ORGOUT(X,Y) - X,Y AS ABOVE E1 12250
C E1 12260
C  IMPLICIT INTEGER*4(A-T,V-Z,$), REAL(U) E1 12270
C  COMMON NEPMAX E1 12280
C  COMMON N , M , A , B E1 12290
C 1 , R , N2 , N1 , NR E1 12300
C 2 , NM , KFLAG , JFLAG , COST E1 12310
C 3 , LFVM , NRN2 , NM1 , NN2 E1 12320
C  COMMON ISJCC(40,40) , LISUCC(40) , IPRED(40,40) , LIPRED(40) E1 12330
C 1 , INC$MX(40,40) , SUP$MX(40,40) , P$(2,1280) , UNAME(40) E1 12340
C 2 , GLEVEL(40) , LGLIST(40) , HLIST(40,40) , TIME E1 12350
C  COMMON T , RTCONN(100) , S , RSCONN(100) E1 12360
C  COMMON TFLAG , POINTA , ESSIS(40) , F$1(32) E1 12370
C 1 , F$UB1 , INPTCV(32) , LISTC(40) , POINTC E1 12380
C 2 , LISTL(40) , POINTL , ORIGIN(40) , IPATH(40) E1 12390
C 3 , POINTR , VF$1(32) , VF$UB1 , GSMALL(40,32) E1 12400

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COMMON  POTAB(200,42),PPOTAB(40)      ,LPOTAB(40)      ,NRPLC(2)      E1 12410
1      ,RPLC(2,40)      ,IDX0(32)      ,IDX0E(32)      ,IDX1(32)      E1 12420
2      ,IDX1E(32)      ,SUMP(32)      ,SETT1(32)      ,NOT1      E1 12430
3      ,SETT1(40)      ,NOS1      ,SETS(40)      ,NOS      E1 12440
4      ,STS      ,SUMS2(32)      ,SETS2(200)      ,NOS2      E1 12450
5      ,LIP      ,NODE      ,KEYA      ,KEYB      E1 12460
6      ,NOD      ,NOL      ,NOL      ,NOL      E1 12470
7      ,LTH      ,PW      ,NOL      ,GI      E1 12480
COMMON      NOT1SV      ,NOS1SV      ,LMTS2      E1 12490
DIMENSION PORDER(40),ONES(40),QINC$M(40,40),NETOUT(40,32),
1  EP(32),ORGOUT(40,32)      E1 12500
THIS SUBROUTINE ASSUMES ALL ARRAYS ARE UPDATED      E1 12510
PREVIOUS TO BEING CALLED      E1 12520
      E1 12530
      E1 12540
$GT = 33      E1 12550
$LTH = 34      E1 12560
$PW = 41      E1 12570
$NOL = 42      E1 12580
WORKED = 0      E1 12590
S = 0      E1 12600
T = 0      E1 12610
      E1 12620
BLOCK  B  B  B  B  B  B  B  B  B  B  B  B  B  B  B  B  B  B  E1 12630
      E1 12640
CALL MINI2(IMPROV)      E1 12650
IN THIS CALL TO MINI2, GORDER WILL BE CALCULATED. GORDER WILL BE      E1 12660
LATER IN EACH CALL TO INITGS (AN ENTRY POINT OF MINI2). NOTE THAT      E1 12670
IS NOT AFFECTED BY THE REMOVAL OF GATES FROM THE ORIGINAL NETWORK.      E1 12680
IF(IMPROV.EQ.0)GO TO 1      E1 12690
GATES = 0      E1 12700
DO 2 I = N1,NR      E1 12710
DO 3 J = N1,NR      E1 12720
IF(INC$MX(I,J).GT.0)GO TO 2      E1 12730
3  CONTINUE      E1 12740
GATES = GATES + 1      E1 12750
2  CONTINUE      E1 12760
PRINT 4,GATES,T      E1 12770
4  FORMAT(' ',I5,' GATES AND',I3,' CONNECTIONS HAVE BEEN REMOVED FROM      E1 12780
1  THE NETWORK DURING THE INITIAL CALCULATION OF THE CSPF SET')      E1 12790
1  CONTINUE      E1 12800
COUNT THE NUMBER OF 1'S IN THE CSPF VECTOR FOR EACH GATE      E1 12810
MAX = 0      E1 12820
DO 5 I = N1,NR      E1 12830
ONECNT = 0      E1 12840
DO 6 J = 1,N2      E1 12850
IF(GSMALL(I,J).LE.0)GO TO 6      E1 12860
ONECNT = ONECNT + 1      E1 12870
6  CONTINUE      E1 12880
IF(ONECNT.GT.MAX)MAX=ONECNT      E1 12890
ONES(I) = ONECNT      E1 12900
5  CONTINUE      E1 12910
MAX = MAX + 1      E1 12920
MIN = -1      E1 12930
PSUB = 1      E1 12940
7  MIN = MIN + 1      E1 12950
IF(MIN.EQ.MAX) GO TO 8      E1 12960
DO 9 I = N1,NR      E1 12970
IF(ONES(I).NE.MIN)GO TO 9      E1 12980
PORDER(PSUB) = I      E1 12990
PSUB = PSUB + 1      E1 13000
9  CONTINUE      E1 13010

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	GOTO 7	E1 13020
	8 CONTINUE	E1 13030
C	SAVE ORIGINAL NETWORK	E1 13040
	DO 10 I = 1,NR	E1 13050
	DO 10 J = 1,NR	E1 13060
	QINC\$M(I,J) = INC\$MX(I,J)	E1 13070
	10 CONTINUE	E1 13080
C	SAVE ORIGINAL OUTPUTS	E1 13090
C	SAVE ORIGINAL OUTPUTS IN (2,1280) FORMAT	E1 13100
	START = (N*N2) + 1	E1 13110
	STOP = (NM*N2)	E1 13120
	DO 13 I = START, STOP	E1 13130
	P\$(2,I) = P\$(1,I)	E1 13140
	13 CONTINUE	E1 13150
C	SAVE ORIGINAL OUTPUTS IN CODED (40,32) FORMAT	E1 13160
	DO 27 I = N1,NM	E1 13170
	X = ( I-1) * N2	E1 13180
	DO 28 J = 1,N2	E1 13190
	Y = P\$(1,X+J)	E1 13200
	IF(Y)30,31,32	E1 13210
C	COMPONENT IS DON'T CARE (I.E., -1)	E1 13220
	30 ORGOUT(I,J) = 0	E1 13230
	GOTO 28	E1 13240
C	COMPONENT IS LOGICAL ZERO	E1 13250
	31 ORGOUT(I,J) = -100	E1 13260
	GO TO 28	E1 13270
C	COMPONENT IS LOGICAL ONE	E1 13280
	32 ORGOUT(I,J) = 1	E1 13290
	28 CONTINUE	E1 13300
	27 CONTINUE	E1 13310
C		E1 13320
C	BLOCK C C C C C C C C C C C C C C C C C C C	E1 13330
C		E1 13340
	PCOUNT = 0	E1 13350
11	PCOUNT = PCOUNT + 1	E1 13360
	IF(PCOUNT.GT. R)GO TO 23	E1 13370
	PCD = PORDER(PCOUNT)	E1 13380
	IF(ONES(PCD).EQ.0)GO TO 11	E1 13390
	IF(PCD.LE.NM)GO TO 11	E1 13400
C	ERRORS UNCORRECTABLE, RESTORE NETWORK, TRY AGAIN	E1 13410
	DO 19 I = 1,NR	E1 13420
	DO 19 J = 1,NR	E1 13430
	INC\$MX(I,J) = QINC\$M(I,J)	E1 13440
	19 CONTINUE	E1 13450
C	REMOVE GATE PCD FROM THE NETWORK	E1 13460
	DO 12 I = 1,NR	E1 13470
	IF(INC\$MX(I,PCD).EQ.0)GO TO 34	E1 13480
	INC\$MX(I,PCD) = 0	E1 13490
34	IF(INC\$MX(PCD,I).EQ.0) GO TO 12	E1 13500
	INC\$MX(PCD,I) = 0	E1 13510
	12 CONTINUE	E1 13520
C	UPDATE GATE OUTPUTS FOR ALTERED NETWORK	E1 13530
C		E1 13540
C	BLOCK D D D D D D D D D D D D D D D D D D D	E1 13550
C		E1 13560
33	CALL SUBNET	E1 13570
	CALL PVALUE	E1 13580
	CALL UNNECE	E1 13590
C	RESTORE GSMALL FOR OUTPUT GATES	E1 13600
	DO 29 I = N1,NM	E1 13610
	DO 29 J = 1, N2	E1 13620



	GSMALL(I,J) = DFGOUT(I,J)	E1	13630
29	CONTINUE	E1	13640
	ERRORS = 0	E1	13650
	DO 24 I=1,N2	E1	13660
24	EP(I) = 0	E1	13670
	DO 14 I = 1,M	E1	13680
	NI = N + I	E1	13690
	X = (NI - 1) * N2	E1	13700
	DO 15 J = 1,N2	E1	13710
	IF(GSMALL(NI,J))16,15,17	E1	13720
	CASE WHERE REQUIREMENT IS A ZERO	E1	13730
16	IF(P\$(1,X+J).EQ.0)GO TO 15	E1	13740
	CASE OF ONE WITH ERROR	E1	13750
	GSMALL(NI,J) = 1001	E1	13760
	ERRORS = ERRORS + 1	E1	13770
	EP(J) = 1	E1	13780
	GO TO 15	E1	13790
	CASE WHERE REQUIREMENT IS A ONE	E1	13800
17	IF(P\$(1,X+J).EQ.1)GO TO 15	E1	13810
	CASE OF ZERO WITH ERROR	E1	13820
	GSMALL(NI,J) = -1100	E1	13830
	ERRORS = ERRORS + 1	E1	13840
	EP(J) = 1	E1	13850
15	CONTINUE	E1	13860
14	CONTINUE	E1	13870
	IF(ERRORS.EQ.0)WORKED = 1	E1	13880
	IF(ERRORS.EQ.0) RETURN	E1	13890
	NEP = 0	E1	13900
	DO 25 I = 1,N2	E1	13910
	IF(EP(I).EQ.0) GO TO 25	E1	13920
	NEP = NEP + 1	E1	13930
25	CONTINUE	E1	13940
	IF(NEP.GT.NEPMAX) GO TO 11	E1	13950
		E1	13960
	BLOCK E	E1	13970
		E1	13980
	CALL POT	E1	13990
	'POT' IS A SUBROUTINE THAT GENERATES THE POTENTIAL OUTPUT TABLE	E1	14000
		E1	14010
	BLOCK F	E1	14020
		E1	14030
	SAVE NEW NETWORK OUTPUTS	E1	14040
	DO 18 J = 1,N2	E1	14050
	DO 18 I = 1,NM	E1	14060
	NETOUT(I,J) = GSMALL(I,J)	E1	14070
18	CONTINUE	E1	14080
	CALL FORMGD	E1	14090
	CALL INITGS	E1	14100
	CALL RCEC(811,833)	E1	14110
		E1	14120
	BLOCK I	E1	14130
		E1	14140
	CASE OF ALL POSSIBLE GATE REMOVALS EXHAUSTED	E1	14150
23	DO 26 I = 1,NR	E1	14160
	DO 26 J = 1,NR	E1	14170
	INC\$MX(I,J) = QINC\$M(I,J)	E1	14180
26	CONTINUE	E1	14190
	CALL SUBNET	E1	14200
	CALL PVALUF	E1	14210
	RETURN	E1	14220
	END	E1	14230

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SUBROUTINE RCEC(*,*)
SUBROUTINE FOR REMOVING CONNECTIONS BY ERROR COMPENSATION
DEFINITIONS OF 'COMMON' VARIABLES CAN BE FOUND IN MAIN PROGRAM.
IMPLICIT INTEGER*4(A-T,V-Z,$), REAL(U)
COMMON NEPMAX
COMMON      N      , M      , A      , B
1      , R      , N2      , N1      , NR
2      , NM      , KFLAG      , JFLAG      , COST
3      , LEVM      , NRN2      , NM1      , NN2
COMMON      TSUCC(40,40) , LISUCC(40) , IPPED(40,40) , LIPRED(40)
1      , INC$MX(40,40) , SUC$MX(40,40) , P$(2,1280) , UNAME(40)
2      , GLEVEL(40) , LGLIST(40) , HLIST(40,40) , TIME
COMMON      T      , RTCONN(100) , S      , RSCONN(100)
COMMON      IFLAG      , POINTA      , ESS1S(40) , F$1(32)
1      , F$U$1      , INPTCV(32) , LISTC(40) , POINTC
2      , LISTL(40) , POINTL      , ORIGIN(40) , IPATH(40)
3      , P$NTR      , VF$1(32) , VF$UB1      , G$SMALL(40,32)
COMMON      POTAB(200,42) , PPOTAB(40) , LPOTAB(40) , NRPLC(2)
1      , RPLC(2,40) , IDX0(32) , IDX0E(32) , IDX1(32)
2      , IDX1E(32) , SJMP(32) , SETT1(32) , NOT1
3      , SETS1(40) , NOS1      , SETS(40) , NOS
4      , STS      , SUMS2(32) , SETS2(200) , NOS2
5      , LIP      , NOOE      , KEYA      , KEYB
6      , NOO      , NO1      , NO1E      , $GT
7      , $LTH      , $PW      , $NOE      , GI
COMMON      NOT1SV      , NOS1SV      , LMTS2
DIMENSION LISCND(200) , DI(10) , DIO(10) , BI(200) ,
1 RIO(200) , IDXK(32)
DIMENSION ORDRP4(40)
SELECT ONE GATE, GI, WHOSE CSPFE IS ALREADY CALCULATED.
ACCORDING TO THE ORDER P1
GI = 0
LEV = 0
100 LEV = LEV + 1
IF(LEV.GE.LEVM) RETURN1
LGL = LGLIST(LEV)
LG = 0
120 LG = LG + 1
IF(LG.GT.LGL)GO TO 100
GI = HLIST(LG,LEV)
GO TO 150
IN CASE THE ORDERING P1 HAS BEEN CHANGED DURING PREVIOUS CALCULA-
TION FOR GI (THE LEVEL OF GI CAN NEVER CHANGE), FIND THE NEXT
GATE OF GI ACCORDING TO THE CURRENT ORDERING P1
130 LGL = LGLIST(LEV)
LG = 0
140 LG = LG + 1
IF(LG.GT.LGL)GO TO 120 CAN NEVER HAPPEN
IF(GI.EQ.HLIST(LG,LEV))GO TO 120
GO TO 140
C**** CONSIDER INPUTS OF GATE GI
150 IF(GI.LE.N) GO TO 120
IF(GI.GT.NM.AND.GLEVEL(GI).LE.1)GO TO 120
LIP =LIPPED(GI)

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E1 14240  
E1 14250  
E1 14260  
E1 14270  
E1 14280  
E1 14290  
E1 14300  
E1 14310  
E1 14320  
E1 14330  
E1 14340  
E1 14350  
E1 14360  
E1 14370  
E1 14380  
E1 14390  
E1 14400  
E1 14410  
E1 14420  
E1 14430  
E1 14440  
E1 14450  
E1 14460  
E1 14470  
E1 14480  
E1 14490  
E1 14500  
E1 14510  
E1 14520  
E1 14530  
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E1 14590  
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E1 14700  
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E1 14720  
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E1 14750  
E1 14760  
E1 14770  
E1 14780  
E1 14790  
E1 14800  
E1 14810  
E1 14820

C**** INITIALIZE ARRAY (VECTOR) SUMP AND SUMS2	E1 14830
DO 160 TH=1,N2	E1 14840
SUMS2(TH) = 0	E1 14850
160 SUMP(TH) = 0	E1 14860
C**** SUM-UP ALL INPUTS OF GI	E1 14870
DO 180 LP =1,LIP	E1 14880
GJ = IPRED(LP,GI)	E1 14890
BSGJ = (GJ-1)*N2	E1 14900
DO 170 TH=1,N2	E1 14910
170 SUMP(TH) = SUMP(TH) + P\$(1,BSGJ + TH)	E1 14920
180 CONTINUE	E1 14930
C**** LIST 0,0-ERROR,1,1-PROP COMPONENTS OF CSPFE OF GI	E1 14940
NOC = 0	E1 14950
N01 = 0	E1 14960
N00E = 0	E1 14970
N01E = 0	E1 14980
DO 230 TH=1,N2	E1 14990
IF(GSMALL(GI,TH))190,230,210	E1 15000
190 IF(GSMALL(GI,TH).LE.-1070)GO TO 200	E1 15010
NCO = NCO + 1	E1 15020
IDX0(NCO) = TH	E1 15030
GO TO 230	E1 15040
200 N00E = N00E + 1	E1 15050
IDX0E(N00E) = TH	E1 15060
GO TO 230	E1 15070
210 IF(GSMALL(GI,TH).GE.1000)GO TO 220	E1 15080
N01 = N01 + 1	E1 15090
IDX1(N01) = TH	E1 15100
GO TO 230	E1 15110
220 N01E = N01E + 1	E1 15120
IDX1E(N01E) = TH	E1 15130
230 CONTINUE	E1 15140
C**** REMOVE REDUNDANT CONNECTIONS BY CALLING FORC ****	E1 15150
LIP = LIPRED(GI)	E1 15160
KEYA = 0	E1 15170
KEYR = 0	E1 15180
DO 350 LP=1,LIP	E1 15190
GJ = IPRED(LP,GI)	E1 15200
CALL FORC(GJ)	E1 15210
350 CONTINUE	E1 15220
C**** COMPARE SUMP WITH THE OUTPUT OF GI TO CHECK WHETHER ALL ERRORS	E1 15230
ARE CORRECTED OR NOT ****	E1 15240
IF(N00E.EQ.0) GO TO 380	E1 15250
N0E=0	E1 15260
DO 370 IDX=1,N00E	E1 15270
IF(SUMP(IDXOE(IDX)).GE.1) GO TO 360	E1 15280
N01=N01+1	E1 15290
IDX1(N01)=IDXOE(IDX)	E1 15300
GO TO 370	E1 15310
360 N0E=N0E+1	E1 15320
IDXOE(N0E)=IDXOE(IDX)	E1 15330
370 CONTINUE	E1 15340
N00E=N0E	E1 15350
C**** ALL ERRORS IN GI ARE CORRECTED ****	E1 15360
380 IF(N00E.EQ.0.AND.N01E.EQ.0) GO TO 1400	E1 15370
C**** CALCULATION OF EFFECTIVELY CONNECTABLE FUNCTIONS FOR GI	E1 15380
C**** ACCESS EACH GATE IN THE NETWORK	E1 15390
400 CNDDT = 0	E1 15400
DO 490 GR =1,NP	E1 15410
IF(GLEVEL(GR).EQ.1.AND.GR.GT.NM)GO TO 490	E1 15420
IF(INC\$MX(GR,GI).GE.1.OR.GI.EQ.GR.OR.SUC\$MX(GI,GR).GE.1)GOTO490	E1 15430

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C**** IF GR IS CONNECTED TO ALL SUCCESSORS OF GI, IT IS NOT A CANDIDATE E1 15440
      LIS = LISUCC(GI) E1 15450
      IF(LIS.EQ.0) GO TO 415 E1 15460
      DO 410 LI=1,LIS E1 15470
        IF(INC$MX(GR,ISUCC(LI,GI)).LE.0) GO TO 415 E1 15480
      410 CONTINUE E1 15490
      GO TO 490 E1 15500
      415 CONTINUE E1 15510
C**** CHECK THE POSSIBLE OUTPUTS OF THAT GATE **** E1 15520
      PA = PPOTAB(GR) E1 15530
      IF(PA.LE.0) GOTO 490 E1 15540
      PB = LPOTAB(GR) E1 15550
      DO 485 PTR=PA,PB E1 15560
        LTH = POTAB(PTR,$LTH) E1 15570
        IF(LTH.LE.0) GO TO 425 E1 15580
        DO 420 LT=1,LTH E1 15590
          IF(SUC$MX(GI,POTAB(PTR,$LTH+LT)).GE.1) GO TO 485 E1 15600
          IF(INC$MX(POTAB(PTR,$LTH+LT),GI).GE.1) GO TO 485 E1 15610
        420 CONTINUE E1 15620
      425 CONTINUE E1 15630
C**** CHECK THE EFFECTIVELY CONNECTABILITY **** E1 15640
      425 IF(N01.EQ.0) GO TO 435 E1 15650
      DO 430 IDX=1,N01 E1 15660
        IF(POTAB(PTR,IDX1(IDX)).NE.0) GO TO 485 E1 15670
      430 CONTINUE E1 15680
C**** POTAB(PTR) IS EFFECTIVELY CONNECTABLE TO GI **** E1 15690
C**** CALCULATE THE PREFERENCE WEIGHT **** E1 15700
      435 PW = 0 E1 15710
      IF(N00.EQ.0) GO TO 445 E1 15720
      DO 440 IDX=1,N00 E1 15730
        IF(POTAB(PTR,IDX0(IDX)).GE.1) PW = PW + 1 E1 15740
      440 CONTINUE E1 15750
      445 IF(N01E.EQ.0) GO TO 460 E1 15760
      DO 450 IDX=1,N01E E1 15770
        IF(POTAB(PTR,IDX1E(IDX)).GE.1) PW = PW + 1 E1 15780
      450 CONTINUE E1 15790
C**** SORT CANDIDATES IN ORDER ACCORDING TO PW E1 15800
      460 IF(CNDDT.EQ.0) GO TO 475 E1 15810
      DO 470 CR=1,CNDDT E1 15820
        CND = CNDDT-CR+1 E1 15830
        IF(PW.LE.POTAB(LISCND(CND),$PW)) GO TO 480 E1 15840
        LISCND(CND+1) = LISCND(CND) E1 15850
      470 CONTINUE E1 15860
      475 CND = 0 E1 15870
      480 LISCND(CND+1) = PTR E1 15880
      POTAB(PTR,$PW) = PW E1 15890
      CNDDT = CNDDT + 1 E1 15900
      485 CONTINUE E1 15910
      490 CONTINUE E1 15920
      IF(CNDDT.EQ.0) GO TO 1490 E1 15930
C**** CLASSIFY CANDIDATES INTO RI,DI,BIO,AND DIO **** E1 15940
      NODI = 0 E1 15950
      NOBI = 0 E1 15960
      NODIO = 0 E1 15970
      NOBIO = 0 E1 15980
      DO 650 NC =1, CNDDT E1 15990
        NOONEE = 0 E1 16000
        PTR = LISCND(NC) E1 16010
        IF(N00E.EQ.0) GO TO 630 E1 16020
        DO 610 NO=1,N00E E1 16030
          IF(POTAB(PTR,IDX0E(NO)).LE.0) GO TO 610 E1 16040
          NOONEE = NOONEE + 1

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610	CONTINUE	E1 16050
	IF(NOCNEE.EQ.0)GO TO 630	E1 16060
	POTAB(PTR,\$NCE) = NOCNEE	E1 16070
	IF(POTAB(PTR,\$GT).GT.N)GO TO 620	E1 16080
	NODI = NODI + 1	E1 16090
	DI(NODI) = PTR	E1 16100
	GO TO 650	E1 16110
C****	PJT PTR INTO TABLE BI AND SORT IT ACCORDING TO ORDER Q1(NOCNEE)	E1 16120
620	CONTINUE	E1 16130
	IF(NOB1.EQ.0) GO TO 624	E1 16140
	DO 623 NO=1,NB1	E1 16150
	NOB = NOB1 - NO + 1	E1 16160
	IF(POTAB( BI(NOB),\$NCE).LE.NOCNEE) GO TO 627	E1 16170
	BI(NOB+ 1) = BI(NOB)	E1 16180
623	CONTINUE	E1 16190
624	NB1 = 0	E1 16200
627	BI(NOB+1) = PTR	E1 16210
	NB1 = NB1 + 1	E1 16220
	GO TO 650	E1 16230
630	POTAB(PTR,\$NCE) = 0	E1 16240
	IF(POTAB(PTR,\$GT).GT.N)GO TO 640	E1 16250
	NODIO = NODIO + 1	E1 16260
	DIO(NODIO) = PTR	E1 16270
	GO TO 650	E1 16280
640	NBIO = NBIO + 1	E1 16290
	BIO(NBIO) = PTR	E1 16300
650	CONTINUE	E1 16310
C****	CALCULATE SET S2 *****	E1 16320
	IF(NBIO.EQ.0)GO TO 770	E1 16330
	DO 760 NO=1,NBIO	E1 16340
760	SETS2(NO) = BIO(NO)	E1 16350
770	IF(NODIO.EQ.0)GO TO 790	E1 16360
	DO 780 NO=1,NODIO	E1 16370
780	SETS2(NBIO+NO) = DIO(NO)	E1 16380
790	IF(NB1.EQ.0)GO TO 810	E1 16390
	DO 800 NO=1,NB1	E1 16400
800	SETS2(NBIO+NODIO+NO) = BI(NO)	E1 16410
C****	CALL PROCEDURES TO REPLACE EXTERNAL VARIABLES *****	E1 16420
810	NOS2 = NBIO + NODIO + NB1	E1 16430
	LMTS2=NOS2	E1 16440
	IF(NOS2.EQ.0) GO TO 1400	E1 16450
	IF(NDOE.EQ.0) GO TO 1210	E1 16460
C****	CALCULATE SUMS2 *****	E1 16470
	DO 832 NO=1,NOS2	E1 16480
	PTR = SETS2(NO)	E1 16490
	DO 830 TH=1,N2	E1 16500
830	SUMS2(TH) = SUMS2(TH) + POTAB(PTR,TH)	E1 16510
832	CONTINUE	E1 16520
C****	REPLACEMENT OF EXTERNAL VARIABLES *****	E1 16530
C****	CALCULATE SET S *****	E1 16540
	NOS = 0	E1 16550
	DO 750 LP = 1,LIP	E1 16560
	GP = IPRED(LP,GI)	E1 16570
	IF(INC\$MX(GP,GI).LE.0.OR.GP.GT.N)GO TO 750	E1 16580
	BSGP = (GP-1)*N2	E1 16590
	DO 700 NO=1,NDOE	E1 16600
	IF(P\$(1,BSGP+IDXOE(NO)).GE.1) GO TO 710	E1 16610
700	CONTINUE	E1 16620
	GO TO 750	E1 16630
710	NOS = NOS + 1	E1 16640
	SETS(NOS) = GP	E1 16650



750	CONTINUE	E1 16660
	STS = 1	E1 16670
	IF(NDS.LE.0)GO TO 835	E1 16680
	NOT1SV=0	E1 16690
	NDS1SV=0	E1 16700
	CALL CALS1	E1 16710
	CALL RPLOC	E1 16720
C****	REPLACEMENT OF INTERNAL FUNCTIONS WITH 1 ERROR *****	E1 16730
835	BRPLC = 0	E1 16740
	NDS1=0	E1 16750
	NOT1=0	E1 16760
	NDS = 0	E1 16770
	CALL ORDRQ2	E1 16780
	P2 = NRPLC(2)	E1 16790
	IF(P2.EQ.0) GO TO 995	E1 16800
	DO 990 PR=1,P2	E1 16810
	GP = RPLC(2,PR)	E1 16820
	NDS = NDS + 1	E1 16830
	SETS(NDS) = GP	E1 16840
	RSGP = (GP - 1)*N2	E1 16850
C****	CALCULATE THE ERROR POSITIONS OF WEIGHT 1 IN GP *****	E1 16860
	NCK = 0	E1 16870
	DO 840 ND=1,NDOE	E1 16880
	TH = IDXOE(ND)	E1 16890
	IF(SUMP(TH).NE.1.DR.P\$(1,RSGP+TH).NE.1)GO TO 840	E1 16900
	NCK = NCK + 1	E1 16910
	IDXK(NCK) = TH	E1 16920
840	CONTINUE	E1 16930
C****	CALCULATE SETS4 *****	E1 16940
	DO 940 NK=1,NCK	E1 16950
	TH = IDXK(NK)	E1 16960
C****	IF SETS2(*)=1000+PTR, IT IS NOT IN SET S4 *****	E1 16970
	DO 870 SR=1,NDS2	E1 16980
	IF(SETS2(SR).GT.1000) GO TO 870	E1 16990
	IF(POTAB(SETS2(SR),TH).EQ.0) GO TO 861	E1 17000
C****	UPDATE SUMS2 *****	E1 17010
	DO 860 TT=1,N2	E1 17020
860	SUMS2(TT) = SUMS2(TT) - POTAB(SETS2(SR),TT)	E1 17030
	SETS2(SR) = SETS2(SR) + 1000	E1 17040
	GO TO 870	E1 17050
C****	TEMPORARILY PUT SETS2(SR) INTO SETS4. PROHIBIT FUNCTIONS WHICH	E1 17060
C	HAS SAME OUTPUT GATE AS SETS2(SR)	E1 17070
861	CONTINUE	E1 17080
	IF(BRPLC.EQ.1) GO TO 870	E1 17090
	GT=POTAB(SETS2(SR),\$GT)	E1 17100
	DO 865 SRR=1,NDS2	E1 17110
	IF(SRR.EQ.SR) GO TO 865	E1 17120
	IF(SETS2(SRR).GT.1000) GO TO 865	E1 17130
	IF(POTAB(SETS2(SRR),\$GT).EQ.GT) GO TO 8617	E1 17140
	IF(POTAB(SETS2(SRR),\$LTH).EQ.0) GO TO 865	E1 17150
	MRUN=PCTAB(SETS2(SRR),\$LTH)	E1 17160
	DO 8615 RUN=1,MRUN	E1 17170
	IF(POTAB(SETS2(SRR),\$LTH+RUN).EQ.GT) GO TO 8617	E1 17180
8615	CONTINUE	E1 17190
	GO TO 865	E1 17200
8617	CONTINUE	E1 17210
C****	UPDATE SUMS2(*)	E1 17220
	DO 863 TT=1,N2	E1 17230
863	SUMS2(TT)=SUMS2(TT)-POTAB(SETS2(SRR),TT)	E1 17240
	SETS2(SRR)=SETS2(SRR)+1000	E1 17250
865	CONTINUE	E1 17260

870	CONTINUE	E1 17270
C****	CHECK WHETHER OR NOT ELEMENTS IN SETS4 COVER ALL ESSENTIAL	E1 17280
C	ONES OF GP *****	E1 17290
	DO 900 NO=1,N00	E1 17300
	TT=IDX0(NO)	E1 17310
	IF(P\$(1,BSGP+TT).NE.1.OR.SUMP(TT).NE.1) GO TO 900	E1 17320
	IF(SUMS2( TT ).EQ.0)GO TO 930	E1 17330
900	CONTINUE	E1 17340
C****	S4 COVERS ALL ESSENTIAL ONES *****	E1 17350
C	UPDATE SETS2,SETS1, AND T1	E1 17360
	BRPLC = 1	E1 17370
	DO 910 SR=1,NOS2	E1 17380
	IF(SETS2(SR).GT.1000.AND.SETS2(SR).LT.2000)	E1 17390
1	SETS2(SR) = SETS2(SR) + 1000	E1 17400
910	CONTINUE	E1 17410
	NOS1 = NOS1 + 1	E1 17420
	SETS1(NOS1) = GP	E1 17430
	SETS(NOS) = 2000 + GP	E1 17440
	DO 915 NO=1,N00	E1 17450
	TT=IDX0(NO)	E1 17460
	IF(P\$(1,BSGP+TT).NE.1.OR.SUMP(TT).NE.1) GO TO 915	E1 17470
	NOT1 = NOT1 + 1	E1 17480
	SETT1(NOT1) = TT	E1 17490
915	CONTINUE	E1 17500
C****	UPDATE SJMP *****	E1 17510
	DO 920 TH=1,N2	E1 17520
	SUMP(TH)=SUMP(TH) - P\$(1,BSGP+TH)	E1 17530
920	CONTINUE	E1 17540
	GO TO 990	E1 17550
C****	RESET SETS4 ****	E1 17560
930	DO 935 NO=1,NOS2	E1 17570
	IF(SETS2(NO).LT.1000.OR.SETS2(NO).GT.2000)GO TO 935	E1 17580
	SETS2(NO) = SETS2(NO) - 1000	E1 17590
C****	UPDATE SETS2 (THIS ELEMENT BECOMES ACTIVE AGAIN)	E1 17600
	DO 932 TH=1,N2	E1 17610
932	SUMS2(TH) = SUMS2(TH) + POTAB(SETS2(NO),TH)	E1 17620
935	CONTINUE	E1 17630
940	CONTINUE	E1 17640
990	CONTINUE	E1 17650
C****	PUT RPLC(1,*) INTO SET S *****	E1 17660
995	STS1 = NOS + 1	E1 17670
	P2 = NRPLC(1)	E1 17680
	IF(P2.EQ.0)GO TO 1005	E1 17690
	DO 1000 PR=1,P2	E1 17700
	NOS = NOS + 1	E1 17710
	SETS(NOS) = RPLC(1,PR)	E1 17720
1000	CONTINUE	E1 17730
1005	IF(BRPLC.EQ.0)GO TO 1010	E1 17740
	NOT1SV=NOT1	E1 17750
	NOS1SV=NOS1	E1 17760
	CALL CALS1	E1 17770
	CALL RPLCF	E1 17780
C****	REPLACEMENT BY BIO AND DIO *****	E1 17790
1010	NOS2 = NOBIO + NODIO	E1 17800
C****	UPDATE SUMS2(*) TO CONTAIN THE CURRENT ELEMENTS IN SETS2 ONLY *****	E1 17810
	IF(NOBI.EQ.0) GO TO 1060	E1 17820
	S1 = NOBIO + NODIO + 1	E1 17830
	S2 = S1 + NOBI - 1	E1 17840
	DO 1050 NO=S1,S2	E1 17850
	PTR = SETS2(NO)	E1 17860
	IF(PTR.GT.1000) GO TO 1050	E1 17870



DO 1040 TH=1,N2	E1 17880
1040 SUMS2(TH) = SUMS2(TH) - POTAB(PTR,TH)	E1 17890
1050 CONTINUE	E1 17900
1060 CONTINUE	E1 17910
IF(NDS2.EQ.0) GO TO 1100	E1 17920
STS = STS1	E1 17930
IF(STS.GT.NDS)GO TO 1100	E1 17940
NDS1SV=0	E1 17950
NOT1SV=0	E1 17960
CALL CALS1	E1 17970
CALL SPLCF	E1 17980
C**** COMPENSATION OF 1 ERRORS OF CSPEE OF GI	E1 17990
1100 CONTINUE	E1 18000
C**** SUM-UP ADDED FUNCTIONS *****	E1 18010
C MODIFY SUMS2 TO CONTAIN ONLY ADDED CONNECTIONS	E1 18020
IF(NDS2.EQ.0) GO TO 1175	E1 18030
DO 1170 ND=1,NDS2	E1 18040
PTR = SETS2(ND)	E1 18050
IF(PTR.GT.1000)GO TO 1170	E1 18060
DO 1150 TH=1,N2	E1 18070
SUMS2(TH) = SUMS2(TH) - POTAB(PTR,TH)	E1 18080
1150 CONTINUE	E1 18090
1170 CONTINUE	E1 18100
1175 CONTINUE	E1 18110
NDS2 = NORID + NODIO + NORI	E1 18120
C**** LIST UNCOVERED 1-ERROR COMPONENTS *****	E1 18130
IF(ND1F.EQ.0)GO TO 1400	E1 18140
NOT0 = 0	E1 18150
DO 1200 ND=1,ND1F	E1 18160
TH = IDX1E(ND)	E1 18170
IF(SUMS2(TH).GE.1)GO TO 1180	E1 18180
NOT0 = NOT0 +1	E1 18190
GO TO 1200	E1 18200
C**** TH IS COVERED ALREADY (IDX1E(*)>1000 : COVERED)	E1 18210
1180 IDX1E(ND) = 1000 + TH	E1 18220
1200 CONTINUE	E1 18230
IF(NOT0.EQ.0)GO TO 1400	E1 18240
GO TO 1220	E1 18250
1210 NOT0 = ND1F	E1 18260
C**** RESTRICTIONS ON S2 AS STATED IN ALGORITHM MAY BE INSERTED HERE	E1 18270
C**** CHECK ACTIVE FUNCTIONS IN S2	E1 18280
1220 CONTINUE	E1 18290
DO 1300 ND=1,NDS2	E1 18300
PTR=SETS2(ND)	E1 18310
IF(PTR.GT.2000)GO TO 1300	E1 18320
C**** CHECK WHETHER THIS FUNCTION CAN COMPENSATE SOME ONE-ERRORS .	E1 18330
C OR NOT	E1 18340
NOT00 = NOT0	E1 18350
DO 1230 NDE=1,ND1E	E1 18360
IF(IDX1E(NDE).GT.1000)GO TO 1230	E1 18370
IF(POTAB(PTR,IDX1E(NDE)).NE.1)GO TO 1230	E1 18380
NOT0 = NOT0 - 1	E1 18390
IDX1E(NDE) = 1000 + IDX1E(NDE)	E1 18400
1230 CONTINUE	E1 18410
IF(NOT0.EQ.NOT00)GO TO 1300	E1 18420
DO 1235 TH=1,N2	E1 18430
1235 SUMP(TH)=SUMP(TH)+POTAB(SETS2(ND),TH)	E1 18440
SETS2(ND) = 5000 + PTR	E1 18450
CALL CONECT(PTR)	E1 18460
C**** PROHIBIT FUNCTIONS WHICH HAS SAME OUTPUT GATE AS PTR ****	E1 18470
IF(ND.EQ.NDS2) GO TO 1300	E1 18480

SR1=NO+1	E1 18490
GT=POTAB(PTP,\$GT)	E1 18500
DO 1250 SR=SR1,NOS2	E1 18510
IF(SETS2(SR).GT.2000) GO TO 1250	E1 18520
IF(POTAB(SETS2(SR),\$GT).EQ.GT) GO TO 1245	E1 18530
IF(POTAB(SETS2(SR),\$LTH).EQ.0) GO TO 1250	E1 18540
MRUN=POTAB(SETS2(SR),\$LTH)	E1 18550
DO 1240 RUN=1,MRUN	E1 18560
IF(POTAB(SETS2(SR),\$LTH+RUN).EQ.GT) GO TO 1245	E1 18570
1240 CONTINUE	E1 18580
GO TO 1250	E1 18590
1245 SETS2(SR)=SETS2(SR)+2000	E1 18600
1250 CONTINUE	E1 18610
1300 CONTINUE	E1 18620
1400 CONTINUE	E1 18630
C**** ADDING EXTERNAL VARIABLES TO GI ****	E1 18640
DO 1480 GP=1,N	E1 18650
IF(INC\$MX(GP,GI).GE.1) GO TO 1480	E1 18660
BSGP=(GP-1)*N2	E1 18670
C**** CHECK CONNECTABILITY ****	E1 18680
DO 1410 IDX=1,N01	E1 18690
IF(P\$(1,BSGP+IDX1(IDX)).GE.1) GO TO 1480	E1 18700
1410 CONTINUE	E1 18710
IF(N00E.EQ.0) GO TO 1430	E1 18720
DO 1420 IDX=1,N00E	E1 18730
TH=IDX0E(IDX)	E1 18740
IF(P\$(1,BSGP+TH).GE.1.AND.SUMP(TH).LE.1) GO TO 1480	E1 18750
1420 CONTINUE	E1 18760
C**** CHECK WHETHER OR NOT IT COVERS 0 OR 1-ERROR COMPONENTS ****	E1 18770
1430 CONTINUE	E1 18780
DO 1440 IDX=1,N00	E1 18790
IF(P\$(1,BSGP+IDX0(IDX)).GE.1) GO TO 1460	E1 18800
1440 CONTINUE	E1 18810
IF(N01E.EQ.0) GO TO 1480	E1 18820
DO 1450 IDX=1,N01E	E1 18830
TH=IDX1E(IDX)	E1 18840
IF(TH.GT.1000) TH=TH-1000	E1 18850
IF(P\$(1,BSGP+TH).GE.1) GO TO 1460	E1 18860
1450 CONTINUE	E1 18870
GO TO 1480	E1 18880
C**** GP CAN BE CONNECTED TO GI ****	E1 18890
1460 S=S+1	E1 18900
RSCOVN(S)=100*GP+GI	E1 18910
INC\$MX(GP,GI)=1	E1 18920
KEYB=1	E1 18930
1480 CONTINUE	E1 18940
1490 IF(KEYA.EQ.0) GO TO 1500	E1 18950
C**** SOME ERRORS WERE COMPENSATED IN GI	E1 18960
RETURN2	E1 18970
C**** CALCULATION OF CSPFE FOR INPUTS	E1 18980
1500 CONTINUE	E1 18990
IF(KEYB.EQ.1)CALL SUBNET	E1 19000
IF(KEYB.EQ.1) CALL UNNECE	E1 19010
C**** PROPAGATE ONE, ONE-ERROR, AND ZERO-ERROR COMPONENTS ****	E1 19020
LIP = LIPRED(GI)	E1 19030
DO 1600 LP=1,LIP	E1 19040
GP = IPRED(LP,GI)	E1 19050
IF(GP.LE.4)GO TO 1600	E1 19060
BSGP = (GP-1)*N2	E1 19070
C**** FOR ONE COMPONENTS ****	E1 19080
IF(N01.EQ.0) GO TO 1530	E1 19090

	DO 1520 NO=1,NO1	E1 19100
	IF(GSMALL(GP,IDX1(NO)).EQ.0)GSMALL(GP,IDX1(NO)) = -100	E1 19110
	IF(GSMALL(GP,IDX1(NO)).LT.-1000.AND.GP.GT.NM)	E1 19120
1	GSMALL(GP,IDX1(NO))= -100	E1 19130
1520	CONTINUE	E1 19140
C****	FOR ONE-ERROR COMPONENTS *****	E1 19150
1530	IF(N01E.EQ.0) GO TO 1550	E1 19160
	DO 1540 NO=1,NO1E	E1 19170
	IF(GSMALL(GP,IDX1E(NO)).EQ.0)GSMALL(GP,IDX1E(NO)) = -1100	E1 19180
1540	CONTINUE	E1 19190
C****	FOR ZERO-ERROR COMPONENTS *****	E1 19200
1550	IF(N00E.EQ.0) GO TO 1600	E1 19210
	DO 1560 NO=1,N00E	E1 19220
	IF(P\$(1,BSGP+IDX0E(NO)).EQ.0)GO TO 1555	E1 19230
	IF(GSMALL(GP,IDX0E(NO)).EQ.0) GSMALL(GP,IDX0E(NO)) = 1001	E1 19240
	GO TO 1560	E1 19250
1555	IF(GSMALL(GP,IDX0E(NO)).EQ.0.OR.GSMALL(GP,IDX0E(NO)).LT.-1000)	E1 19260
1	GSMALL(GP,IDX0E(NO))=-100	E1 19270
1560	CONTINUE	E1 19280
1600	CONTINUE	E1 19290
C****	PROPAGATE ZERO COMPONENTS *****	E1 19300
	IF(N00.EQ.0) GO TO 1800	E1 19310
C*****	*****	E1 19320
C	CALCULATION OF ORDERING P4	E1 19330
C	SINCE NO CONNECTIONS ARE ADDED DURING CALCULATION FOR GATE GI,	E1 19340
C	SUMP(*) HAS THE INPUT SUM OF GATE GI	E1 19350
C****	CALCULATE NUMBER OF 1-ERRORS AND NUMBER OF ESSENTIAL 1-ERRORS IN	E1 19360
C	EACH INPUT OF GI	E1 19370
	DO 1610 LP=1,LIP	E1 19380
	GP=IPRED(LP,GI)	E1 19390
	IF(GP.LE.N) GO TO 1610	E1 19400
	N00NEE=0	E1 19410
	N01EES=0	E1 19420
	IF(N00E.EQ.0) GO TO 1606	E1 19430
	BSGP=(GP-1)*N2	E1 19440
	DO 1605 NO=1,N00E	E1 19450
	TH=IDX0E(NO)	E1 19460
	IF(P\$(1,BSGP+TH).GE.1) N00NEE=N00NEE+1	E1 19470
	IF(SUMP(TH).EQ.1.AND.P\$(1,BSGP+TH).EQ.1)N01EES=N01EES+1	E1 19480
1605	CONTINUE	E1 19490
1606	ORDPP4(GP)=1000000-N01EES*10000-N00NEE*100+LISJCC(GP)	E1 19500
1610	CONTINUE	E1 19510
C****	END OF CALCULATION OF ORDERING P4	E1 19520
	DO 1700 NO=1,N00	E1 19530
	TH = IDX0(NO)	E1 19540
	PODRMX = 0	E1 19550
	DO 1620 LP=1,LIP	E1 19560
	GP = IPRED(LP,GI)	E1 19570
	BSGP = (GP-1)*N2	E1 19580
	IF(P\$(1,BSGP+TH).EQ.0)GO TO 1620	E1 19590
C****	IF SOME GATES ARE ALREADY ASSIGNED ONE, DO NOTHING	E1 19600
	IF(GSMALL(GP,TH).EQ.1.OR.GP.LE.N) GO TO 1700	E1 19610
C****	COMPARE PRIORITY OF CURRENT GATE WITH PREVIOUS HIGHEST ONE	E1 19620
	IF(GP.LE.NM.AND.GSMALL(GP,TH).GT.1000)GO TO 1620	E1 19630
	PODR=ORDPP4(GP)	E1 19640
	IF(PODR.LE.PODRMX)GO TO 1620	E1 19650
	PODRMX = PODR	E1 19660
	PODRGT = GP	E1 19670
1620	CONTINUE	E1 19680
	IF(PODRMX.EQ.0)GO TO 1630	E1 19690
	GSMALL(PODRGT,TH) = 1	E1 19700

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GO TO 1700
C**** THIS ZERO IS COVERED ONLY BY OUTPUT GATES WHICH ARE ALREADY
C ASSIGNED ONE-ERROR. PROPAGATE ZERO ERROR TO PREDECESSORS WHICH
C HAVE ZERO COMPONENT
1630 DO 1650 LP=1,LIP
      GP = IPRED(LP,GI)
      BSGP = (GP-1)*N2
      IF(P$(1,BSGP+TH).EQ.1.OR.GP.LE.N)GO TO 1650
      IF(GSMALL(GP,TH).EQ.?)GSMALL(GP,TH) = -1100
1650 CONTINUE
1700 CONTINUE
1800 CONTINUE
GO TO 130
END

SUBROUTINE RPLCF
C**** CALCULATE A SUBSET OF SET S2 WHICH CAN REPLACE SET S1
C SET S2 IS LISTED ACCORDING TO ORDER Q1
C
C DEFINITIONS OF 'COMMON' VARIABLES CAN BE FOUND IN MAIN PROGRAM.
C
IMPLICIT INTEGER*4(A-T,V-Z,$), REAL(U)
COMMON NEPMAX
COMMON
  V , M , A , B
1 , R , N2 , N1 , NR
2 , NM , KFLAG , JFLAG , COST
3 , LEVM , NRN2 , NM1 , NN2
COMMON ISJCC(40,40) , LISUCC(40) , IPRED(40,40) , LIPRED(40)
1 , INC$MX(40,40) , SUC$MX(40,40) , P$(2,1280) , UNAME(40)
2 , GLEVEL(40) , LGLIST(40) , HLIST(40,40) , TIME
COMMON T , RTCONN(100) , S , RSCONN(100)
COMMON IFLAG , POINTA , ESS1S(40) , F$1(32)
1 , F$UB1 , INPTCV(32) , LISTC(40) , POINTC
2 , LISTL(40) , POINTL , ORIGIN(40) , IPATH(40)
3 , POINTR , VF$1(32) , VF$UB1 , GSMALL(40,32)
COMMON POTAB(200,42) , PPOTAB(40) , LPOTAB(40) , NRPLC(2)
1 , RPLC(2,40) , IDX0(32) , IDX0E(32) , IDX1(32)
2 , IDX1E(32) , SUMP(32) , SETT1(32) , NOT1
3 , SETS1(40) , NOS1 , SETS(40) , NOS
4 , STS , SUMS2(32) , SETS2(200) , NOS2
5 , LIP , NOOE , KEYA , KEYB
6 , NOO , NO1 , NO1E , $GT
7 , $LTH , $PW , $NOE , GI
COMMON NOT1SV , NOS1SV , LMTS2
6000 IF (NOS1.EQ.0) RETURN
DO 6200 NO=1,NOS2
  IF(SETS2(NO).GT.2000) GO TO 6200
  PTR=SETS2(NO)
C**** CHECK ESSENTIAL ONES IN S1
  KEYT2=0
DO 6100 NOT=1,NOT1
  IF(SETT1(NOT).GT.1000) GO TO 6100
  TH=SETT1(NOT)
  IF(POTAB(PTR,TH).LE.0)GO TO 6100
  KEYT2=1
  SETT1(NOT)=1000+TH
6100 CONTINUE
C**** IF T2 IS EMPTY THIS FUNCTION IS NOT IN SET S3
  IF(KEYT2.LE.0) GO TO 6200
C**** PROHIBIT FNCTIONS IN S2 WHICH USE SAME OUTPUT GATE AS PTR *****

```



IF(NOT1SV.GT.0) GOTO 6180	E1 20300
NOT1SV>0 => BRPLC>0 => SETS2 IS COMPATIBLE	E1 20310
SW=0	E1 20320
GT=PTRAB(PTR,\$GT)	E1 20330
DO 6150 IDX=1,LMTS2	E1 20340
IF(IDX.EQ.NO) GO TO 6150	E1 20350
IF(SETS2(IDX).GT.2000) GO TO 6150	E1 20360
IF(POTAB(SETS2(IDX),\$GT).EQ.GT) GO TO 6140	E1 20370
IF(POTAB(SETS2(IDX),\$LTH).EQ.0) GO TO 6150	E1 20380
MPJIN=PTRAB(SETS2(IDX),\$LTH)	E1 20390
DO 6120 RUN=1,MRUN	E1 20400
IF(POTAB(SETS2(IDX),\$LTH+RUN).EQ.GT) GO TO 6140	E1 20410
6120 CONTINUE	E1 20420
GO TO 6150	E1 20430
6140 CONTINUE	E1 20440
IF(IDX.GT.NOS2) GO TO 6145	E1 20450
SW=1	E1 20460
C**** UPDATE SUMS2 ****	E1 20470
DO 6130 TH=1,N2	E1 20480
6130 SUMS2(TH)=SUMS2(TH)-PTRAB(SETS2(IDX),TH)	E1 20490
6145 SETS2(IDX)=3000+SETS2(IDX)	E1 20500
6150 CONTINUE	E1 20510
IF(SW.EQ.0) GO TO 6180	E1 20520
C**** CHECK WHETHER S2 STILL COVER ALL ESSENTIAL ONES OR NOT ****	E1 20530
C IF YES CONTINUE THIS PROCEDURE, OTHERWISE RECALCULATE S1 AND	E1 20540
C REPEAT THIS PROCEDURE	E1 20550
DO 6160 NOT1=1,NOT1	E1 20560
IF(SETT1(NOT).GT.1000) GO TO 6160	E1 20570
IF(SUMS2(SETT1(NOT)).LE.0) GO TO 6190	E1 20580
6160 CONTINUE	E1 20590
6180 SETS2(NO)=4000+PTR	E1 20600
GO TO 6200	E1 20610
6190 CONTINUE	E1 20620
DO 6195 TH=1,N2	E1 20630
6195 SUMS2(TH)=SUMS2(TH)-PTRAB(PTR,TH)	E1 20640
SETS2(NO)=PTR+2000	E1 20650
GO TO 6500	E1 20660
6200 CONTINUE	E1 20670
C**** REPLACE FUNCTIONS IN S1 BY FUNCTIONS IN S2 ****	E1 20680
DO 6300 NO=1,LMTS2	E1 20690
PTR=SETS2(NO)	E1 20700
IF(PTR.GT.5000) GO TO 6300	E1 20710
IF(PTR.GT.4000) GO TO 6210	E1 20720
IF(PTR.GT.3000) GO TO 6290	E1 20730
GO TO 6300	E1 20740
C**** ADDING POTAB(PTR,*) TO SUMP ****	E1 20750
6210 CONTINUE	E1 20760
DO 6220 TH=1,N2	E1 20770
6220 SUMP(TH)=SUMP(TH)+POTAB(PTR-4000,TH)	E1 20780
C**** RECORD THIS FUNCTION HAS BEEN CONNECTED (SETS2(*)>5000) ****	E1 20790
SETS2(NO)=1000+PTR	E1 20800
C**** CONNECT THIS FUNCTION TO GI AND MAKE OTHER CONNECTIONS NECESSARY	E1 20810
C FOR REALIZING THIS FUNCTION *****	E1 20820
CALL CONECT(PTR-4000)	E1 20830
GO TO 6300	E1 20840
C**** THIS FUNCTION CAN NO LONGER BE USED TO REPLACE OTHER FUNCTIONS ***	E1 20850
6290 SETS2(NO)=PTR-1000	E1 20860
6300 CONTINUE	E1 20870
C**** DISCONNECTION PREDECESSORS OF GI IN SET S1	E1 20880
DO 6400 NO=1,NOS1	E1 20890
GP=SETS1(NO)	E1 20900

	INC\$MX(GP,GI)=0	E1 20910
	T=T+1	E1 20920
	RTCONN(T)=100*GP+GI	E1 20930
6400	CONTINUE	E1 20940
C****	MAKE PERMANENT CHANGES IN SET S (S>2000 : REMOVED )	E1 20950
C****	(2000>S>1000 : TEMPORARILY REMOVED )	E1 20960
	DO 6450 NO=1,NOS	E1 20970
	IF(SETS(NO).LT.2000.AND.SETS(NO).GT.1000) SETS(NO)=SETS(NO)+1000	E1 20980
6450	CONTINUE	E1 20990
	RETURN	E1 21000
C****	SET S1 IS NOT REPLACABLE BY S2. RECALCULATE SET S1	E1 21010
C	PROHIBIT THE FUNCTIONS WHICH CAUSED THIS SITUATION ****	E1 21020
6500	CONTINUE	E1 21030
	DO 6550 NO=1,LMTS2	E1 21040
	PTR=SETS2(NO)	E1 21050
	IF(PTR.GT.5000) GO TO 6550	E1 21060
	IF(PTR.GT.4000) GO TO 6510	E1 21070
	IF(PTR.GT.3000) GO TO 6520	E1 21080
	GO TO 6550	E1 21090
C****	MAKE THIS ELEMENT ACTIVE *****	E1 21100
6510	SETS2(NO)=PTR-4000	E1 21110
	GO TO 6550	E1 21120
6520	SETS2(NO)=PTR-3000	E1 21130
C****	MAKE THIS ELEMENT ACTIVE ****	E1 21140
	IF(NO.GT.NOS2)GO TO 6550	E1 21150
	DO 6530 TH=1,N2	E1 21160
6530	SUMS2(TH)=SUMS2(TH)+POTAB(SETS2(NO),TH)	E1 21170
6550	CONTINUE	E1 21180
C****	ADD BACK ELEMENTS OF SET S1 TO SUMP	E1 21190
	DO 6600 NO=1,NOS1	E1 21200
	GP=SETS1(NO)	E1 21210
	BSP=(GP-1)*N2	E1 21220
	DO 6580 TH=1,N2	E1 21230
6580	SUMP(TH)=SUMP(TH)+P\$(1,BSP+TH)	E1 21240
6600	CONTINUE	E1 21250
C****	RECONSTRUCT SET S *****	E1 21260
	DO 6650 NO=1,NOS	E1 21270
	IF(SETS(NO).GT.2000) GO TO 6650	E1 21280
	IF(SETS(NO).LT.1000) GO TO 6650	E1 21290
	SETS(NO)=SETS(NO)-1000	E1 21300
6650	CONTINUE	E1 21310
	CALL CALS1	E1 21320
	GO TO 6000	E1 21330
	END	E1 21340
SUBROUTINE SUBNET		E1 21350
C		E1 21360
C	DEFINITIONS OF 'COMMON' VARIABLES CAN BE FOUND IN MAIN PROGRAM	E1 21370
C		E1 21380
	IMPLICIT INTEGER*4(A-T,V-Z,\$), REAL(U)	E1 21390
	COMMON NEPMAX	E1 21400
	COMMON N , M , A , B	E1 21410
1	, R , N2 , N1 , NR	E1 21420
2	, NM , KFLAG , JFLAG , COST	E1 21430
3	, LEVM , NRN2 , NM1 , NN2	E1 21440
	COMMON ISJCC(40,40) , LISUCC(40) , IPRED(40,40) , LIPRED(40)	E1 21450
1	, INC\$MX(40,40) , SUC\$MX(40,40) , P\$(2,1280) , UNAME(40)	E1 21460
2	, GLEVEL(40) , LGLIST(40) , HLIST(40,40) , TIME	E1 21470
	COMMON T , RTCONN(100) , S , RSCONN(100)	E1 21480
	COMMON IFLAG , POINTA , ESSIS(40) , F\$1(32)	E1 21490

1	,F\$UR1	,INPTCV(32)	,LISTC(40)	,POINTC	E1 21500
2	,LISTL(40)	,POINTL	,ORIGIN(40)	,IPATH(40)	E1 21510
3	,POINTR	,VF\$1(32)	,VF\$UB1	,GSMALL(40,32)	E1 21520
COMMON	POTAB(200,42),PPOTAB(40)		,LPOTAB(40)	,NRPLC(2)	E1 21530
1	,RPLC(2,40)	,IDX0(32)	,IDX0E(32)	,IDX1(32)	E1 21540
2	,IDX1F(32)	,SUMP(32)	,SETT1(32)	,NOT1	E1 21550
3	,SET\$1(40)	,NOS1	,SETS(40)	,NOS	E1 21560
4	,STS	,SUMS2(32)	,SETS2(200)	,NOS2	E1 21570
5	,LIP	,NO1E	,KEYA	,KEYB	E1 21580
6	,NDO	,NO1	,NO1E	,\$GT	E1 21590
7	,\$LTH	,\$PW	,\$NOE	,G\$\$\$\$\$	E1 21600
COMMON	NOT1SV	,NOS1SV	,LMTS2		E1 21610
	DIMENSION X(40),LX(40,2),OUTO(40)				E1 21620
C	ENTRY PRESUC				E1 21630
1	CONTINUE				E1 21640
	DO 10 GI=1,NR				E1 21650
	LS=0				E1 21660
	LP=0				E1 21670
	DO 5 GJ=1,NR				E1 21680
	IF(INC\$MX(GI,GJ).EQ.0) GO TO 3				E1 21690
	LS=LS+1				E1 21700
	ISUCC(LS,GI)=GJ				E1 21710
	GO TO 5				E1 21720
3	IF(INC\$MX(GJ,GI).EQ.0) GO TO 5				E1 21730
	LP=LP+1				E1 21740
	IPRED(LP,GI)=GJ				E1 21750
5	CONTINUE				E1 21760
	LISUCC(GI)=LS				E1 21770
	LIPRED(GI)=LP				E1 21780
10	CONTINUE				E1 21790
C					E1 21800
	ENTRY SUCCS				E1 21810
	DO 21 GI=1,NR				E1 21820
	DO 21 GJ=1,NR				E1 21830
	SUC\$MX(GI,GJ)=0				E1 21840
21	CONTINUE				E1 21850
	DO 30 GJ=N1,NR				E1 21860
	DO 22 GS=1,NR				E1 21870
	X(GS)=0				E1 21880
22	CONTINUE				E1 21890
	X(GJ)=1				E1 21900
	LO=1				E1 21910
	LX(1,1)=GJ				E1 21920
	V=1				E1 21930
23	CONTINUE				E1 21940
	V=1-V				E1 21950
	SW0=1+V				E1 21960
	SW1=2-V				E1 21970
	L1=0				E1 21980
	DO 28 LL=1,LO				E1 21990
	GM=LX(LL,SW0)				E1 22000
	LIP=LIPRED(GM)				E1 22010
	IF(LIP.EQ.0) GO TO 28				E1 22020
	DO 26 LP=1,LIP				E1 22030
	GP=IPRED(LP,GM)				E1 22040
	IF(X(GP).GT.0) GO TO 25				E1 22050
	SUC\$MX(GP,GJ)=1				E1 22060
	L1=L1+1				E1 22070
	LX(L1,SW1)=GP				E1 22080
	X(GP)=1				E1 22090
26	CONTINUE				E1 22100



28	CONTINUE	E1 22110
	IF(L1.EQ.0) GO TO 30	E1 22120
	L1=L1	E1 22130
	GO TO 23	E1 22140
30	CONTINUE	E1 22150
	ENTRY LEVEL	E1 22160
	DO 40 GJ=1,NR	E1 22170
	OUTO(GJ)=LISUCC(GJ)	E1 22180
	GLEVEL(GJ)=-1	E1 22190
40	CONTINUE	E1 22200
	LEV=0	E1 22210
45	LEV=LEV+1	E1 22220
	G=0	E1 22230
	DO 50 GJ=1,NR	E1 22240
	IF(OUTO(GJ).GT.0 .OR. GLEVEL(GJ).GT.0) GO TO 50	E1 22250
	G=G+1	E1 22260
	HLIST(G,LEV)=GJ	E1 22270
	GLEVEL(GJ)=LEV	E1 22280
50	CONTINUE	E1 22290
	IF(G.EQ.0) RETURN	E1 22300
	LGLIST(LEV)=G	E1 22310
	DO 60 GG=1,G	E1 22320
	GJ=HLIST(GG,LEV)	E1 22330
	LIP=LIPRED(GJ)	E1 22340
	IF(LIP.EQ.0) GO TO 60	E1 22350
	DO 55 LP=1,LIP	E1 22360
	GP=IPRED(LP,GJ)	E1 22370
	OUTO(GP)=OUTO(GP)-1	E1 22380
55	CONTINUE	E1 22390
60	CONTINUE	E1 22400
	LEVM=LEV	E1 22410
	GO TO 45	E1 22420
		E1 22430
		E1 22440
		E1 22450
	ENTRY PVALUE	E1 22460
	DO 100 L=NN2,NRN2	E1 22470
	P\$(1,L)=1	E1 22480
100	CONTINUE	E1 22490
		E1 22500
		E1 22510
	LEV=LEVM	E1 22520
110	CONTINUE	E1 22530
	LO=LGLIST(LEV)	E1 22540
	DO 130 L=1,LO	E1 22550
	GI=HLIST(L,LEV)	E1 22560
	LIS=LISUCC(GI)	E1 22570
	BSGI=(GI-1)*N2	E1 22580
	LJTH=0	E1 22590
	DO 115 JTH=1,N2	E1 22600
	IF(P\$(1,BSGI+JTH).EQ.0) GO TO 115	E1 22610
	LJTH=LJTH+1	E1 22620
	X(LJTH)=JTH	E1 22630
115	CONTINUE	E1 22640
	IF(LJTH.EQ.0) GO TO 130	E1 22650
	DO 125 LS=1,LIS	E1 22660
	GS=ISUCC(LS,GI)	E1 22670
	BSGS=(GS-1)*N2	E1 22680
	DO 120 LJ=1,LJTH	E1 22690
	P\$(1,X(LJ)+BSGS)=0	E1 22700
120	CONTINUE	E1 22710

125	CONTINUE	E1	22720
130	CONTINUE	E1	22730
	LEV=LEV-1	E1	22740
	IF(LEV.GE.2) GO TO 110	E1	22750
	RETURN	E1	22760
		E1	22770
		E1	22780
		E1	22790
	ENTRY RSTRT(KEYRST)	E1	22800
	KEYRST=0	E1	22810
	IF(LEVM.GT.LMAX)GO TO 160	E1	22820
	DO 150 GT=N1,NP	E1	22830
	IF(LIPRED(GI).GT.FANIN)GO TO 160	E1	22840
	IF(LISUCG(GI).GT.FANOUT)GO TO 160	E1	22850
150	CONTINUE	E1	22860
	RETURN	E1	22870
160	KEYRST=1	E1	22880
	RETURN	E1	22890
	ENTRY UNNECE	E1	22900
C*****	THIS ENTRY DISCONNECT ALL GATES FROM WHICH THERE IS NO PATH	E1	22910
C	TO OUTPUT GATES *****	E1	22920
	TS=T	E1	22930
	DO 209 GI=NM1,NR	E1	22940
	IF(GLEVEL(GI).EQ.1) GO TO 207	E1	22950
	DO 205 GJ=N1,NM	E1	22960
	IF(SUC\$MX(GI,GJ).GT.0) GO TO 209	E1	22970
205	CONTINUE	E1	22980
C*****	GI IS REDUNDANT *****	E1	22990
207	CONTINUE	E1	23000
	LIP=LIPRED(GI)	E1	23010
	IF(LIP.EQ.0) GO TO 206	E1	23020
	DO 203 LI=1,LIP	E1	23030
	GK=IPRED(LI,GI)	E1	23040
	IF(INC\$MX(GK,GI).LE.0) GO TO 203	E1	23050
	T=T+1	E1	23060
	RTCONN(T)=100*GK+GI	E1	23070
	INC\$MX(GK,GI)=0	E1	23080
203	CONTINUE	E1	23090
206	LIS=LISUCG(GI)	E1	23100
	IF(LIS.EQ.0) GO TO 209	E1	23110
	DO 204 LI=1,LIS	E1	23120
	GK=ISUCG(LI,GI)	E1	23130
	IF(INC\$MX(GI,GK).LE.0) GO TO 204	E1	23140
	T=T+1	E1	23150
	RTCONN(T)=100*GI+GK	E1	23160
	INC\$MX(GI,GK)=0	E1	23170
204	CONTINUE	E1	23180
209	CONTINUE	E1	23190
	IF(T.GT.TS) GO TO 1	E1	23200
	RETURN	E1	23210
	END	E1	23220

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*****
P P P P   R R P R   O O O   G G G   R R P R   A   M   M
P   P   R   R   O   O   G   G   R   R   A   A   M M   M M
P   P   R   R   O   O   G   G   P R R R   A A A A   M   M   M
P   R   R   O   O   G   G   R   R   A   A   M   M   M
P   R   R   O O O   G G G   R   R   A   A   M   M

N   N   E E E E   T T T T   T T T T   R R R R   A   E E E E   2 2 2
N N   E   T   T   R   R   A   A   E   2   2
N N N   F   T   T   R   R   A   A   E   2
N   N N   E E E   T   T   R R R R   A A A A   X X X X   E E E   2
N   N   E   T   T   R   R   A   A   E   2
N   N   E E E E   T   T   P   R   A   A   E E E E   2 2 2 2
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SUBROUTINE MAIN                                E2 00010
EDITION BBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBB E2 00020
NOTE: ALL COMMON VARIABLES MIGHT NOT BE USED IN THIS PROGRAM. E2 00030
COMMON VARIABLES:                             E2 00040
$GT: POINTS TO A 'COLUMN' OF POTAB. FOR EACH 'ROW' THE ENTRY E2 00050
      IN THIS COL. TELLS GATE WHERE FN. IS REALIZED.           E2 00060
$LTH: POINTS TO A 'COLUMN' OF POTAB. FOR EACH 'ROW' THE ENTRY E2 00070
      IN THIS COL. TELLS HOW MANY CONNECTIONS MUST BE ADDED.   E2 00080
$NDE: POINTS TO A 'COLUMN' OF POTAB. FOR EACH 'ROW' THE ENTRY E2 00090
      IN THIS COL. TELLS THE NUMBER OF 1-ERRORS CREATED IF THIS E2 00100
      ROW IS USED.                                              E2 00110
$PW: POINTS TO A 'COLUMN' OF POTAB. FOR EACH 'ROW' THE ENTRY E2 00120
      IN THIS COLUMN TELLS THE PREFERENCE WEIGHT.               E2 00130
A: WEIGHT FOR NO. OF GATES IN COMPUTING COST FUNCTION.         E2 00140
B: WEIGHT FOR NO. OF CONNECTIONS IN COMPUTING COST FUNCTION.   E2 00150
COST: COST OF NETWORK - A MEASURE OF NETWORK SIZE.             E2 00160
ESSIS: RECORDS NO. OF ESSENTIAL 1'S IN EVERY INPUT TO CURRENT GCOE2 00170
      (POSITIONS IN ESSIS CORRES. TO GATES NOT FEEDING GCO ARE E2 00180
      IGNORED).                                                 E2 00190
F$UR1: POINTS TO LAST ELEMENT IN F$1.                          E2 00200
F$1: LISTS (CONSECUTIVELY) POSITIONS OF DESIRABLE 1'S (FOR E2 00210
      COVERING) IN A CONNECTIBLE FUNCTION.                      E2 00220
GI: LABEL OF A PARTICULAR GATE.                                E2 00230
GLEVEL: GLEVEL(GI) TELLS WHICH LEVEL OF THE NETWORK GI IS IN. E2 00240
GSMALL: STORES INTERMEDIATE AND FINAL CALCULATED CSPF'S.      E2 00250
HLIST: HLIST(I,J) GIVES NAME OF I-TH GATE (OR EX. VAR.) IN NET- E2 00260
      WORK LEVEL J.                                             E2 00270
IDX0: LIST OF 0-COORDINATES IN CSPFE OF THE GATE UNDER E2 00280
      CONSIDERATION.                                           E2 00290
IDX0E: LIST OF 0-ERROR-COORDINATES IN CSPFE OF THE GATE UNDER E2 00300
      CONSIDERATION.                                           E2 00310
IDX1: LIST OF 1-COORDINATES IN CSPFE OF THE GATE UNDER E2 00320
      CONSIDERATION.                                           E2 00330
IDX1E: LIST OF 1-ERROR-COORDINATES IN CSPFE OF THE GATE UNDER E2 00340

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C	CONSIDERATION.	E2	00370
C	IFLAG: SAME AS EYFFLG IN SUBROUTINE PROCII.	E2	00380
C	INC\$MX: INC\$MX(GI,GJ)>0 MEANS THERE EXISTS A CONNECTION FROM GATE	E2	00390
C	(OR EX. VAR.) GI TO GATE GJ. INC\$MX(GI,GJ)=0 IF NOT.	E2	00400
C	INPTCV: LISTS FOR EACH CORRESPONDING ENTRY OF F\$1, HOW MANY INPUTS	E2	00410
C	HAVE A '1' IN THE POSITION INDICATED BY F\$1.	E2	00420
C	IPATH: IPATH(GI)=1 MEANS GATE GI IS ON A PATH FROM A CERTAIN GATE	E2	00430
C	TO AN OUTPUT GATE. OTHERWISE IPATH(GI) = 0.	E2	00440
C	IPRED: IPRED(I,GJ) GIVES THE NAME OF THE I-TH GATE OR EX. VAR. IN	E2	00450
C	A LIST OF GATES AND EX. VAR. FEEDING GJ.	E2	00460
C	ISUCC: ISUCC(I,GJ) GIVES THE NAME OF THE I-TH GATE FED BY GJ.	E2	00470
C	JFLAG: SAME AS JAYFLG IN SUBROUTINE PROCII.	E2	00480
C	<FYA: A FLAG INDICATING IF ANY ERROR COMPENSATION HAS BEEN	E2	00490
C	PERFORMED.	E2	00500
C	KEY3: A FLAG INDICATING IF ANY PRIMARY D-ERROR-COORDINATES HAS	E2	00510
C	BEEN COMPENSATED.	E2	00520
C	KFLAG: SAME AS KEIFLG IN PROCII.	E2	00530
C	LFVM: NUMBER OF LEVELS IN THE NETWORK (NOTE EX. VAR. ARE ALSO	E2	00540
C	ASSIGNED LEVELS JUST LIKE GATES).	E2	00550
C	LGLIST: LGLIST(J) TELLS NO. OF GATES AND EX. VAR. IN LEVEL J OF	E2	00560
C	NETWORK.	E2	00570
C	LIP: NUMBER OF PREDECESSORS FOR THE GATE UNDER CONSIDERATION.	E2	00580
C	LIPRED: LIPRED(GI) TELLS NO. OF IMMEDIATE PREDECESSORS OF GATE GI.	E2	00590
C	LISTC: ORDERED LIST OF CONNECTIBLE INPUTS TO GCD. ORDERED BY	E2	00600
C	DECREASING NO. OF 0'S IN GCD COVERED.	E2	00610
C	LISTL: ORDERED LIST OF GATES AND EX. VAR. WHICH ORIGINALLY FED	E2	00620
C	GCD AND WHICH HAVE NOT YET BEEN DISCONNECTED. ORDERED BY	E2	00630
C	DECREASING NO. OF ESSENTIAL 1'S.	E2	00640
C	LISUCC: LISUCC(GI) TELLS NO. OF IMMEDIATE SUCCESSORS OF GATE (OR	E2	00650
C	EX. VAR.) GI.	E2	00660
C	LMTS2: UPPER LIMIT OF THE NUMBER OF ELEMENTS IN SET S2.	E2	00670
C	LPOTAB: FOR GATE GI, LPOTAB(GI) POINTS TO LAST ROW OF POTAB	E2	00680
C	CONCERNING GI.	E2	00690
C	M: NUMBER OF NETWORK OUTPUT GATES.	E2	00700
C	N: NUMBER OF EXTERNAL VARIABLES (OR INPUT FNC.) AVAILABLE.	E2	00710
C	NEPMAX: FOR ERROR COMPENSATION PROGRAMS. IF MORE THAN NEPMAX	E2	00720
C	ERROR POSITIONS OCCUR WHEN A PARTICULAR GATE IS REMOVED,	E2	00730
C	PROGRAM SKIPS ATTEMPT TO COMPENSATE FOR THAT GATE'S	E2	00740
C	REMOVAL. VALUE CAN BE SPECIFIED BY USER, OTHERWISE EQUAL	E2	00750
C	TO ONE HALF OF N2 BY DEFAULT.	E2	00760
C	NM: SUM OF N PLUS M	E2	00770
C	NM1: SUM OF NM PLUS 1.	E2	00780
C	NM2: PRODUCT OF N AND N2.	E2	00790
C	NOS: NUMBER OF ELEMENTS IN SET S.	E2	00800
C	NOS1: NUMBER OF ELEMENTS IN SET S1.	E2	00810
C	NOS1SV: NUMBER OF ELEMENTS IN SET S1 BEFORE ENTERING SUBROUTINE	E2	00820
C	RPLCF.	E2	00830
C	NOS2: NUMBER OF ELEMENTS IN SET S2.	E2	00840
C	NOT1: NUMBER OF ELEMENTS IN SET T1.	E2	00850
C	NOT1SV: NUMBER OF ELEMENTS IN SET T1 BEFORE ENTERING SUBROUTINE	E2	00860
C	RPLCF.	E2	00870
C	NOO: NUMBER OF ACTIVE ELEMENTS IN ARRAY IDXO.	E2	00880
C	NOOE: NUMBER OF ACTIVE ELEMENTS IN ARRAY IDXOE.	E2	00890
C	NO1: NUMBER OF ACTIVE ELEMENTS IN ARRAY IDX1.	E2	00900
C	NO1F: NUMBER OF ACTIVE ELEMENTS IN ARRAY IDX1E.	E2	00910
C	NR: SUM OF N PLUS R.	E2	00920
C	NRN2: PRODUCT OF NR AND N2.	E2	00930
C	NRPLC: NRPLC(I) STORES THE NUMBER OF ELEMENTS IN RPLC(I,*)	E2	00940
C	FOR I=1,2.	E2	00950
C	N1: SUM OF N PLUS 1.	E2	00960
C	N2: NUMBER OF DIFFERENT INPUT COMBINATIONS TO BE CONSIDERED	E2	00970



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C      (USUALLY 2 TO THE POWER N). E2 00980
C ORIGIN: ORIGIN(GI)=1 MEANS GI ORIGINALLY CONNECTED TO GCO. E2 00990
C ORIGIN(GI)=0 MEANS GI DID NOT FEED GCO ORIGINALLY. E2 01000
C P$: P$(1,-) CONSECUTIVELY LISTS OUTPUTS OF EVERY EX. VAR. AND E2 01010
C EVERY GATE (FOR EVERY INPUT COMBINATION): P$(1,1),..., E2 01020
C P$(1,N2) FOR FIRST EX VAR; P$(1,N2+1),...,P$(1,2*N2) FOR E2 01030
C SECOND EX VAR; ... ; P$(1,N*N2+1),..., P$(1,N*N2+N2) FOR E2 01040
C FIRST GATE; ETC. P$(2,-) IS USED AS WORK SPACE FOR E2 01050
C CALCULATIONS ASSOCIATED WITH P$(1,-). E2 01060
C PCO: FOR ERROR COMPENSATION PROCEDURES. PCO IS THE GATE E2 01070
C REMOVED FROM ORIGINAL NETWORK TO OBTAIN CURRENT ALTERED E2 01080
C NETWORK. E2 01090
C POINTA: NOT USED. E2 01100
C POINTC: POINTS TO LAST ELEMENT IN LISTC. E2 01110
C POINTL: POINTS TO LAST ELEMENT IN LISTL. E2 01120
C POINTR: POINTS TO LAST ELEMENT IN RNEC1 (IN SUBROUTINE SUBST1). E2 01130
C POTAB: POTENTIAL OUTPUT TABLE. HOLDS INFORMATION ABOUT ALL E2 01140
C COMBINATIONS OF CONNECTIONS TO FORM NEW (AND HOPEFULLY E2 01150
C USEFUL) FUNCTIONS. E2 01160
C PPOTAB: FOR GATE GI, PPOTAB(GI) POINTS TO FIRST OF A SEQUENCE OF E2 01170
C ROWS OF POTAB CONCERNING GI. E2 01180
C R: NUMBER OF GATES IN THE NETWORK (EXCLUDES EX VAR, ALSO E2 01190
C NOTE SOME OF R GATES MAY BE ISOLATED). E2 01200
C RPLC: RPLC(1,*) STORES THE SELECTED GATE'S IP GATES WHICH HAVE E2 01210
C ERROR-COORDINATES OF WEIGHT 2 OR ABOVE. E2 01220
C RPLC(2,*) STORES THE SELECTED GATE'S IP GATES WHICH HAVE E2 01230
C AT LEAST ONE ERROR-COORDINATE OF WEIGHT 1. E2 01240
C RSCONN: LIST OF CONNECTIONS ADDED TO A NETWORK (IN CODED FORM). E2 01250
C RTCONN: LIST OF CONNECTIONS REMOVED FROM A NETWORK (CODED FORM). E2 01260
C S: NO. OF CONNECTIONS ADDED TO A NETWORK. POINTS TO LAST E2 01270
C ENTRY IN RSCONN. E2 01280
C SETS: SET S CONSISTING OF INPUTS OF THE GATE UNDER CONSIDERATION E2 01290
C WHICH ARE TO BE REPLACED IF POSSIBLE. E2 01300
C SETS1: SET S1 CONSISTING OF ELEMENTS OF SET S WHICH CAN BE E2 01310
C REPLACED BY ELEMENTS IN SET S2. E2 01320
C SETS2: SET S2 CONSISTING OF FUNCTIONS WHICH ARE CANDIDATES FOR E2 01330
C REPLACING ELEMENTS IN SET S. E2 01340
C SETT1: SET T1 CONSISTING OF ESSENTIAL ONES COVERED BY ELEMENTS IN E2 01350
C SET S1. E2 01360
C STS: STARTING ELEMENT OF SET S. E2 01370
C SUC$MX: SUC$MX(GI,GJ)>0 MEANS GATE GJ IS A SUCCESSOR OF GATE GI. E2 01380
C SUC$MX(GI,GJ)=0 IF NOT. E2 01390
C SUMP: SUM OF ALL ACTIVE INPUTS OF THE GATE UNDER CONSIDERATION. E2 01400
C SUMS2: SUM OF ALL ACTIVE ELEMENTS OF SET S2. E2 01410
C T: NUMBER OF CONNECTIONS REMOVED FROM A NETWORK. POINTS TO E2 01420
C LAST ENTRY IN RTCONN. E2 01430
C TIME: USED TO STORE AMOUNT OF ELAPSED COMPUTATION TIME. E2 01440
C UNAME: MNEMONIC NAMES FOR EXTERNAL VARIABLES AND GATES. E2 01450
C VF$UB1: POINTS TO LAST ELEMENT IN VF$1. E2 01460
C VF$1: SIMILAR TO F$1, EXCEPT THIS LISTS JUST COMPONENT POSITIONS E2 01470
C (OF 0'S IN CSPF VECTOR OF GCO) COVERED ONLY BY REMAINING E2 01480
C ORIGINALLY CONNECTED INPUTS TO GCO. E2 01490
C E2 01500
C E2 01510
C E2 01520
C IMPLICIT INTEGER*4(A-T,V-Z,$), REAL(U) E2 01530
C COMMON NEPMAX E2 01540
C COMMON N , M , A , B E2 01550
C 1 , R , N2 , N1 , NR E2 01560
C 2 , NM , KFLAG , JFLAG , COST E2 01570
C 3 , LEVM , NRN2 , NM1 , NN2 E2 01580

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COMMON ISUCC(40,40) , LISUCC(40) , IPRED(40,40) , LIPRED(40) E2 01590
1 , INC$MX(40,40), SUP$MX(40,40), P$(2,1280) , UNAME(40) E2 01600
2 , GLEVEL(40) , LGLIST(40) , HLIST(40,40) , TIME E2 01610
COMMON T , RTCONN(100) , S , RSCONN(100) E2 01620
COMMON IFLAG , POINTA , FSSIS(40) , F$1(32) E2 01630
1 , F$UB1 , INPTCV(32) , LISTC(40) , POINTC E2 01640
2 , LISTL(40) , POINTL , ORIGIN(40) , IPATH(40) E2 01650
3 , POINTR , VF$1(32) , VF$UB1 , GSMALL(40,32) E2 01660
COMMON POTAB(200,42), PPOTAB(40) , LPOTAB(40) , NRPLC(2) E2 01670
1 , RPLC(2,40) , IDX0(32) , IDX0E(32) , IDX1(32) E2 01680
2 , IDX1E(32) , SUMP(32) , SETT1(32) , NOT1 E2 01690
3 , SETS1(40) , NOS1 , SETS(40) , NOS E2 01700
4 , STS , SUMS2(32) , SETS2(200) , NOS2 E2 01710
5 , LIP , NOOE , KEYA , KEYB E2 01720
6 , NOO , NO1 , NO1E , $GT E2 01730
7 , $LTH , $PW , $NDE , GI E2 01740
COMMON NOT1SV , NOS1SV , LMTS2 E2 01750
DIMENSION INTLIS(144),UGATE(40),UHEAD(20) E2 01760
DATA KOUNT5 /0/, UBLANK/' ' E2 01770
NEPMAX IS THE MAXIMUM ALLOWABLE NUMBER OF ERROR POSITIONS E2 01780
990 READ(5,1000,END=500) UHEAD, N, M, R, A, B, UC, NEPMAX E2 01790
1000 FORMAT(20A4/5I4,A4,I4) E2 01800
KEYXC=0 E2 01810
IF(UC.NE.UBLANK) KEYXC=1 E2 01820
CALL PAGE E2 01830
CALL LINE(10) E2 01840
KOUNT5=KOUNT5+1 E2 01850
PRINT 2, KOUNT5 E2 01860
2 FORMAT(20X,'*** OPTIMAL NOR NETWORK ***',50X,'PROBLEM NO.= ',I4 ) E2 01870
CALL LINE(4) E2 01880
PRINT 1005, UHEAD E2 01890
1005 FORMAT(25X,20A4) E2 01900
CALL LINE(4) E2 01910
PRINT 10, N,M,A,B E2 01920
10 FORMAT(30X,'NUMBER OF VARIABLES =',I4 // E2 01930
1 30X,'NUMBER OF FUNCTIONS =',I4 // E2 01940
2 30X,'COST COEFFICIENT A =',I4// E2 01950
3 47X, 'B =',I4) E2 01960
CALL LINE(1) E2 01970
IF(KEYXC.NE.0) GO TO 25 E2 01980
PRINT 21 E2 01990
21 FORMAT(1H0,29X,'--- UNCOMPLEMENTED VARIABLES X ---') E2 02000
GO TO 30 E2 02010
25 CONTINUE E2 02020
PRINT 28 E2 02030
28 FORMAT(1H0,29X,'--- BOTH COMPLEMENTED AND UNCOMPLEMENTED VARIABLES E2 02040
1 X, Y ---') E2 02050
30 CONTINUE E2 02060
CALL LINE(5) E2 02070
C***** SET UP EXTERNAL VARIABLES ***** E2 02080
N2=2**N E2 02090
IF(NEPMAX.EQ.0)NEPMAX = N2/2 E2 02100
H=N*N2 E2 02110
J=N2 E2 02120
L= 1 E2 02130
I=0 E2 02140
DO 1011 I=1,N E2 02150
J=J/2 E2 02160
L=L*2 E2 02170
SN= 1 E2 02180
DO 1010 LL=1,L E2 02190

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SN=-SN	E2 02200
V=(1+SN)/2	E2 02210
DO 1009 JJ=1,J	E2 02220
I=I+1	E2 02230
P\$(1,I)=V	E2 02240
IF (KEYXC.NE.0) P\$(1,I+H)=1-V	E2 02250
1009 CONTINUE	E2 02260
1010 CONTINUE	E2 02270
1011 CONTINUE	E2 02280
IF (KEYXC.NE.0) N=N+N	E2 02290
N1=N+1	E2 02300
NM=N+M	E2 02310
NM1=NM+1	E2 02320
NN2=V*N2+1	E2 02330
NR=N+R	E2 02340
NRN2=NR*N2	E2 02350
CALL OUTPUT(INC\$MX,KEYXC)	E2 02360
C***** READ IN NETWORK INFORMATION AND SET UP INC\$MX *****	E2 02370
READ 1001, CNTLIS	E2 02380
1001 FORMAT(16I5)	E2 02390
DO 1115 GI=1,NR	E2 02400
DO 1115 GJ=1,NR	E2 02410
1115 INC\$MX(GI,GJ)=0	E2 02420
DO 1120 I=1,144	E2 02430
ITEM=CNTLIS(I)	E2 02440
IF (ITEM.EQ.0) GO TO 1119	E2 02450
GI=ITEM/100	E2 02460
GJ=ITEM-100*GI	E2 02470
INC\$MX(GI,GJ)=1	E2 02480
GO TO 1120	E2 02490
1119 COST=A*R+B*(I-1)	E2 02500
GO TO 1130	E2 02510
1120 CONTINUE	E2 02520
1130 CONTINUE	E2 02530
CALL SUBNET	E2 02540
CALL PVALUE	E2 02550
CALL LINE(4)	E2 02560
PRINT 1140, COST	E2 02570
1140 FORMAT(20X, ' ORIGINAL NETWORK COST=', I5)	E2 02580
CALL LINE(4)	E2 02590
CALL TRUTH(P\$,1)	E2 02600
CALL LINE(4)	E2 02610
CALL CKT(INC\$MX,GLEVEL)	E2 02620
C	E2 02630
C***** ENTRY REDUNDANCY CHECK *****	E2 02640
S = 0	E2 02650
T = 0	E2 02660
CALL UNNECE	E2 02670
GATES = M	E2 02680
C = 0	E2 02690
DO 4 GI = 1,NR	E2 02700
C = C + LISUCC(GI)	E2 02710
IF (GI.LE.NM) GOTO 4	E2 02720
IF (LISUCC(GI).GT.0) GATES=GATES+1	E2 02730
4 CONTINUE	E2 02740
NEWCST = A*GATES + B*(C)	E2 02750
ITERTN = 0	E2 02760
3 OLDGST = NEWCST	E2 02770
ITERTN = ITERTN + 1	E2 02780
PRINT 2444,ITERTN	E2 02790
2444 FORMAT('1',5X,'***** BEGIN ',I3,'-TH APPLICATION OF PROCCE :	E2 02800



1	*****'/////////)	E2	02810
	T=0	E2	02820
	S=0	E2	02830
C	INITIALIZE TIMER TO 10 MINUTES	E2	02840
	CALL STIMEZ(60000)	E2	02850
	TIME = KTIMEZ(0)	E2	02860
C****	PROCEDURE COMPENSATE ERRORS *****	E2	02870
	CALL PROCCE(WORKED)	E2	02880
C	CALL FOR ELAPSED TIME	E2	02890
	TIME = KTIMEZ(0) - TIME	E2	02900
	CALL LINE(4)	E2	02910
	PRINT 3915	E2	02920
3916	FORMAT(20X,'TIME ELAPSED =',I8,' CENTISECONDS')	E2	02930
3915	FORMAT(20X,'NETWORK DERIVED BY PROCCE')	E2	02940
	PRINT 3916,TIME	E2	02950
	CALL LINE(4)	E2	02960
	CALL TRUTH(P\$,1)	E2	02970
	CALL LINE(4)	E2	02980
	CALL CKT(INC\$MX,GLEVEL)	E2	02990
	GATES = M	E2	03000
	C = 0	E2	03010
	DO 36 GI = 1,NR	E2	03020
	C = C + LISJCC(GI)	E2	03030
	IF(GI.LE.NM) GO TO 36	E2	03040
	IF(LISJCC(GI).GT.0) GATES = GATES + 1	E2	03050
36	CONTINUE	E2	03060
	NEWGST = A*GATES + B*C	E2	03070
	IF(NEWGST.LT.OLDGST)GO TO 37	E2	03080
	PRINT 105	E2	03090
105	FORMAT(1H ,10X,'NO REDUNDANCY FOUND.')	E2	03100
	GO TO 990	E2	03110
37	CALL LINE(3)	E2	03120
	PRINT 320,NEWGST	E2	03130
320	FORMAT(9X,'* A NETWORK DERIVED BY PROCCE'/9X,' COST=',I5,'.')	E2	03140
	IF(WORKED.EQ.1)GO TO 3	E2	03150
	CALL LINE(3)	E2	03160
	PRINT 301	E2	03170
301	FORMAT(9X,'* PROCCE CANNOT REDUCE THE PRECEDING NETWORK FURTHER')	E2	03180
	GO TO 990	E2	03190
500	STOP	E2	03200
	END	E2	03210

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*****
*
*           THE REST OF THE SUBROUTINES
*           REQUIRED FOR NETTRA-E2 ARE:
*
*   CALS1, CONECT, FORC, MINI2, ORDRQ2, OUTPUT,
*   POT, PROCCE, RCEC, RPLCF, AND SUBNET.
*
*
*   (THESE ARE IDENTICAL TO SUBROUTINES OF
*   THE SAME NAMES LISTED FOR NETTRA-E1.)
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P P P P R R R R O O O G G G R R R R A M M
P P R R O O G G R R A A M M M M
P P R R O O G G R R A A M M M M
P P R R O O G G R R A A M M M M
P R R O O G G R R A A M M M M

V V E E E E T T T T T T T R R R R A E E E E 3 3 3 3
N N E T T R R A A E 3
N N N E T T R R A A E 3
V V E E E T T R R R R A A A A A X X X X E E E 3
N N E T T R R A A A A E 3 3
N N E E E E T T R R A A E E E E 3 3 3
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SUBROUTINE MAIN E3 00010
EDITION AAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAA E3 00020
E3 00030
NOTE: ALL COMMON VARIABLES MIGHT NOT BE USED IN THIS PROGRAM. E3 00040
E3 00050
COMMON VARIABLES: E3 00060
$GT: POINTS TO A 'COLUMN' OF POTAB. FOR EACH 'ROW' THE ENTRY E3 00070
IN THIS COL. TELLS GATE WHERE FN. IS REALIZED. E3 00080
$LTH: POINTS TO A 'COLUMN' OF POTAB. FOR EACH 'ROW' THE ENTRY E3 00090
IN THIS COL. TELLS HOW MANY CONNECTIONS MUST BE ADDED. E3 00100
$NDE: POINTS TO A 'COLUMN' OF POTAB. FOR EACH 'ROW' THE ENTRY E3 00110
IN THIS COL. TELLS THE NUMBER OF 1-ERRORS CREATED IF THIS E3 00120
ROW IS USED. E3 00130
$PW: POINTS TO A 'COLUMN' OF POTAB. FOR EACH 'ROW' THE ENTRY E3 00140
IN THIS COLUMN TELLS THE PREFERENCE WEIGHT. E3 00150
A: WEIGHT FOR NO. OF GATES IN COMPUTING COST FUNCTION. E3 00160
B: WEIGHT FOR NO. OF CONNECTIONS IN COMPUTING COST FUNCTION. E3 00170
COST: COST OF NETWORK - A MEASURE OF NETWORK SIZE. E3 00180
ESS1S: RECORDS NO. OF ESSENTIAL 1'S IN EVERY INPUT TO CURRENT GCOE3 00190
(PPOSITIONS IN ESS1S CORRES. TO GATES NOT FEEDING GCD ARE E3 00200
IGNORED). E3 00210
F$UB1: POINTS TO LAST ELEMENT IN F$1. E3 00220
F$1: LISTS (CONSECUTIVELY) POSITIONS OF DESIRABLE 1'S (FOR E3 00230
COVERING) IN A CONNECTIBLE FUNCTION. E3 00240
GI: LABEL OF A PARTICULAR GATE. E3 00250
GLEVEL: GLEVEL(GI) TELLS WHICH LEVEL OF THE NETWORK GI IS IN. E3 00260
GSMALL: STORES INTERMEDIATE AND FINAL CALCULATED CSPF'S. E3 00270
HLIST: HLIST(I,J) GIVES NAME OF I-TH GATE (OR EX. VAR.) IN NET- E3 00280
WORK LEVEL J. E3 00290
IDX0: LIST OF 0-COORDINATES IN CSPFE OF THE GATE UNDER E3 00300
CONSIDERATION. E3 00310
IDX0E: LIST OF 0-ERROR-COORDINATES IN CSPFE OF THE GATE UNDER E3 00320
CONSIDERATION. E3 00330
IDX1: LIST OF 1-COORDINATES IN CSPFE OF THE GATE UNDER E3 00340
CONSIDERATION. E3 00350
IDX1E: LIST OF 1-ERROR-COORDINATES IN CSPFE OF THE GATE UNDER E3 00360

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C	CONSIDERATION.	E3 00370
C	IFLAG: SAME AS EYFFLG IN SUBROUTINE PROCII.	E3 00380
C	INC\$MX: INC\$MX(GI,GJ)>0 MEANS THERE EXISTS A CONNECTION FROM GATE	E3 00390
C	(OR EX. VAR.) GI TO GATE GJ. INC\$MX(GI,GJ)=0 IF NOT.	E3 00400
C	INPTCV: LISTS FOR EACH CORRESPONDING ENTRY OF F\$1, HOW MANY INPUTSE	E3 00410
C	HAVE A '1' IN THE POSITION INDICATED BY F\$1.	E3 00420
C	IPATH: IPATH(GI)=1 MEANS GATE GI IS ON A PATH FROM A CERTAIN GATE	E3 00430
C	TO AN OUTPUT GATE. OTHERWISE IPATH(GI) = 0.	E3 00440
C	IPRED: IPRED(I,GJ) GIVES THE NAME OF THE I-TH GATE OR EX. VAR. INE	E3 00450
C	A LIST OF GATES AND EX. VAR. FEEDING GJ.	E3 00460
C	ISUCC: ISUCC(I,GJ) GIVES THE NAME OF THE I-TH GATE FED BY GJ.	E3 00470
C	JFLAG: SAME AS JAYFLG IN SUBROUTINE PROCII.	E3 00480
C	KEYA: A FLAG INDICATING IF ANY ERROR COMPENSATION HAS BEEN	E3 00490
C	PERFORMED.	E3 00500
C	KEYB: A FLAG INDICATING IF ANY PRIMARY O-ERROR-COORDINATES HAS	E3 00510
C	BEEN COMPENSATED.	E3 00520
C	KFLAG: SAME AS KEIFLG IN PROCII.	E3 00530
C	LEVNM: NUMBER OF LEVELS IN THE NETWORK (NOTE EX. VAR. ARE ALSO	E3 00540
C	ASSIGNED LEVELS JUST LIKE GATES).	E3 00550
C	LGLIST: LGLIST(J) TELLS NO. OF GATES AND EX. VAR. IN LEVEL J OF	E3 00560
C	NETWORK.	E3 00570
C	LIP: NUMBER OF PREDECESSORS FOR THE GATE UNDER CONSIDERATION.	E3 00580
C	LIPRED: LIPRED(GI) TELLS NO. OF IMMEDIATE PREDECESSORS OF GATE GI.	E3 00590
C	LISTC: ORDERED LIST OF CONNECTIBLE INPUTS TO GCD. ORDERED BY	E3 00600
C	DECREASING NO. OF 0'S IN GCD COVERED.	E3 00610
C	LISTL: ORDERED LIST OF GATES AND EX. VAR. WHICH ORIGINALLY FED	E3 00620
C	GCD AND WHICH HAVE NOT YET BEEN DISCONNECTED. ORDERED BY	E3 00630
C	DECREASING NO. OF ESSENTIAL 1'S.	E3 00640
C	LISUCC: LISUCC(GI) TELLS NO. OF IMMEDIATE SUCCESSORS OF GATE (OR	E3 00650
C	EX. VAR.) GI.	E3 00660
C	LMTS2: UPPER LIMIT OF THE NUMBER OF ELEMENTS IN SET S2.	E3 00670
C	LPOTAB: FOR GATE GI, LPOTAB(GI) POINTS TO LAST ROW OF POTAB	E3 00680
C	CONCERNING GI.	E3 00690
C	M: NUMBER OF NETWORK OUTPUT GATES.	E3 00700
C	N: NUMBER OF EXTERNAL VARIABLES (OR INPUT FNC.) AVAILABLE.	E3 00710
C	NEPMAX: FOR ERROR COMPENSATION PROGRAMS. IF MORE THAN NEPMAX	E3 00720
C	ERROR POSITIONS OCCUR WHEN A PARTICULAR GATE IS REMOVED,	E3 00730
C	PROGRAM SKIPS ATTEMPT TO COMPENSATE FOR THAT GATE'S	E3 00740
C	REMOVAL. VALUE CAN BE SPECIFIED BY USER, OTHERWISE EQUAL	E3 00750
C	TO ONE HALF OF N2 BY DEFAULT.	E3 00760
C	NM: SUM OF N PLUS M	E3 00770
C	NM1: SUM OF NM PLUS 1.	E3 00780
C	NM2: PRODUCT OF N AND N2.	E3 00790
C	NDS: NUMBER OF ELEMENTS IN SET S.	E3 00800
C	NCS1: NUMBER OF ELEMENTS IN SET S1.	E3 00810
C	NDS1SV: NUMBER OF ELEMENTS IN SET S1 BEFORE ENTERING SUBROUTINE	E3 00820
C	RPLCF.	E3 00830
C	NCS2: NUMBER OF ELEMENTS IN SET S2.	E3 00840
C	NOT1: NUMBER OF ELEMENTS IN SET T1.	E3 00850
C	NOT1SV: NUMBER OF ELEMENTS IN SET T1 BEFORE ENTERING SUBROUTINE	E3 00860
C	RPLCF.	E3 00870
C	NDO: NUMBER OF ACTIVE ELEMENTS IN AFRAY IDX0.	E3 00880
C	NDOE: NUMBER OF ACTIVE ELEMENTS IN ARRAY IDXOE.	E3 00890
C	NDI: NUMBER OF ACTIVE ELEMENTS IN AFRAY IDX1.	E3 00900
C	NDOIE: NUMBER OF ACTIVE ELEMENTS IN AFRAY IDX1E.	E3 00910
C	NR: SUM OF N PLUS R.	E3 00920
C	NRN2: PRODUCT OF NR AND N2.	E3 00930
C	NRPLC: NRPLC(I) STORES THE NUMBER OF ELEMENTS IN RPLC(I,*)	E3 00940
C	FOR I=1,2.	E3 00950
C	N1: SUM OF N PLUS 1.	E3 00960
C	N2: NUMBER OF DIFFERENT INPUT COMBINATIONS TO BE CONSIDERED	E3 00970

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C      (USUALLY 2 TO THE POWER N). E3 00980
C ORIGIN: ORIGIN(GI)=1 MEANS GI ORIGINALLY CONNECTED TO GCO. E3 00990
C          ORIGIN(GI)=0 MEANS GI DID NOT FEED GCO ORIGINALLY. E3 01000
C      P$: P$(1,-) CONSECUTIVELY LISTS OUTPUTS OF EVERY EX. VAR. AND E3 01010
C          EVERY GATE (FOR EVERY INPUT COMBINATION): P$(1,1),..., E3 01020
C          P$(1,N2) FOR FIRST EX VAR; P$(1,N2+1),...,P$(1,2*N2) FOR E3 01030
C          SECOND EX VAR; ... ; P$(1,N*N2+1),..., P$(1,N*N2+N2) FOR E3 01040
C          FIRST GATE; ETC. P$(2,-) IS USED AS WORK SPACE FOR E3 01050
C          CALCULATIONS ASSOCIATED WITH P$(1,-). E3 01060
C      PCO: FOR ERROR COMPENSATION PROCEDURES. PCO IS THE GATE E3 01070
C          REMOVED FROM ORIGINAL NETWORK TO OBTAIN CURPENT ALTERED E3 01080
C          NETWORK. E3 01090
C      POINTA: NOT USED. E3 01100
C      POINTC: POINTS TO LAST ELEMENT IN LISTC. E3 01110
C      POINTL: POINTS TO LAST ELEMENT IN LISTL. E3 01120
C      POINTR: POINTS TO LAST ELEMENT IN RNEC1 (IN SUBROUTINE SUBST1). E3 01130
C      POTAB: POTENTIAL OUTPUT TABLE. HOLDS INFORMATION ABOUT ALL E3 01140
C          COMBINATIONS OF CONNECTIONS TO FORM NEW (AND HOPEFULLY E3 01150
C          USEFUL) FUNCTIONS. E3 01160
C      PPOTAB: FOR GATE GI, PPOTAB(GI) POINTS TO FIRST OF A SEQUENCE OF E3 01170
C          ROWS OF POTAB CONCERNING GI. E3 01180
C      R: NUMBER OF GATES IN THE NETWORK (EXCLUDES EX VAR, ALSO E3 01190
C          NOTE SOME OF R GATES MAY BE ISOLATED). E3 01200
C      RPLC: RPLC(1,*) STORES THE SELECTED GATE'S IP GATES WHICH HAVE E3 01210
C          ERROR-COORDINATES OF WEIGHT 2 OR ABOVE. E3 01220
C          RPLC(2,*) STORES THE SELECTED GATE'S IP GATES WHICH HAVE E3 01230
C          AT LEAST ONE ERROR-COORDINATE OF WEIGHT 1. E3 01240
C      RSCONN: LIST OF CONNECTIONS ADDED TO A NETWORK (IN CODED FORM). E3 01250
C      RTCONN: LIST OF CONNECTIONS REMOVED FROM A NETWORK (CODED FORM). E3 01260
C      S: NO. OF CONNECTIONS ADDED TO A NETWORK. POINTS TO LAST E3 01270
C          ENTRY IN RSCONN. E3 01280
C      SETS: SET S CONSISTING OF INPUTS OF THE GATE UNDER CONSIDERATION E3 01290
C          WHICH ARE TO BE REPLACED IF POSSIBLE. E3 01300
C      SFTS1: SET S1 CONSISTING OF ELEMENTS OF SET S WHICH CAN BE E3 01310
C          REPLACED BY ELEMENTS IN SET S2. E3 01320
C      SFTS2: SET S2 CONSISTING OF FUNCTIONS WHICH ARE CANDIDATES FOR E3 01330
C          REPLACING ELEMENTS IN SET S. E3 01340
C      SETT1: SET T1 CONSISTING OF ESSENTIAL ONES COVERED BY ELEMENTS IN E3 01350
C          SET S1. E3 01360
C      STS: STARTING ELEMENT OF SET S. E3 01370
C      SUC$MX: SUC$MX(GI,GJ)>0 MEANS GATE GJ IS A SUCCESSOR OF GATE GI. E3 01380
C          SUC$MX(GI,GJ)=0 IF NOT. E3 01390
C      SUMP: SUM OF ALL ACTIVE INPUTS OF THE GATE UNDER CONSIDERATION. E3 01400
C      SUMS2: SUM OF ALL ACTIVE ELEMENTS OF SET S2. E3 01410
C      T: NUMBER OF CONNECTIONS REMOVED FROM A NETWORK. POINTS TO E3 01420
C          LAST ENTRY IN RTCONN. E3 01430
C      TIME: USED TO STORE AMOUNT OF ELAPSED COMPUTATION TIME. E3 01440
C      JNAME: MNEMONIC NAMES FOR EXTERNAL VARIABLES AND GATES. E3 01450
C      VF$UB1: POINTS TO LAST ELEMENT IN VF$1. E3 01460
C      VF$1: SIMILAR TO F$1, EXCEPT THIS LISTS JUST COMPONENT POSITIONS E3 01470
C          (OF 0'S IN CSPF VECTOR OF GCO) COVERED ONLY BY REMAINING E3 01480
C          ORIGINALLY CONNECTED INPUTS TO GCO. E3 01490
C          E3 01500
C          E3 01510
C          E3 01520
C      IMPLICIT INTEGER*4(A-T,V-Z,$), REAL(U) E3 01530
C      COMMON NEPMAX E3 01540
C      COMMON      N          , M          , A          , B          E3 01550
C      1          , R          , N2         , N1         , NR         E3 01560
C      2          , NM         , KFLAG      , JFLAG      , COST       E3 01570
C      3          , LEV4       , NRN2       , NM1        , NN2       E3 01580

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COMMON  ISUCC(40,40) , LISUCC(40) , IPRED(40,40) , LIPRED(40) E3 01590
1      , INC$MX(40,40), SUC$MX(40,40), P$(2,1280) , UNAME(40) E3 01600
2      , GLEVL(40) , LGLIST(40) , HLIST(40,40) , TIME E3 01610
COMMON  T , RTCONN(100) , S , RSCONN(100) E3 01620
COMMON  IFLAG , POINTA , ESSIS(40) , F$1(32) E3 01630
1      , F$UR1 , INPTCV(32) , LISTC(40) , POINTC E3 01640
2      , LISTL(40) , POINTL , CRIGIN(40) , IPATH(40) E3 01650
3      , PCINTR , VF$1(32) , VF$UB1 , G$SMALL(40,32) E3 01660
COMMON  POTAB(200,42), PPOTAB(40) , LPOTAB(40) , NRPLC(2) E3 01670
1      , RPLC(2,40) , IDX0(32) , IDX0E(32) , IDX1(32) E3 01680
2      , IDX1E(32) , SUMP(32) , SETT1(32) , NOT1 E3 01690
3      , SETS1(40) , NOS1 , SETS(40) , NOS E3 01700
4      , STS , SUMS2(32) , SETS2(200) , NOS2 E3 01710
5      , LIP , NOOF , KEYA , KEYB E3 01720
6      , NOO , NO1 , NO1E , $GT E3 01730
7      , $LTH , $PW , $NCE , GI E3 01740
COMMON  NOT1SV , NOS1SV , LMTS2 E3 01750
DIMENSION ENTLIS(144),UGATE(40),UHEAD(20) E3 01760
DATA KOUNT5 /0/, UBLANK/' ' E3 01770
990 READ(5,1000,END=500) UHEAD, N, M, R, A, B, UC, NEPMAX E3 01780
C NEPMAX IS THE MAXIMUM ALLOWABLE NUMBER OF ERROR POSITIONS E3 01790
1000 FORMAT(20A4/5I4,A4,I4) E3 01800
KEYXC=0 E3 01810
IF(UC.NE.UBLANK) KEYXC=1 E3 01820
CALL PAGE E3 01830
CALL LINE(10) E3 01840
KOUNT5=KOUNT5+1 E3 01850
PRINT 2, KOUNT5 E3 01860
2 FORMAT(20X,'*** OPTIMAL NOR NETWORK ***',50X,'PROBLEM NO.= ',I4 ) E3 01870
CALL LINE(4) E3 01880
PRINT 1005, UHEAD E3 01890
1005 FORMAT(25X,20A4) E3 01900
CALL LINE(4) E3 01910
PRINT 10, N,M,A,B E3 01920
10 FORMAT(30X,'NUMBER OF VARIABLES =',I4 // E3 01930
1      30X,'NUMBER OF FUNCTIONS =',I4 // E3 01940
2      30X,'COST COEFFICIENT A =',I4// E3 01950
3      47X, 'B =',I4) E3 01960
CALL LINE(1) E3 01970
IF(KEYXC.NE.0) GO TO 25 E3 01980
PRINT 21 E3 01990
21 FORMAT(1H0,29X,'--- UNCOMPLEMENTED VARIABLES X ---') E3 02000
GO TO 30 E3 02010
25 CONTINUE E3 02020
PRINT 28 E3 02030
28 FORMAT(1H0,29X,'--- BOTH COMPLEMENTED AND UNCOMPLEMENTED VAPIABLES E3 02040
1 X, Y ---') E3 02050
30 CONTINUE E3 02060
CALL LINE(5) E3 02070
C***** SET JP EXTERNAL VARIABLES ***** E3 02080
N2=2**N E3 02090
IF(NEPMAX.EQ.0)NEPMAX = N2/2 E3 02100
H=N*N2 E3 02110
J=N2 E3 02120
L= 1 E3 02130
I=0 E3 02140
DO 1011 II=1,N E3 02150
J=J/2 E3 02160
L=L*2 E3 02170
$V= 1 E3 02180
DO 1010 LL=1,L E3 02190

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SN=-SN	E3 02200
V=(1+SN)/2	E3 02210
DO 1009 JJ=1,J	E3 02220
I=I+1	E3 02230
PS(1,I)=V	E3 02240
IF(KEYXC.NE.0)PS(1,I+1)=1-V	E3 02250
1009 CONTINUE	E3 02260
1010 CONTINUE	E3 02270
1011 CONTINUE	E3 02280
IF(KEYXC.NE.0) N=N+N	E3 02290
N1=N+1	E3 02300
NM=N+1	E3 02310
NM1=N+1	E3 02320
NN2=N*N2+1	E3 02330
NR=N+R	E3 02340
NRN2=NR*N2	E3 02350
CALL OUTPUT(INC\$MX,KEYXC)	E3 02360
C***** READ IN NETWORK INFORMATION AND SET UP INC\$MX *****	E3 02370
READ 1001, CNTLIS	E3 02380
1001 FORMAT(16I5)	E3 02390
DO 1115 GI=1,NR	E3 02400
DO 1115 GJ=1,NR	E3 02410
1115 INC\$MX(GI,GJ)=0	E3 02420
DO 1120 I=1,144	E3 02430
ITEM=CNTLIS(I)	E3 02440
IF(ITEM.EQ.0) GO TO 1119	E3 02450
GI=ITEM/100	E3 02460
GJ=ITEM-100*GI	E3 02470
INC\$MX(GI,GJ)=1	E3 02480
GO TO 1120	E3 02490
1119 COST=A*R+B*(I-1)	E3 02500
GO TO 1130	E3 02510
1120 CONTINUE	E3 02520
1130 CONTINUE	E3 02530
CALL SUBNET	E3 02540
CALL PVALUE	E3 02550
CALL LINE(4)	E3 02560
PRINT 1140, COST	E3 02570
1140 FORMAT(20X,' ORIGINAL NETWORK COST=', I5)	E3 02580
CALL LINE(4)	E3 02590
CALL TRUTH(PS,1)	E3 02600
CALL LINE(4)	E3 02610
CALL CKT(INC\$MX,GLEVEL)	E3 02620
C	E3 02630
C***** ENTRY REDUNDANCY CHECK *****	E3 02640
T=0	E3 02650
S=0	E3 02660
CALL UNNECE	E3 02670
C INITIALIZE TIMER TO 10 MINUTES	E3 02680
CALL STIMEZ(60000)	E3 02690
TIME = KTIMEZ(0)	E3 02700
C***** PROCEDURE COMPENSATE ERRORS *****	E3 02710
CALL ALPATH	E3 02720
C CALL FOR ELAPSED TIME	E3 02730
TIME = KTIMEZ(0) - TIME	E3 02740
CALL LINE(4)	E3 02750
PRINT 3915	E3 02760
3916 FORMAT(20X,'TIME ELAPSED =',I8,' CENTISECONDS')	E3 02770
3915 FORMAT(20X,'NETWORKS DERIVED BY ALL-PATH PROCCE')	E3 02780
PRINT 3916,TIME	E3 02790
GO TO 990	E3 02800

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C**** FOR THIS SPECIAL ALL-PATH VERSION OF MAIN, THE REST OF THE PROGRAME3 02810
C ST.S ARE BYPASSED E3 02820
CALL LINE(4) E3 02830
CALL TRUTH(P$,1) E3 02840
CALL LINE(4) E3 02850
CALL CKT(INC$MX,GLEVEL) E3 02860
C***** PRINT OUT NETWORK DERIVED BY REDUNDANCY CHECK ***** E3 02870
IF(T.GT.0) GO TO 110 E3 02880
CALL LINE(2) E3 02890
102 PRINT 105 E3 02900
105 FORMAT(1H ,10X,'NO REDUNDANCY FOUND.') E3 02910
GO TO 990 E3 02920
C E3 02930
110 CONTINUE E3 02940
C***** PRINT OUT REDUNDANT GATES ***** E3 02950
CALL UNNECE E3 02960
G=0 E3 02970
115 KEY=0 E3 02980
DO 125 GJ=NM1,NR E3 02990
IF(LISUCC(GJ).GT.0) GO TO 125 E3 03000
G=G+1 E3 03010
UGATE(G)=UNAME(GJ) E3 03020
LISUCC(GJ)=9999 E3 03030
LIP=LIPRED(GJ) E3 03040
IF(LIP.EQ.0) GO TO 125 E3 03050
DO 120 LP=1,LIP E3 03060
GP=IPRED(LP,GJ) E3 03070
T=T+1 E3 03080
INC$MX(GP,GJ)=0 E3 03090
RTCONN(T)=100*GP+GJ E3 03100
LISUCC(GP)=LISUCC(GP)-1 E3 03110
KEY=1 E3 03120
120 CONTINUE E3 03130
125 CONTINUE E3 03140
IF(KEY.GT.0) GO TO 115 E3 03150
CALL SURNET E3 03160
CALL PVALUE E3 03170
CALL LINE(3) E3 03180
301 PRINT 302 E3 03190
302 FORMAT(1H ,10X,'THE FOLLOWING RECONFIGUFATION DONE.'//) E3 03200
IF(G.EQ.0) GO TO 310 E3 03210
PRINT 303,( UGATE(GG),GG=1,G) E3 03220
303 FORMAT(1H ,15X,'REDUNDANT GATE(S) '//20X,10(3X,A3)) E3 03230
CALL LINE(2) E3 03240
C***** PRINT OUT REMOVED AND ADDED CONNECTIONS ***** E3 03250
310 IF(T.EQ.0)GOTO401 E3 03260
PRINT 311 E3 03270
311 FORMAT(1H ,15X,'REMOVED CONNECTION(S)') E3 03280
DO 315 TT=1,T E3 03290
ITEM=RTCONN(TT) E3 03300
GI=ITEM/100 E3 03310
GJ=ITEM-GI*100 E3 03320
UI=UNAME(GI) E3 03330
UJ=JNAME(GJ) E3 03340
PRINT 314,UI,UJ E3 03350
314 FORMAT(1H0,19X,'(' ,2X,A3,', ' ,2X,A3,')') E3 03360
315 CONTINUE E3 03370
401 IF(S.EQ.0) GO TO 319 E3 03380
CALL LINE(2) E3 03390
PRINT 316 E3 03400
316 FORMAT(1H ,15X,'ADDED CONNECTION(S)') E3 03410

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DD 318 SS=1,S	E3 03420
ITEM=RSCONN(SS)	E3 03430
GI=ITEM/100	E3 03440
GJ=ITEM-GI*100	E3 03450
JI=JNAME(GI)	E3 03460
UJ=JNAME(GJ)	E3 03470
PRINT 314,UI,UJ	E3 03480
318 CONTINUE	E3 03490
319 CONTINUE	E3 03500
COUNTC = 0	E3 03510
DD 6447 I = 1,NR	E3 03520
6447 COUNTC = COUNTC + LISUCC(I)	E3 03530
CJNEW = A * (R - G) + B * (COUNTC)	E3 03540
COST = CUNEW	E3 03550
CALL LINE(3)	E3 03560
PRINT 320, CUNEW	E3 03570
320 FORMAT(9X,'* A NETWORK DERIVED BY PROCCE'/9X,' COST=',I5,'.')	E3 03580
NEWCOST = CUNEW	E3 03590
IF(NEWCOST.LT.OLD COST)GO TO 3	E3 03600
GO TO 990	E3 03610
500 STOP	E3 03620
END	E3 03630
	E3 03640
SUBROUTINE PROCCE(WORKED)	E3 03650
PROCCE FOR MULTI-PATH PROGRAM *****	E3 03660
EDITION AAA	E3 03670
IF PROCCE SUCCESSFULLY COMPENSATES ERRORS, 'WORKED' IS SET TO 1, O	E3 03680
'WORKED' IS SET TO 0	E3 03690
	E3 03700
DEFS. OF MOST 'COMMON' VARIABLES CAN BE FOUND IN MAIN PROGRAM.	E3 03710
	E3 03720
SPECIAL COMMON VARIABLES:	E3 03730
	E3 03740
INC\$AV: STACK CONTAINING BLOCKS. EACH BLOCK CONTAINS A NETWORK'S	E3 03750
NAME, PARENT, COST, AND LIST OF CONNECTIONS.	E3 03760
NAME: NAME OF NETWORK UNDER CONSIDERATION.	E3 03770
NETLST: LIST (STACK) OF POINTERS TO TOP OF BLOCKS IN INC\$AV.	E3 03780
NLSPT: POINTS TO TOP OF STACK NETLST.	E3 03790
NTCNTR: NETWORK COUNTER - COUNTS NUMBER OF NETWORKS GENERATED SO	E3 03800
FAR.	E3 03810
NTCOST: COST OF A PARTICULAR NETWORK.	E3 03820
PARENT: NAME OF NETWORK FROM WHICH NETWORK 'NAME' WAS DERIVED.	E3 03830
\$AVPT: POINTS TO FIRST FREE LOCATION IN STACK INC\$AV.	E3 03840
	E3 03850
VARIABLE DEFINITIONS:	E3 03860
EP: EP(I)=1 MEANS AT LEAST ONE NETWORK OUTPUT GATE HAS AN	E3 03870
ERRONEOUS OUTPUT IN THE I-TH COMPONENT WHEN PCO IS REMOVED	E3 03880
FROM THE NETWORK. EP(I)=0 OTHERWISE.	E3 03890
ERRORS: TOTAL NO. OF ERRORS IN NETWORK OUTPUTS WHEN PCO REMOVED.	E3 03900
GATES: NUMBER OF GATES REMOVED FROM NETWORK BY CALL TO MINI2.	E3 03910
IMPROV: A PARAMETER RETURNED BY MINI2. '=1' MEANS MINI2 WAS ABLE	E3 03920
TO REDUCE COST OF NETWORK.	E3 03930
MAX: MAXIMUM NUMBER OF REQUIRED 1'S IN A CSPF VECTOR (AFTER	E3 03940
CALLING MINI2) PLUS 1.	E3 03950
MIN: ORIGINALLY SET TO ZERO, MIN IS INCREMENTED EACH TIME BY 1	E3 03960
UNTIL ITS VALUE EQUALS MAX.	E3 03970
NEP: NO. OF ERROR POSITIONS FOR A GIVEN NETWORK AFTER A SE-	E3 03980
LECTED GATE HAS BEEN REMOVED. AN ERROR POSITION IS A	E3 03990
COMPONENT POSITION WHICH IS IN ERROR FOR AT LEAST ONE	E3 04000

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C      OUTPUT. E3 04010
C      NEPMAX: READ FROM INPUT CARDS, THIS PARAMETER IS PASSED TO PROCCE E3 04020
C      WHEN IT IS CALLED BY MAIN. IT REPRESENTS THE MAXIMUM E3 04030
C      ALLOWABLE NUMBER OF ERROR POSITIONS. IF AN ALTERED (I.E., E3 04040
C      SOME PCO REMOVED) NETWORK EXCEEDS THIS MAXIMUM, ERROR E3 04050
C      COMPENSATION IS NOT ATTEMPTED FOR THAT NETWORK. E3 04060
C      NETOUT: STORES OUTPUTS OF GATES IN ALTERED (PCO REMOVED) NETWORK. E3 04070
C      ONECNT: USED IN COUNTING NO. OF 1'S IN CSPF VECTOR OF A GATE. E3 04080
C      ONES: AFTER THE INITIAL CALCULATION OF THE CSPF SETS IN THE E3 04090
C      BEGINNING, ONES(GI) GIVES THE NUMBER OF 1'S IN THE CSPF E3 04100
C      VECTOR OF GI. THIS INFORMATION IS REQUIRED FOR GENERATING E3 04110
C      PORDER. E3 04120
C      ORGOUT: USED TO STORE ORIGINAL (UNALTERED) NETWORK OUTPUTS IN E3 04130
C      CODED FORM (SAME CODE AS IN GSMALL) AND (40,32) FORMAT. E3 04140
C      PCO: CURRENT GATE REMOVED FROM ORIGINAL NETWORK TO OBTAIN E3 04150
C      CURRENT ALTERED NETWORK. PCO = PORDER(PCOUNT). E3 04160
C      PCOUNT: A POINTER TO PORDER. E3 04170
C      PORDER: ORDERING OF GATES ACCORDING TO NUMBER OF 1'S IN THEIR E3 04180
C      CSPF VECTORS. GATES ARE INDIVIDUALLY REMOVED FROM ORIGI- E3 04190
C      NAL NETWORK IN THIS ORDER E3 04200
C      PSUB: USED AS A POINTER TO PORDER DURING ITS INITIALIZATION. E3 04210
C      QINC$M: STORES A COPY OF INC$MX FOR THE ORIGINAL NETWORK. E3 04220
C      START: POINTS TO BEGINNING OF LIST OF NETWORK OUTPUTS IN P$. E3 04230
C      STOP: POINTS TO END OF LIST OF NETWORK OUTPUTS IN P$. E3 04240
C      E3 04250
C      I,J,NI,X,Y ARE USED AS JUST TEMPORARY VARIABLES. E3 04260
C      E3 04270
C      HOW TO INCREASE CAPACITY OF SUBROUTINE. E3 04280
C      DIMENSION: PORDER(X) E3 04290
C      ONES(X) E3 04300
C      QINC$M(X,X) - X EQUAL TO MAX NO. OF GATES PLUS EX. VAR. E3 04310
C      EP(Y) - Y EQUAL TO: 2** (MAX ALLOWED NO OF EX VAR) E3 04320
C      NETOUT(X,Y) E3 04330
C      ORGOUT(X,Y) - X,Y AS ABOVE E3 04340
C      E3 04350
C      IMPLICIT INTEGER*4(A-T,V-Z,$), REAL(U) E3 04360
C      COMMON NEPMAX E3 04370
C      COMMON N , M , A , B E3 04380
C      1 , R , N2 , N1 , NR E3 04390
C      2 , NM , KFLAG , JFLAG , COST E3 04400
C      3 , LEVM , NRN2 , NM1 , NN2 E3 04410
C      COMMON ISJCC(40,40) , LISUCC(40) , IPRED(40,40) , LIPRED(40) E3 04420
C      1 , INC$MX(40,40) , SUC$MX(40,40) , P$(2,1280) , UNAME(40) E3 04430
C      2 , GLEVEL(40) , LGLIST(40) , HLIST(40,40) , TIME E3 04440
C      COMMON T , RTCONN(100) , S , RSCONN(100) E3 04450
C      COMMON IFLAG , POINTA , ESSIS(40) , F$1(32) E3 04460
C      1 , F$UB1 , INPTCV(32) , LISTC(40) , POINTC E3 04470
C      2 , LISTL(40) , POINTL , ORIGIN(40) , IPATH(40) E3 04480
C      3 , POINTR , VF$1(32) , VF$UB1 , GSMALL(40,32) E3 04490
C      COMMON POTAB(200,42) , PPOTAB(40) , LPOTAB(40) , NRPLC(2) E3 04500
C      1 , RPLC(2,40) , IDX0(32) , IDX0E(32) , IDX1(32) E3 04510
C      2 , IDX1E(32) , SUMP(32) , SETT1(32) , NOT1 E3 04520
C      3 , SETS1(40) , NOS1 , SETS(40) , NOS E3 04530
C      4 , STS , SUMS2(32) , SETS2(200) , NOS2 E3 04540
C      5 , LIP , NOOE , KEYA , KEYB E3 04550
C      6 , NOO , NO1 , NO1E , $GT E3 04560
C      7 , $LTH , $PW , $NCE , GI E3 04570
C      COMMON NOT1SV , NOS1SV , LMTS2 E3 04580
C      COMMON NTCNTR , PARENT , NAME , INC$AV(10000) E3 04590
C      1 , $AVPT , NETLST(500) , NLSTPT , NTCOST E3 04600
C      DIMENSION PORDER(40),ONES(40),QINC$M(40,40),NETOUT(40,32), E3 04610

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	1	FP(32),CRGOUT(40,32)	E3	04620
C		THIS SUBROUTINE ASSUMES ALL ARRAYS ARE UPDATED	E3	04630
C		PREVIOUS TO BEING CALLED	E3	04640
C			E3	04650
		\$GT = 33	E3	04660
		\$LTH = 34	E3	04670
		\$PW = 41	E3	04680
		\$NDE = 42	E3	04690
		WORKED = 0	E3	04700
		S = 0	E3	04710
		T = 0	E3	04720
		\$PSAVE = \$AVPT	E3	04730
C			E3	04740
C		BLOCK B	E3	04750
C			E3	04760
		CALL MINI2(IMPROV)	E3	04770
C		IN THIS CALL TO MINI2, GORDER WILL BE CALCULATED. GORDER WILL BE	E3	04780
C		LATER IN EACH CALL TO INITGS (AN ENTRY POINT OF MINI2). NOTE THAT	E3	04790
C		IS NOT AFFECTED BY THE REMOVAL OF GATES FROM THE ORIGINAL NETWORK.	E3	04800
		IF(IMPROV.EQ.0)GO TO 1	E3	04810
		GAFTER = M	E3	04820
		C = 0	E3	04830
		DO 2 I = 1,NR	E3	04840
		C = C + LISUCC(I)	E3	04850
		IF(I.LE.NM)GOTO2	E3	04860
		IF(LISUCC(I).GT.0) GAFTER = GAFTER + 1	E3	04870
	2	CONTINUE	E3	04880
		GREFOR = (INTCOST-B*(C+T))/A	E3	04890
		GATES = R - GAFTER	E3	04900
		PRINT 4,GATES,T	E3	04910
	4	FORMAT(' ',15,' GATES AND',13,' CONNECTIONS HAVE BEEN REMOVED FROM	E3	04920
	1	THE NETWORK DURING THE INITIAL CALCULATION OF THE CSPF SET')	E3	04930
	1	CONTINUE	E3	04940
C		COUNT THE NUMBER OF 1'S IN THE CSPF VECTOR FOR EACH GATE	E3	04950
		MAX = 0	E3	04960
		DO 5 I = N1,NP	E3	04970
		ONECNT = 0	E3	04980
		DO 6 J = 1,N2	E3	04990
		IF(GSMALL(I,J).LE.0)GO TO 6	E3	05000
		ONECNT = ONECNT + 1	E3	05010
	6	CONTINUE	E3	05020
		IF(ONECNT.GT.MAX)MAX=ONECNT	E3	05030
		ONES(I) = ONECNT	E3	05040
	5	CONTINUE	E3	05050
		MAX = MAX + 1	E3	05060
		MIN = -1	E3	05070
		PSUB = 1	E3	05080
	7	MIN = MIN + 1	E3	05090
		IF(MIN.EQ.MAX) GO TO 8	E3	05100
		DO 9 I = N1,NP	E3	05110
		IF(ONES(I).NE.MIN)GO TO 9	E3	05120
		PORDER(PSUB) = I	E3	05130
		PSUB = PSUB + 1	E3	05140
	9	CONTINUE	E3	05150
		GOTO7	E3	05160
	8	CONTINUE	E3	05170
C		SAVE ORIGINAL NETWORK	E3	05180
		DO 10 I = 1,NR	E3	05190
		DO 10 J = 1,NR	E3	05200
		QINC\$(I,J) = INC\$(MX(I,J))	E3	05210
	10	CONTINUE	E3	05220

C	SAVE ORIGINAL OUTPUTS	E3 05230
C	SAVE ORIGINAL OUTPUTS IN (2,1280) FORMAT	E3 05240
	START = (N*N2) + 1	E3 05250
	STOP = (NM*N2)	E3 05260
	DO 13 I = START, STOP	E3 05270
	PS(2,I) = PS(1,I)	E3 05280
13	CONTINUE	E3 05290
C	SAVE ORIGINAL OUTPUTS IN CODED (40,32) FORMAT	E3 05300
	DO 27 I = N1,NM	E3 05310
	X = ( I-1) * N2	E3 05320
	DO 28 J = 1,N2	E3 05330
	Y = PS(1,X+J)	E3 05340
	IF(Y)30,31,32	E3 05350
C	COMPONENT IS DON'T CARE (I.E., -1)	E3 05360
30	ORROUT(I,J) = 0	E3 05370
	GOTO 28	E3 05380
C	COMPONENT IS LOGICAL ZERO	E3 05390
31	ORROUT(I,J) = -100	E3 05400
	GO TO 28	E3 05410
C	COMPONENT IS LOGICAL ONE	E3 05420
32	ORROUT(I,J) = 1	E3 05430
28	CONTINUE	E3 05440
27	CONTINUE	E3 05450
C		E3 05460
C	BLOCK C C C C C C C C C C C C C C C C C C	E3 05470
C		E3 05480
	PCOUNT = 0	E3 05490
11	PCOUNT = PCOUNT + 1	E3 05500
	IF(PCOUNT.GT. R)GOTO42	E3 05510
	PCO = PORDER(PCOUNT)	E3 05520
	IF(ONES(PCO).EQ.0)GO TO 11	E3 05530
	IF(PCO.LE.NM)GO TO 11	E3 05540
C	ERRORS UNCORRECTABLE, RESTORE NETWORK, TRY AGAIN	E3 05550
	DO 19 I = 1,NR	E3 05560
	DO 19 J = 1,NP	E3 05570
	INC\$MX(I,J) = QINC\$M(I,J)	E3 05580
19	CONTINUE	E3 05590
C	REMOVE GATE PCO FROM THE NETWORK	E3 05600
	DO 12 I = 1,NR	E3 05610
	IF(INC\$MX(I,PCO).EQ.0)GO TO 34	E3 05620
	INC\$MX(I,PCO) = 0	E3 05630
34	IF(INC\$MX(PCO,I).EQ.0) GO TO 12	E3 05640
	INC\$MX(PCO,I) = 0	E3 05650
12	CONTINUE	E3 05660
C	UPDATE GATE OUTPUTS FOR ALTERED NETWORK	E3 05670
C		E3 05680
C	BLOCK D D D D D D D D D D D D D D D D D D	E3 05690
C		E3 05700
33	CALL SUBNET	E3 05710
	CALL PVALUE	E3 05720
	CALL UNNECE	E3 05730
	S = 0	E3 05740
	T = 0	E3 05750
C	RESTORE GSMALL FOR OUTPUT GATES	E3 05760
	DO 29 I = N1,NM	E3 05770
	DO 29 J = 1, N2	E3 05780
	GSMALL(I,J) = ORROUT(I,J)	E3 05790
29	CONTINUE	E3 05800
	ERRORS = 0	E3 05810
	DO 24 I=1,N2	E3 05820
24	EP(I) = 0	E3 05830

	DO 14 I = 1,M	E3 05840
	NI = N + I	E3 05850
	X = (NI - 1) * M2	E3 05860
	DO 15 J = 1,M2	E3 05870
	IF(GSMALL(NI,J))16,15,17	E3 05880
C	CASE WHERE REQUIREMENT IS A ZERO	E3 05890
16	IF(P\$(1,X+J).EQ.0)GO TO 15	E3 05900
C	CASE OF ONE WITH ERROR	E3 05910
	GSMALL(NI,J) = 1001	E3 05920
	ERRORS = ERRORS + 1	E3 05930
	EP(J) = 1	E3 05940
	GO TO 15	E3 05950
C	CASE WHERE REQUIREMENT IS A ONE	E3 05960
17	IF(P\$(1,X+J).EQ.1)GO TO 15	E3 05970
C	CASE OF ZERO WITH ERROR	E3 05980
	GSMALL(NI,J) = -1100	E3 05990
	ERRORS = ERRORS + 1	E3 06000
	EP(J) = 1	E3 06010
15	CONTINUE	E3 06020
14	CONTINUE	E3 06030
	IF(ERRORS.EQ.0)WORKED = 1	E3 06040
	IF(ERRORS.EQ.0) GO TO 23	E3 06050
	NEP = 0	E3 06060
	DO 25 I = 1,M2	E3 06070
	IF(EP(I).EQ.0) GO TO 25	E3 06080
	NEP = NEP + 1	E3 06090
25	CONTINUE	E3 06100
	IF(NEP.GT.NEPMAX) GO TO 11	E3 06110
C		E3 06120
C	BLOCK E E E F E E E E E E E E E E E E E E E E	E3 06130
C		E3 06140
	CALL POT	E3 06150
C	'POT' IS A SUBROUTINE THAT GENERATES THE POTENTIAL OUTPUT TABLE	E3 06160
C		E3 06170
C	BLOCK F	E3 06180
C		E3 06190
C	SAVE NEW NETWORK OUTPUTS	E3 06200
	DO 18 J = 1,M2	E3 06210
	DO 18 I = N1,NM	E3 06220
	NETOUT(I,J) = GSMALL(I,J)	E3 06230
18	CONTINUE	E3 06240
	CALL FORMGO	E3 06250
	CALL INITGS	E3 06260
	CALL RDEC(811,833)	E3 06270
C	NEW NETWORK HAS BEEN FOUND, PUT IT IN STACK	E3 06280
23	NTNTR = NTNTR + 1	E3 06290
	TEMP = \$AVPT	E3 06300
	DO 26 I=1,NR	E3 06310
	DO 26 J=N1,NR	E3 06320
	IF(INC\$MX(I,J).LE.0)GO TO 26	E3 06330
	X = 1000*I + J	E3 06340
	INC\$AV(\$AVPT) = X	E3 06350
	\$AVPT = \$AVPT + 1	E3 06360
26	CONTINUE	E3 06370
	NLSTPT = NLSTPT + 1	E3 06380
	NETLST(NLSTPT) = \$AVPT + 2	E3 06390
C	INSERT COST OF NETWORK	E3 06400
	GATES = M	E3 06410
	DO 41 GATEI = NML,NR	E3 06420
	IF(LISUCG(GATEI).GT.0)GATES=GATES+1	E3 06430
41	CONTINUE	E3 06440



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C = $AVPT - TEMP
CST = A*GATES + B*C
INC$AV($AVPT) = CST
INSERT PARENT'S NAME
INC$AV($AVPT+1) = NAME
INSERT THIS NEW NETWORK'S NAME INTO STACK
INC$AV($AVPT+2) = NTCNTR
$AVPT = $AVPT + 3
PRINT 2345,NTCNTR,NAME
2345 FORMAT('1',19X,'NETWORK NUMBER ',I4,' DERIVED BY PROCCE. THE PAREE
INT OF THIS NETWORK IS NUMBER ',I4/////))
CALL LINE(4)
CALL TRUTH(P$,1)
CALL LINE(4)
CALL CKT(INC$MX,GLEVEL)
S = 0
T = 0
CALL LINE(4)
PRINT 2346,CST
2346 FORMAT(20X,'THIS NETWORK HAS A COST OF:',I5)
GO TO 11
42 IF(IMPROV.EQ.0.OR.$PSAVE.NE.$AVPT)RETURN
IF(GBEFOR.EQ.GAFTER)RETURN
IF HERE, AN IMPROVED NETWORK OF FEWER GATES HAS BEEN FOUND, BUT ONE
THE OPERATION OF MINIZ (I.E., NOT BY RCEC). SO NEW NETWORK WILL BE
PRINTED OUT, BUT NOT STORED IN THE STACK.
RESTORE NETWORK DERIVED BY MINIZ
DO 43 I=1,NR
DO 43 J=1,NR
INC$MX(I,J) = QINC$M(I,J)
43 CONTINUE
CALL SUBNET
CALL PVALUE
NTCNTR = NTCNTR + 1
PRINT 2348,NTCNTR,NAME
2348 FORMAT('1',19X,'NETWORK NUMBER ',I4,' DERIVED BY MINIZ. THE PAREE
IT OF THIS NETWORK IS NUMBER ',I4/////))
CALL LINE(4)
CALL TRUTH(P$,1)
CALL LINE(4)
CALL CKT(INC$MX,GLEVEL)
CST = A*GAFTER + B*C
PRINT 2346,CST
RETURN
END

SUBROUTINE ALPATH
EDITIOY AAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAA
DEFINITIONS OF MOST 'COMMON' VARIABLES CAN BE FOUND IN MAIN PROG.
SPECIAL COMMON VARIABLES:
INC$AV: STACK CONTAINING BLOCKS. EACH BLOCK CONTAINS A NETWORK'S
NAME, PARENT, COST, AND LIST OF CONNECTIONS.
NAME: NAME OF NETWORK UNDER CONSIDERATION.
NETLST: LIST (STACK) OF POINTERS TO TOP OF BLOCKS IN INC$AV.
NLSTPT: POINTS TO TOP OF STACK NETLST.

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C      NTCNTR: NETWORK COUNTER - COUNTS NUMBER OF NETWORKS GENERATED SO E3 07040
C      FAR. E3 07050
C      NTCOST: COST OF A PARTICULAR NETWORK. E3 07060
C      PARENT: NAME OF NETWORK FROM WHICH NETWORK 'NAME' WAS DERIVED. E3 07070
C      $AVPT: POINTS TO FIRST FREE LOCATION IN STACK INC$AV. E3 07080
C E3 07090
C      HOW TO INCREASE CAPACITY OF SUBROUTINE. E3 07100
C      DIMENSION: NETLST(X) - X EQUAL TO MAX NUMBER OF NETWORKS ALLOWED E3 07110
C      IN STACK + 1 E3 07120
C      INC$AV(Y) - Y EQUAL TO 3*X PLUS MAX TOTAL NUMBER OF E3 07130
C      CONNECTIONS IN ALL NETWORKS STORED IN STACK E3 07140
C E3 07150
C      IMPLICIT INTEGER*4(A-T,V-Z,$), REAL(U) E3 07160
C      COMMON NEPMAX E3 07170
C      COMMON N , M , A , B E3 07180
C      1 , R , N2 , N1 , NR E3 07190
C      2 , NM , KFLAG , JFLAG , CCOST E3 07200
C      3 , LEVM , NRN2 , NM1 , NN2 E3 07210
C      COMMON ISJCC(40,40) , LISJCC(40) , IPRED(40,40) , LIPRED(40) E3 07220
C      1 , INC$MX(40,40) , SUC$MX(40,40) , P$(2,1280) , UNAME(40) E3 07230
C      2 , GLEVEL(40) , LGLIST(40) , HLIST(40,40) , TIME E3 07240
C      COMMON T , RTCONN(100) , S , RSCONN(100) E3 07250
C      COMMON IFLAG , POINTA , ESSIS(40) , F$1(32) E3 07260
C      1 , F$UB1 , INPTCV(32) , LISTC(40) , POINTC E3 07270
C      2 , LISTL(40) , POINTL , ORIGIN(40) , IPATH(40) E3 07280
C      3 , POINTR , VF$1(32) , VF$UB1 , GSMALL(40,32) E3 07290
C      COMMON POTAB(200,42) , PPOTAB(40) , LPOTAB(40) , NRPLC(2) E3 07300
C      1 , RPLC(2,40) , IDX0(32) , IDX0E(32) , IDX1(32) E3 07310
C      2 , IDX1E(32) , SUMP(32) , SETT1(32) , NOT1 E3 07320
C      3 , SETS1(40) , NOS1 , SETS(40) , NOS E3 07330
C      4 , STS , SUMS2(32) , SETS2(200) , NOS2 E3 07340
C      5 , LIP , NOOE , KEYA , KEYB E3 07350
C      6 , NOO , NOI , NOIE , $GT E3 07360
C      7 , $LTH , $PW , $NOE , GI E3 07370
C      COMMON NOT1SV , NOS1SV , LMTS2 E3 07380
C      COMMON NTCNTR , PARENT , NAME , INC$AV(10000) E3 07390
C      1 , $AVPT , NETLST(500) , NLSTPT , NTCOST E3 07400
C      NETLST(1) = 0 E3 07410
C      NTCNTR = 1 E3 07420
C      NLSTPT = 1 E3 07430
C      $AVPT = 1 E3 07440
C      PARENT = 0 E3 07450
C      NAME = 1 E3 07460
C      NTCOST = COST E3 07470
C      1 CALL PROCDE(WORKED) E3 07480
C      IF STACK EMPTY, RETURN E3 07490
C      IF(NETLST(NLSTPT).EQ.0)RETURN E3 07500
C      CHOOSE NEW NETWORK FROM TOP OF STACK, SET UP E3 07510
C      DO 2 I=1,NR E3 07520
C      DO 2 J=1,NR E3 07530
C      2 INC$MX(I,J) = 0 E3 07540
C      X = NETLST(NLSTPT) E3 07550
C      NLSTPT = NLSTPT - 1 E3 07560
C      Y = NETLST(NLSTPT) + 1 E3 07570
C      $AVPT = Y E3 07580
C      NAME = INC$AV(X) E3 07590
C      PARENT = INC$AV(X-1) E3 07600
C      NTCOST = INC$AV(X-2) E3 07610
C      X = X - 3 E3 07620
C      DO 3 I=Y,X E3 07630
C      Z = INC$AV(I) E3 07640

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IGATE = Z/1000
JGATE = Z - 1000*IGATE
3 INC$MX(IGATE,JGATE) = 1
CALL SUBNET
CALL PVALUE
GO TO 1
END

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E3 07650
E3 07660
E3 07670
E3 07680
E3 07690
E3 07700
E3 07710

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*****
*
*           THE REST OF THE SUBROUTINES
*           REQUIRED FOR NETTRA-E3 ARE:
*
*   CALS1, CONECT, FORC, MINI2, OPDRQ2, OUTPUT,
*           POT, RCEC, RPLCF, AND SUBNET.
*
*
*   (THESE ARE IDENTICAL TO SUBROUTINES OF
*   THE SAME NAMES LISTED FOR NETTRA-E1.)
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16. Abstracts <p>Three NOR network transduction procedures based on error-compensation were implemented in the FORTRAN computer programs NETTRA-E1, NETTRA-E2, and NETTRA-E3. The general principles on which these programs are based are discussed in a separate report. The present report, however, describes the specific implementations of the three programs and serves as a reference manual for the program user. Preparation of input data is discussed in detail.</p> <p>Transduction (<u>transformation and reduction</u>) procedures attempt to reduce given, non-optimal, multiple-output, multiple-level, loop-free, NOR-gate networks to "near-optimal" networks of fewer gates. The three programs described in this report, based on the sophisticated "error-compensation" concept, remove gates one at a time from the network and, after each removal, try to reconfigure the network, without employing additional gates, to compensate for any resultant errors caused in the network output(s).</p>			
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